

# Blind Techniques

Pierre Comon

I3S, CNRS, Univ. of Nice, Sophia-Antipolis, France

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# Part I

## Principle & Tools

# Introduction

- Modeling
- General concepts, a  $2 \times 2$  example
- Historical survey, Origins

# Observation model

$$\mathbf{x} = \mathbf{H} \mathbf{s} + \mathbf{v} \quad (1)$$

- $\mathbf{x}$ : observed, dim  $K$
- $P$ : source vector, dim  $P$
- $\mathbf{H}$ :  $K \times P$  mixing matrix
- $\mathbf{v}$ : additive noise

# Taxonomy (1)

**Static/Dynamic** and **Noisy/Noiseless**:

$$\mathbf{x}[n] = \mathbf{H} \star \mathbf{s}[n] + \mathbf{v}[n] \quad (2)$$

**Over/Under-Determined**:

*Number of sources :  $P \leq^{\text{Underdet}} K$  : Number of sensors*

# Taxonomy (2)

Transmit/Receive diversity:

| Sources | Sensors     |             |
|---------|-------------|-------------|
|         | 1           | $K > 1$     |
| 1       | <b>SISO</b> | <b>SIMO</b> |
| $P > 1$ | <b>MISO</b> | <b>MIMO</b> |

# Taxonomy (3)

**One additional assumption required on sources:**

- mutually independent sources
- discrete sources
- colored sources
- nonstationary sources



# Principal Component Analysis (PCA)

## Goal

Given a  $K$ -dimensional r.v.,  $\mathbf{x}$ , find  $\mathbf{U}$  and  $\mathbf{z}$  such that

- Observation

$$\mathbf{x} = \mathbf{U} \mathbf{z}$$

- $\mathbf{z}$  has uncorrelated components  $z_i$

**NB:** Because of lack of uniqueness,  $\mathbf{U}$  is often assumed to be unitary.

# Independent Component Analysis (ICA)

## Goal

Given a  $K$ -dimensional r.v.,  $\mathbf{x}$ , find  $\mathbf{H}$  and  $\mathbf{s}$  such that

- Observation

$$\mathbf{x} = \mathbf{H}\mathbf{s} \quad (3)$$

- $\mathbf{s}$  has mutually statistically independent components  $s_i$

➡ “*Blind*” Source Separation: only outputs  $x_i$  are observed.

# Uniqueness

## Inherent indeterminations

if  $\mathbf{s}$  has independent components  $s_i$ , so has  $\mathbf{\Lambda P s}$   
 where  $\mathbf{\Lambda}$  is invertible diagonal and  $\mathbf{P}$  permutation

## Solutions

If  $(\mathbf{A}, \mathbf{s})$  solution, then  $(\mathbf{A}\mathbf{\Lambda P}, \mathbf{P}^T \mathbf{\Lambda}^{-1} \mathbf{s})$  also is.

- “*Essential uniqueness*”: unique up to a *trivial filter*, i.e. a scale-permutation (cf. slide 67)
- Whole equivalence class of solutions  $\Rightarrow$  Look for one representative.

# Decorrelation vs Independence

## Example 1: Mixture of 2 identically distributed sources

Consider the mixture of two independent sources

$$\begin{pmatrix} x_1 \\ x_2 \end{pmatrix} = \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \cdot \begin{pmatrix} s_1 \\ s_2 \end{pmatrix}$$

where  $E\{s_i^2\} = 1$  and  $E\{s_i\} = 0$ . Then  $x_i$  are *uncorrelated*:

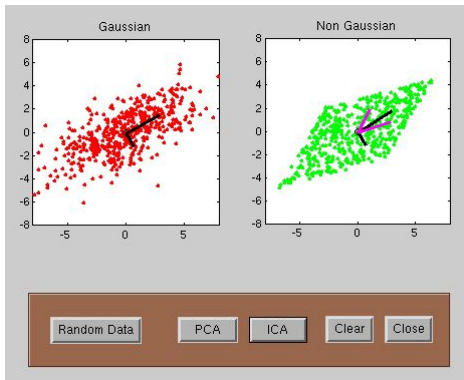
$$E\{x_1 x_2\} = E\{s_1^2\} - E\{s_2^2\} = 0$$

But  $x_i$  are *not independent* since, for instance:

$$E\{x_1^2 x_2^2\} - E\{x_1^2\}E\{x_2^2\} = E\{s_1^4\} + E\{s_2^4\} - 6 \neq 0$$

# PCA vs ICA

## Example 2: 2 sources and 2 sensors



# Historical survey: Static MIMO

- **The ancestors:** Dugué'51, Darmois'53, Feller'66, Friedman'74, Donoho'80
- **The first shy steps in ICA:** Bar-Ness'82, Jutten'83, Fety'88
- **The first steps in Multi-Way:** Carroll-Chang'70, Harshman'70, Kruskal'77
- **First closed-form solutions:** Comon'89, Cardoso'92
- **First IT frameworks:** Comon'91, Cardoso'93, Comon'94, Bell'95, Delfosse-Loubaton'95
- **Specific applications:** Hyvarinen'97, Pajunen'97, Amari'98, Grellier'98, Parra'2000
- **Discrete/CM:** Talwar'96, VanderVeen'97, Grellier'00

# Historical survey: Static MIMO (cont'd)

- **Other:** Cao-Liu'96, VanDerVeen-Paulraj'96, Moreau-Pesquet'97, Taleb-Jutten'97, Comon'96, Ferreol-Chevalier'98, Belouchrani'98, Lee-Lewicki'99, deLathauwer'00, Pham-Cardoso'2000, Yeredor'2000, Sidiropoulos-Bro'00, Albera'04, Comon-Rajih'05, deLathauwer'05...

