A FORWARD-BACKWARD VIEW OF SOME PRIMAL-DUAL OPTIMIZATION METHODS IN IMAGE RECOVERY

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ABSTRACT

A wide array of image recovery problems can be abstracted into the problem of minimizing a sum of composite convex functions in a Hilbert space. To solve such problems, primal-dual proximal approaches have been developed which provide efficient solutions to large-scale optimization problems. The objective of this paper is to show that a number of existing algorithms can be derived from a general form of the forward-backward algorithm applied in a suitable product space. Our approach also allows us to develop useful extensions of existing algorithms by introducing a variable metric. An illustration to image restoration is provided.

Index Terms— convex optimization, duality, parallel computing, proximal algorithm, variational methods, image recovery.

1. INTRODUCTION

Many image recovery problems can be formulated in Hilbert spaces $\mathcal{H}$ and $(\mathcal{G}_i)_{1 \leq i \leq m}$ as structured optimization problems of the form

$$\sum_{i=1}^m g_i(L_i x),$$

where, for every $i \in \{1, \ldots, m\}$, $g_i$ is a proper lower semi-continuous convex function from $\mathcal{G}_i$ to $[-\infty, +\infty]$ and $L_i$ is a bounded linear operator from $\mathcal{H}$ to $\mathcal{G}_i$. For example, the functions $(g_i \circ L_i)_{1 \leq i \leq m}$ may model data fidelity terms, smooth or nonsmooth measures of regularity, or hard constraints on the solution. In recent years, many algorithms have been developed to solve such a problem by taking advantage of recent advances in convex optimization, especially in the development of proximal tools (see [9] and the references therein). In image processing, however, solving such a problem still poses a number of conceptual and numerical challenges. First of all, one often looks for methods which have the ability to split the problem by activating each of the functions through elementary processing steps which can be computed in parallel. This makes it possible to reduce the complexity of the original problem and to benefit from existing parallel computing architectures. Secondly, it is often useful to design algorithms which can exploit, in a flexible manner, the structure of the problem. In particular, some of the functions may be Lipschitz differentiable in which case they should be exploited through their gradient rather than through their proximity operator, which is usually harder to implement (examples of proximity operators with closed-form expression can be found in [6, 12]). In some problems, the functions $(g_i)_{1 \leq i \leq m}$ can be expressed as the infimal convolution of simpler functions (see [9] and the references therein). Last but not least, in image recovery, the operators $(L_i)_{1 \leq i \leq m}$ may be of very large size so that their inversions are costly (e.g., in reconstruction problems). Finding algorithms which do not require to perform inversions of these operators is thus of paramount importance.

Note that all the existing convex optimization algorithms do not have these desirable properties. For example, the Alternating Direction Method of Multipliers (ADMM) [18, 17, 20] requires a stringent assumption of invertibility of the involved linear operator. Parallel versions of ADMM [28] and related Parallel Proximal Algorithm (PPXA) [11, 25] usually necessitate a linear inversion to be performed at each iteration. Also, early primal-dual algorithms [4, 5, 7, 10, 16, 21] did not make it possible to handle smooth functions through their gradients. Only recently, have primal-dual methods been proposed with this feature. Such a work was initiated in [13] in the line of [4] and subsequent developments can be found in [2, 3, 8, 9, 15, 27, 30]. As will be seen in the present paper, another advantage of these approaches is that they can be coupled with variable metric strategies which can potentially accelerate their convergence.

In Section 2, we provide some background on convex analysis and monotone operator theory. In Section 3, we introduce a general form of the forward-backward algorithm which uses a variable metric. This algorithm is employed in Section 4 to develop a versatile family of primal-dual proximal methods. Several particular instances of this framework are discussed. Finally, we provide illustrative numerical results in Section 5.

2. NOTATION AND BACKGROUND

Monotone operator theory [1] provides a both insightful and elegant framework for dealing with convex optimization problems and developing new solution algorithms that could not be devised using purely variational tools. We summarize a number of related concepts that will be needed.

