Robust self-triggered coordination by ternary controllers

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The simplest and best known example of coordination:

• Consider *n* systems (integrators)

$$\dot{x}_i = u_i \qquad i \in I := \{1, \ldots, n\}$$

linked by an undirected connected graph G = (I, E). N_i is the set of neighbors of system *i*

- **Control problem:** Design inputs u_i , $i \in I$,
 - which depend on x_i and $\{x_j : j \in \mathcal{N}_i\}$ (local information),
 - such that

$$x_i - x_j \rightarrow 0 \quad \forall i, j$$

Why (still) studying consensus?

• It is a prototypical problem:

solutions can help us to understand more complex problems

- It is useful in many application fields:
 - robotic networks
 - sensors networks
 - distribution networks
 - opinion dynamics
 - load balancing
- It is well studied:

Proposition (Standard consensus)

If the graph G is connected, the control law
$$u_i = \sum_{j \in \mathcal{N}_i} (x_j - x_i)$$

guarantees that $\lim_{t \to \infty} x_i(t) = c$ for all i, where $c = \sum_{j=1}^n \frac{x_j(0)}{n}$.

Standard consensus requires continuous flow of information from neighbors

this is too demanding!

We instead want a scenario in which

- sensors collect information only upon need \longrightarrow discrete event times!
- the continuous-time systems "naturally" interacts with the discrete-time information acquisition
- the whole system is robust against network uncertainties (delays, poor synchronization of local clocks, limited data rate communication)

Hybrid system definition and main result

State variables $(i \in I)$

- consensus variables: $x_i \in \mathbb{R}$
- control variables: $u_i \in \{-1, 0, +1\}$ (ternary controls)
- local clock variables: $\theta_i \in \mathbb{R}$

Continuous evolution when no information exchange occurs

$$\begin{cases} \dot{x}_i = u_i \\ \dot{u}_i = 0 \\ \dot{\theta}_i = -1 \end{cases}$$

Jumps occur at every *t* such that the set

$$\mathcal{I}(\theta, t) = \{i \in I : \theta_i = 0\} \neq \emptyset$$

A hybrid coordination system II

Discrete evolution: how the exchange of information affects the systems

$$\begin{cases} x_i(t^+) = x_i(t) \quad \forall i \in I \\ u_i(t^+) = \begin{cases} \operatorname{sign}_{\varepsilon}(\operatorname{ave}_i(t)) & \text{if } i \in \mathcal{I}(\theta, t) \\ u_i(t) & \text{otherwise} \\ \theta_i(t^+) = \begin{cases} f_i^{\alpha}(x(t)) & \text{if } i \in \mathcal{I}(\theta, t) \\ \theta_i(t) & \text{otherwise} \end{cases} \end{cases}$$

•
$$\operatorname{ave}_i(t) := \sum_{j \in \mathcal{N}_i} (x_j(t) - x_i(t))$$
 is the "consensus feedback"
• $\operatorname{sign}_{\varepsilon}(z) = \begin{cases} \operatorname{sign}(z) & \text{if } |z| \ge \varepsilon \\ 0 & \text{otherwise} \end{cases}$
• $\varepsilon > 0$ is a *sensitivity* parameter
• $\alpha \in (0, 1)$ is a *robustness* parameter

A hybrid coordination system III

Next sampling time is chosen by $\theta_i(t^+) = \begin{cases} f_i^{\alpha}(x(t)) & \text{if } i \in \mathcal{I}(\theta, t) \\ \theta_i(t) & \text{otherwise} \end{cases}$

$$f_i^{\alpha}(x) = \begin{cases} \frac{\alpha}{2 \deg_i} \left| \sum_{j \in \mathcal{N}_i} (x_j - x_i) \right| & \text{if } \left| \sum_{j \in \mathcal{N}_i} (x_j - x_i) \right| \ge \varepsilon \\ \frac{\alpha}{2 \deg_i} \varepsilon & \text{otherwise} \end{cases}$$

so that

- sign(ave_i) is constant during inter-sampling interval (from t_k^i to t_{k+1}^i)
- "dwell time" property holds: $t_{k+1}^i t_k^i \ge \frac{\alpha \varepsilon}{2 \deg_{\max}}$
- $\varepsilon > 0$ is a *sensitivity* parameter
- $\alpha \in (0,1)$ is a *robustness* parameter