Throughout, $\mathcal{H}$, $\mathcal{G}$, and $(\mathcal{G}_i)_{1 \leq i \leq m}$ are real Hilbert spaces. We denote the scalar product of a Hilbert space by $(\cdot \mid \cdot)$ and the associated norm by $\| \cdot \|$. The symbol $\rightharpoonup$ denotes weak convergence, and $\Id$ denotes the identity operator. We denote by $\mathcal{B}(\mathcal{H}, \mathcal{G})$ the space of bounded linear operators from $\mathcal{H}$ to $\mathcal{G}$, we set $S(\mathcal{H}) = \{ \mathcal{H} \to \mathcal{H} \}$.

This work was supported by the CNRS MASTODONS project (grant 2013 MesureHD).

1In a finite dimensional space, weak convergence is equivalent to strong convergence.
\[ \{ L \in B(H, H) \mid L = L^* \} \], where \( L^* \) denotes the adjoint of \( L \).

The Loewner partial ordering on \( \mathcal{S}(H) \) is denoted by \( \succeq \). For every \( \alpha \in [0, +\infty) \), we set \( \mathcal{P}_\alpha(H) = \{ U \in \mathcal{S}(H) \mid U \succeq \alpha \text{Id} \} \), and we denote by \( \sqrt{U} \) the square root of \( U \in \mathcal{P}_\alpha(H) \). Moreover, for every \( U \in \mathcal{P}_\alpha(H) \) and \( \alpha > 0 \), we define the norm \( \|x\|_U = \sqrt{U(x, x)} \).

We denote by \( \mathcal{G} = G_1 \oplus \cdots \oplus G_m \) the Hilbert direct sum of the Hilbert spaces \( (G_i)_{i \in \mathcal{E}} \), i.e., their product space equipped with the scalar product \( (x, y) \mapsto \sum_{i=1}^m (x_i, y_i) \) where \( x = (x_i)_{i \in \mathcal{E}} \) and \( y = (y_i)_{i \in \mathcal{E}} \) denote generic elements in \( \mathcal{G} \).

Let \( A : H \to 2^H \) be a set-valued operator. We denote by \( \text{gra} A = \{ (x, u) \in H \times H \mid u \in Ax \} \) the graph of \( A \), by \( \text{zer} A = \{ x \in H \mid 0 \in Ax \} \) the set of zeros of \( A \), and by \( \text{ran} A = \{ u \in H \mid \exists x \in H \text{ such that } u \in Ax \} \) its range. The inverse of \( A \) is \( A^{-1} : H \to 2^H : u \mapsto \{ x \in H \mid u \in Ax \} \), and the resolvent of \( A \) is \( J_A = (\text{Id} + A)^{-1} \). Moreover, \( A \) is monotone if
\[ (\forall (x, y) \in H \times H)(\forall(u, v) \in Ax \times Ay) \quad \langle x - y \mid u - v \rangle \geq 0, \] and maximally monotone if it is monotone and there exists no monotone operator \( B : H \to 2^H \) such that \( \text{gra} A \subseteq \text{gra} B \) and \( A \neq B \).

An operator \( B : H \to H \) is \( \beta \)-coercive for some \( \beta \in [0, +\infty) \) if
\[ (\forall y \in H)(\forall(\gamma > 0)) \quad \langle \gamma y - x \mid Bx - By \rangle \geq \beta \gamma^2 \|Bx - By\|^2. \]

The conjugate of a function \( f : H \to [-\infty, +\infty) \) is
\[ f^* : H \to [-\infty, +\infty) : u \mapsto \sup_{x \in H} \{ \langle x \mid u \rangle - f(x) \}, \]
and the infimal convolution of \( f \) with \( g : H \to [-\infty, +\infty) \) is
\[ f \boxdot g : H \to [-\infty, +\infty) : x \mapsto \inf_{y \in H} \{ f(y) + g(x - y) \}. \]

The class of lower semicontinuous functions \( f : H \to [-\infty, +\infty] \) such that \( \text{dom} f = \{ x \in H \mid f(x) < +\infty \} = \emptyset \) is denoted by \( \Gamma_0(H) \). If \( f \in \Gamma_0(H) \), then \( f^* \in \Gamma_0(H) \) and the subdifferential of \( f \) is the maximally monotone operator
\[ \partial f : H \to 2^H : x \mapsto \{ u \in H \mid (\forall y \in H) \langle y - x \mid u \rangle + f(x) \leq f(y) \}. \]

Let \( U \in \mathcal{P}_\alpha(H) \) for some \( \alpha \in [0, +\infty] \). The proximity operator of \( f \in \Gamma_0(H) \) relative to the metric induced by \( U \) is [22, Section XV.4]
\[ \text{prox}_U f : H \to H : x \mapsto \text{argmin}_{y \in H} f(y) + \frac{1}{2} \|x - y\|^2_U. \]

When \( U = \text{Id} \), we retrieve the standard definition of the proximity operator [1, 24]. Let \( C \) be a nonempty subset of \( H \). The indicator function of \( C \) is defined on \( H \) as
\[ 1_C : x \mapsto \begin{cases} 0, & \text{if } x \in C; \\ +\infty, & \text{if } x \notin C. \end{cases} \]

Finally, \( \ell_2^+ (\mathbb{N}) \) denotes the set of summable sequences in \([0, +\infty)\).

3. A GENERAL FORM OF FORWARD-BACKWARD ALGORITHM

Optimization problems can often be reduced to finding a zero of a sum of two maximally monotone operators \( A \) and \( B \) acting on \( H \). When \( B \) is \( \beta \)-coercive (see (3)), a useful algorithm to solve this problem is the forward-backward algorithm, which can be formulated in a general form involving a variable metric as shown in the next result.

\[ \text{Theorem 3.1} \quad \text{Let } \alpha \in [0, +\infty], \beta \in [0, +\infty], \text{let } A : H \to 2^H \text{ be maximally monotone, and let } B : H \to H \text{ be } \beta \text{-coercive. Let } (\gamma_n)_{n \in \mathbb{N}} \in \ell_2^+ (\mathbb{N}) \text{ and let } (V_n)_{n \in \mathbb{N}} \text{ be a sequence in } \mathcal{P}_\alpha(H) \text{ such that}
\]
\[ \left\{ \begin{array}{l}
\sup_{n \in \mathbb{N}} \|V_n\| < +\infty \\
(\forall n \in \mathbb{N}) \quad (1 + \eta_n) V_{n+1} \succeq V_n
\end{array} \right. \]
Note that in the special case when $\ell_i = \epsilon_{(0)}$, $g_i \square \ell_i$ reduces to $g_i$ in (12).
Let us now examine how Problem 4.1 can be reformulated from the standpoint of monotone operators. To this end, let us define $g \in \Gamma_0(\mathcal{G})$, $\ell \in \Gamma_0(\mathcal{G})$ and $L \in \mathcal{B}(\mathcal{H}, \mathcal{G})$ by
\[
g(v) = \sum_{i=1}^{m} g_i(v), \quad \ell(v) = \sum_{i=1}^{m} \ell_i(v) \quad \text{and} \quad L(x) = (L_1x, \ldots, L_mx).
\]
Let us now introduce the product space $\mathcal{K} = \mathcal{H} \oplus \mathcal{G}$ and the operators
\[
A : \mathcal{K} \rightarrow 2^\mathcal{K} \quad \text{(14)}
\]
and
\[
B : \mathcal{K} \rightarrow \mathcal{K} \quad \text{(15)}
\]
The operator $A$ can be shown to be maximally monotone, whereas $B$ is cocoercive. A key observation in this context is that, if there exists $(\mathcal{X}, \mathcal{T}) \in \mathcal{K}$ such that $(\mathcal{X}, \mathcal{T}) \in \text{zer}(A + B)$, then $(\mathcal{X}, \mathcal{T})$ is a pair of primal-dual solutions to Problem 4.1 [13]. This connection with the construction for a zero of $A + B$ makes it possible to apply a forward-backward algorithm as discussed in Section 3, by using a linear operator $V_n \in \mathcal{B}(\mathcal{K}, \mathcal{K})$ to change the metric at each iteration $n$. Depending on the form of this operator various algorithms can be obtained.