Protocol A

- 1: initialization: for all $i \in I$, set $u_i(0) \in \{-1, 0, +1\}$ and $\theta_i(0) = 0$;
- 2: for all $i \in I$ do
- 3: while $\theta_i(t) > 0$ do
- 4: *i* applies the control $u_i(t)$;
- 5: end while
- 6: **if** $\theta_i(t) = 0$ **then**
- 7: for all $j \in \mathcal{N}_i$ do
- 8: *i* polls *j* and collects the information $x_j(t) x_i(t)$;
- 9: end for
- 10: *i* computes $ave_i(t)$;
- 11: $i \text{ computes } \theta_i(t^+) = f_i^{\alpha}(x(t));$
- 12: $i \text{ computes } u_i(t^+) = \operatorname{sign}_{\varepsilon}(\operatorname{ave}_i(t));$
- 13: end if
- 14: end for

Theorem (Practical consensus)

For every initial condition \bar{x} , let x(t) be the solution to Protocol A such that $x(0) = \bar{x}$. Then x(t) converges in finite time to a point x^* belonging to the set

$$\mathcal{E} = \{x \in \mathbb{R}^n : |\sum_{j \in \mathcal{N}_i} (x_j - x_i)| < \varepsilon \ \forall \ i \in I\}$$

Proof sketch:

- Lyapunov-like argument $V(x) = \frac{1}{2}x^T L x = \frac{1}{2}\sum_{\{i,j\}\in E} (x_i x_j)^2$
- we approximate the dynamics x_i = sign(ave_i), which is known to imply finite-time convergence

J. Cortés. Finite-time convergent gradient flows with applications to network consensus. *Automatica*, 42(11):1993–2000, 2006



Sample evolutions of states x and corresponding controls u in Protocol A on a ring with n = 5 nodes, $\varepsilon = 0.02$

Robustness

Robustness: clock skews and quantized measurements

Continuous dynamics

$$\begin{cases} \dot{x}_i = u_i \\ \dot{u}_i = 0 \\ \dot{\theta}_i = -\mathbf{R_i} \end{cases}$$

where $R_i > 0$ are local *(skewed)* clock rates

Quantized measurements: each system measures $q(x_i - x_j)$



Discrete dynamics

$$\begin{cases} x_i(t^+) = x_i(t) \quad \forall i \in I \\ u_i(t^+) = \begin{cases} \operatorname{sign}_{\varepsilon}(\operatorname{qave}_i(t)) & \text{if } i \in \mathcal{I}(\theta, t) \\ u_i(t) & \text{otherwise} \end{cases} \\ \theta_i(t^+) = \begin{cases} f_i^{\alpha}(x(t)) & \text{if } i \in \mathcal{I}(\theta, t) \\ \theta_i(t) & \text{otherwise} \end{cases} \end{cases}$$

where

•
$$qave_i(t) := \sum_{j \in \mathcal{N}_i} q(x_j(t) - x_i(t))$$

• $f_i^{\alpha}(x) = \begin{cases} \frac{\alpha}{2 \deg_i} |\sum_{j \in \mathcal{N}_i} qave_i| & \text{if } |\sum_{j \in \mathcal{N}_i} qave_i| \ge \varepsilon \\ \frac{\alpha}{2 \deg_i} \varepsilon & \text{otherwise} \end{cases}$

Theorem (Clock skew & quantization robustness)

Assume that $R_i \ge R_{\min} > 0$ for all $i \in I$. If $\varepsilon > \frac{1}{2}d_{\max}\Delta$ and

$$\alpha < \frac{2\varepsilon - d_{\max}\Delta}{2\varepsilon} R_{\min},$$

then x(t) converges in finite time to a point in

$$\mathcal{E}_2 = \{x \in \mathbb{R}^n \, : \, |\sum_{j \in \mathcal{N}_i} (x_j - x_i)| < 2\varepsilon\}$$

• Size of the region of convergence depends on quantizer resolution Δ • $\alpha \in (0, R_{\min})$ quantifies the stability margin

Conclusion

Work done & its positive features:

- Coordination with self-triggered information collection (upon need)
- Coordination using coarse controllers and relative measurements
- Convergence of solutions with guaranteed dwell-time
- Finite-time convergence (with an explicit estimate)
- Other good properties in distributed systems:
 - No need for knowledge of absolute time
 - Robust against delays, quantization, clock skews, parameter uncertainties

Available extensions/variations (not presented here)

- Independent polling of neighbors & self-triggered gossiping
- Self-triggered protocols for asymptotical consensus $(x_i x_j \rightarrow 0)$

Future work

- Application to network flow control
- Extension to saturated controllers and other constrained controllers
- Extension to higher-dimensional systems (cf. joint work with J.M. Hendrickx)
- Extension to more complex coordination tasks (e.g., formation control of autonomous systems)

Details and related literature available in

C. De Persis and P. Frasca. Robust self-triggered coordination with ternary controllers. *IEEE Transactions on Automatic Control*, provisionally accepted. revised Dec. 2012

Thank you for your attention