4.2. A first class of primal-dual algorithms

Let $\alpha \in (0, +\infty)$, let $(U_n)_{n \in \mathbb{N}}$ be a sequence in $\mathcal{P}_\alpha(\mathcal{H})$ such that $(\forall n \in \mathbb{N}) \ U_n+1 \geq U_n$. For every $i \in \{1, \ldots, m\}$, let $(U_i(n))_{n \in \mathbb{N}}$ be a sequence in $\mathcal{P}_\alpha(\mathcal{G}_i)$ such that $(\forall n \in \mathbb{N}) \ U_{i,n+1} \geq U_{i,n}$. A first possible choice for $(V_n)_{n \in \mathbb{N}}$ is given by
\[
(V_n)_{n \in \mathbb{N}}^{-1} : (x, v) \mapsto (U_n^{-1}x - L^*v, -Lx + \tilde{U}_n^{-1}v) \quad \text{(17)}
\]
where
\[
\tilde{U}_n : \mathcal{G} \rightarrow \mathcal{G} : (v_1, \ldots, v_m) \mapsto (U_{1,n}v_1, \ldots, U_{m,n}v_m). \quad \text{(18)}
\]
The following result constitutes a direct extension of [14, Example 6.4]:

Proposition 4.2 Let $x_0 \in \mathcal{H}$, and let $(a_n)_{n \in \mathbb{N}}$ and $(c_n)_{n \in \mathbb{N}}$ be absolutely summable sequences in $\mathcal{H}$. For every $i \in \{1, \ldots, m\}$, let $\nu_i, 0 \in \mathcal{G}_i$, let $(b_i(n))_{n \in \mathbb{N}}$ and $(d_i(n))_{n \in \mathbb{N}}$ be absolutely summable sequences in $\mathcal{G}_i$. For every $n \in \mathbb{N}$, let $\mu_n \in (0, +\infty)$ be a Lipschitz constant of $U_n^{1/2} \circ \nabla h \circ U_n^{1/2}$ and, for every $i \in \{1, \ldots, m\}$, let $\nu_i, n \in (0, +\infty)$ be a Lipschitz constant of $U_i^{1/2,n} \circ \nabla \ell_i \circ U_i^{1/2,n}$. Let $(\lambda_n)_{n \in \mathbb{N}}$ be a sequence in $[0, 1]$ such that $\inf_{n \in \mathbb{N}} \lambda_n > 0$. For every $n \in \mathbb{N}$, set
\[
\delta_n = \left( \sum_{i=1}^{m} \|U_i^{-1}L_i\sqrt{U_n}\|^2 \right)^{-1/2} - 1, \quad \text{(19)}
\]
and suppose that
\[
\inf_{n \in \mathbb{N}} (1 + \delta_n) \max_{\mu \in \mathcal{G}_i} (\nu_1, n, \ldots, \nu_m,n) > \frac{1}{2} \quad \text{(20)}
\]
Set
\[
\ell_n = \begin{cases} 
1 & \text{if } 0 = 1, \ldots, 0, \\
\|x_n - U_n\| & \text{if } f(x_n) \geq 0, \\
0 & \text{otherwise}.
\end{cases}
\]
Then $(x_n)_{n \in \mathbb{N}}$ converges weakly to a solution to (12), for every $i \in \{1, \ldots, m\}$, $(v_i(n))_{n \in \mathbb{N}}$ converges weakly to some $v_i \in \mathcal{G}_i$, and $(\nu_1, \ldots, \nu_m) \in \mathcal{H}$ is a solution to (13).

In the special case when $U_n \equiv \tau I_d$, $\tau \in [0, +\infty)$ and, for every $i \in \{1, \ldots, m\}$, $U_i(n) \equiv \sigma_i I_d$ with $\sigma_i \in [0, +\infty)$, we recover the parallel algorithm proposed in [30]. Variants of this algorithm where, for every $i \in \{1, \ldots, m\}$, $\ell_i = \epsilon_{(0)}$ are also investigated in [15]. In this case, less restrictive assumptions on the choice of $(\tau, \sigma_1, \ldots, \sigma_m)$ can be made. Note that this algorithm itself can be viewed as a generalization of the algorithm which constitutes the main topic of [5, 16, 21] (designated by some authors as PDHG). A preconditioned version of this algorithm was proposed in [25] corresponding to the case when $m = 1$, $(\forall n \in \mathbb{N}) U_n$ and $U_1$ are constant matrices, and no error term is taken into account. Algorithm (21) when, for every $n \in \mathbb{N}$, $\lambda_n \equiv 1$, $U_n$ and $(U_i(n))_{i \in \mathbb{N}}$ are diagonal matrices, $h = 0$, and $(\forall i \in \{1, \ldots, m\}) \ e_i = \epsilon_{(0)}$ appears also to be closely related to the adaptive method in [19].

4.3. A second class of primal-dual algorithms

Let $\alpha \in [0, +\infty)$, let $(U_n)_{n \in \mathbb{N}}$ be a sequence in $\mathcal{P}_\alpha(\mathcal{H})$ such that $(\forall n \in \mathbb{N}) \ U_{n+1} \geq U_n$. For every $i \in \{1, \ldots, m\}$, let $(U_i(n))_{n \in \mathbb{N}}$ be a sequence in $\mathcal{P}_\alpha(\mathcal{G}_i)$ such that $(\forall n \in \mathbb{N}) \ U_{i,n+1} \geq U_{i,n}$. A second possible choice for $(V_n)_{n \in \mathbb{N}}$ is given by the following diagonal form:
\[
(V_n)_{n \in \mathbb{N}}^{-1} : (x, v) \mapsto (U_n^{-1}x, (\tilde{U}_n^{-1}L_nU_n^*)v) \quad \text{(22)}
\]
where $\tilde{U}_n$ is given by (18).

The following result can then be deduced from Theorem 3.1. Its proof is skipped due to the lack of space.

Proposition 4.3 Let $x_0 \in \mathcal{H}$, and let $(a_n)_{n \in \mathbb{N}}$ be an absolutely summable sequence in $\mathcal{H}$. For every $i \in \{1, \ldots, m\}$, let $\nu_i, 0 \in \mathcal{G}_i$, let $(b_i(n))_{n \in \mathbb{N}}$ and $(d_i(n))_{n \in \mathbb{N}}$ be absolutely summable sequences in $\mathcal{G}_i$. For every $n \in \mathbb{N}$, let $\mu_n \in (0, +\infty)$ be a Lipschitz constant of $U_n^{1/2} \circ \nabla h \circ U_n^{1/2}$ and, for every $i \in \{1, \ldots, m\}$, let $\nu_i, n \in (0, +\infty)$ be a Lipschitz constant of $U_i^{1/2,n} \circ \nabla \ell_i \circ U_i^{1/2,n}$. Let $(\lambda_n)_{n \in \mathbb{N}}$ be a sequence in $[0, 1]$ such that $\inf_{n \in \mathbb{N}} \lambda_n > 0$. For every $n \in \mathbb{N}$, set
\[
\zeta_n = 1 - \sum_{i=1}^{m} \|U_i^{-1}L_i\sqrt{U_n}\|^2 \quad \text{(23)}
\]
and suppose that
\[
\inf_{n \in \mathbb{N}} \max_{\mu \in \mathcal{G}_i} (\nu_1, n, \ldots, \nu_m,n) > \frac{1}{2} \quad \text{(24)}
\]
Assume that $U(c)$. The first one is corrupted with a noise with a variance the same scene $v$ on a image recovery example. Two observed images $w_1$ (SNR = 5.87 dB) (b), blurred image $w_2$ (SNR = 16.63 dB) (c), and restored image $\tilde{x}$ (SNR = 21.61 dB) (d).

Set

$$s_n = x_n - U_n(\nabla \theta(x_n) + c_n - z)$$

For $i = 1, \ldots, m$

$$q_{i,n} = \text{prox}_{\tau_1}^{-1}(v_{i,n} + U_{i,n}(L_1y_n - \nabla \ell_1(v_{i,n} - d_{i,n} - r_i)) + b_{i,n})$$

$$v_{i,n+1} = v_{i,n} + \lambda_n(q_{i,n} - v_{i,n})$$

$$p_n = s_n - U_n\sum_{i=1}^m L_i q_{i,n}$$

$$x_{n+1} = x_n + \lambda_n(p_n - x_n).$$

(25)

Assume that $f = 0$. Then $(x_n)_{n \in \mathbb{N}}$ converges weakly to a solution to (12) for every $i \in \{1, \ldots, m\}$, $(v_{i,n})_{n \in \mathbb{N}}$ converges weakly to some $\mathcal{H} \in G_1$, and $(\mathcal{H}_1, \ldots, \mathcal{H}_m)$ is a solution to (13).

The algorithm proposed in [23, 8] is a special case of the previous one, in the absence of errors, when $m = 1$, $\mathcal{H}$ and $G_1$ are finite dimensional spaces, $\ell_1 = \ell_2$ is used with $\tau \equiv 0$, $U_1 \equiv \text{Id}$ with $\tau \equiv 0$, $U_2 \equiv \sigma \text{Id}$ with $\tau \equiv 0$, and no relaxation ($\lambda_n \equiv 1$) or a constant one ($\lambda_n \equiv \kappa < 1$) is performed.

5. APPLICATION TO IMAGE RESTORATION

We illustrate the flexibility of the proposed primal-dual algorithms on an image recovery example. Two observed images $w_1$ and $w_2$ of the same scene $x \in \mathbb{R}^N$ ($N = 256^2$) are available (see Fig. 1(a)-(c)). The first one is corrupted with a noise with a variance $\theta_1^2 = 576$,

while the second one has been degraded by a linear operator $H \in \mathbb{R}^{N \times N}$ ($7 \times 7$ uniform blur) and a noise with variance $\theta_2^2 = 25$. The noise components are mutually statistically independent, additive, zero-mean, white, and Gaussian distributed. Note that this kind of multivariate restoration problem is encountered in some push-broom satellite imaging systems.

An estimate $\tilde{x}$ of $x$ is computed as a solution to (12) where $m = 2$, $z = 0$, $r_1 = 0$, $r_2 = 0$,

$$h = \frac{1}{\theta_1^2} \| \cdot - w_1 \|^2 + \frac{1}{\theta_2^2} \| H \cdot w_2 \|^2,$$

$$g_1 = \ell_{[0,255]^N}, \quad g_2 = \kappa \| \cdot \|_{1,2},$$

$$f = 0, \quad \ell_1 = \ell_2 = \ell_{(0)}$$

(26)

(27)

(28)

where the second function in (27) denotes the $\ell_{1,2}$-norm and $\kappa \in [0, +\infty]$. In addition, $L_1 = \text{Id}$ and $L_2 = [G_1^T, G_2^T]^T$ where $G_1 \in \mathbb{R}^{N \times N}$ and $G_2 \in \mathbb{R}^{N \times N}$ are horizontal and vertical discrete gradient operators. Function $g_1$ introduces some a priori constraint on the range values in the target image, while function $g_2 \circ L_2$ corresponds to a classical total variation regularization. The minimization problem is solved numerically by using Algorithm (25) with $\lambda_n \equiv 1$. In a first experiment, standard choices of the algorithm parameters are made by setting $U_n = \tau \text{Id}$, $U_1 = \sigma \text{Id}$, and $U_2 = \sigma_2 \text{Id}$ with $(\tau, \sigma, \sigma_2) \in [0, +\infty]^3$. In a second experiment, a more sophisticated choice of the metric is made. The operators $(U_n)_{n \in \mathbb{N}}$, $(U_1,n)_{n \in \mathbb{N}}$ and $(U_2,n)_{n \in \mathbb{N}}$ are still chosen diagonal and constant in order to facilitate the implementation of the algorithm, but the diagonal values are optimized in an empirical manner. A similar strategy was applied in [26] in the case of Algorithm (21). The regularization parameter $\kappa$ has been set so as to get the highest value of the resulting signal-to-noise ratio (SNR).

The restored image is displayed in Fig. 1(d). Fig. 2 shows the convergence profile of the algorithm. We plot the evolution of the normalized Euclidean distance (in log scale) between the iterates and $\tilde{x}$ in terms of computational time (Matlab R2011b codes running on a single-core Intel i7-2620M CPU@2.7 GHz with 8 GB of RAM). An approximation of $\tilde{x}$ obtained after 5000 iterations is used. This result illustrates the fact that an appropriate choice of the metric may be beneficial in terms of speed of convergence.
6. REFERENCES