# Historical survey: Convolutive SISO

## ■ Identification

- **Kurtosis** Benveniste-Ruget'80, Tugnait'89
- **Non circularity/Alphabet:** Yellin-Porat'93, Grellier-Comon'99, Ciblat-Loubaton'02, Lebrun-Comon'03

## ■ Equalization

- **CMA:** Sato'75, Godard'80, Treichler'85
- **Kurtosis:** Benveniste-Goursat'84, Donoho'81, Shalvi-Weinstein'90
- **Bispectrum:** Marron'90, Matsuoka'84, LeRoux'93

**NB:** Earlier equalization algorithms, e.g. Decision-Directed, need the eye to be open.



# Historical survey: Convolutional SIMO

- **Subspace:** Slock'94, Xu-Tong'95, Moulines-Duhamel'95, Xu-Liu-Tong'95, Gurelli-Nikias'95, Gesbert-Duhamel'97
- **Linear Prediction:** AbedMeraim-Moulines-Loubaton'97, Gesbert-Duhamel'00

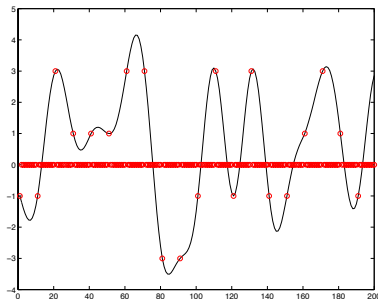
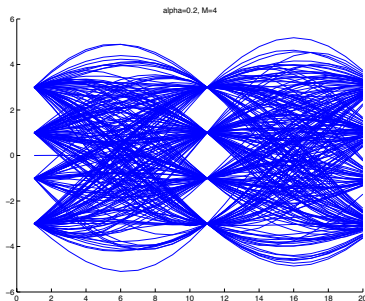
# Historical survey: Convolutional MIMO

- **Subspace:** Gorokhov-Loubaton'97, Chevreuil-Loubaton'97, Loubaton-Moulines'01
- **Linear Prediction:** Comon'90, Ding'96, AbedMeraim-Loubaton'97, Gorokhov-Loubaton'99
- **Kurtosis:** Comon'96, Tugnait'97, Simon-Loubaton'98, Touzni'98
- **Discrete/CM:** Touzni-Fijalkow'98, VanDerVeen-Talwar'95, Ayadi-Slock'98

# Origins of “Blind” Techniques

**Pulse Amplitude Modulation (PAM)** in baseband:

$$x(t) = a \sum_k h(t - k T) u_k$$



PAM4: symbols  $u_k \in \{-3, -1, 1, 3\}$

# General bibliography

## ■ Books on HOS, ICA, or Multi-Way:

Lacoume-Amblard-Comon'97 [LAC97] (freely downloadable, but in French)

Hyvarinen-Karhunen-Oja'01 (but dedicated only to FastICA)

Smilde-Bro-Geladi'04 [SBG04] (but dedicated only to Factor Analysis)

Cichocki-Amari'02 [CA02] (but Neural Networks oriented)

Comon-Jutten'06 [CJ07] [JC07] (but in French)

Comon-Jutten'08 (will cover more topics, but you have to wait!)

## ■ Other related books:

Kagan-Linnik-Rao'73 [KLR73]

McCullagh'87 [McC87]

Nikias-Petropulu'93 [NP93]

Haykin'2000 [HAY00a] [HAY00b]

# Algebraic tools

- Singular Value decomposition (SVD)
- Spatial whitening (Standardization)
- PCA by pair sweeping
- Filter decomposition
- Time Whitening
- Space-time Whitening
- Matched filter

# Singular Value Decomposition (SVD)

Every matrix  $\mathbf{M}$  may be decomposed into:

$$\mathbf{M} = \mathbf{U} \mathbf{\Sigma} \mathbf{V}^H$$

where

- $\mathbf{U}$  and  $\mathbf{V}$  are unitary
- $\mathbf{\Sigma}$  is positive real diagonal
- $\mathbf{u}_i$  and  $\mathbf{v}_i$  of  $\mathbf{U}$  and  $\mathbf{V}$  are the left and right singular vectors:

$$\mathbf{M} \mathbf{v}_i = \mathbf{u}_i \sigma_i \quad \mathbf{M}^H \mathbf{u}_i = \mathbf{v}_i \sigma_i$$

- $\mathbf{u}_i$  are eigenvectors of  $\mathbf{M} \mathbf{M}^H$ , and  $\mathbf{v}_i$  those of  $\mathbf{M}^H \mathbf{M}$ , associated with  $\sigma_i^2$ .

# Spatial whitening (1)

**Standardization via Cholesky or QR** Let  $\mathbf{x}$  be a zero-mean r.v. with covariance matrix:

$$\mathbf{\Gamma}_x \stackrel{\text{def}}{=} \mathbb{E}\{\mathbf{x} \mathbf{x}^H\}$$

Then Cholesky yields:

$$\exists \mathbf{L} / \quad \mathbf{L} \mathbf{L}^H = \mathbf{\Gamma}_x$$

Consequence:  $\mathbf{L}^{-1}\mathbf{x}$  has a unit variance.

Variable  $\tilde{\mathbf{x}} \stackrel{\text{def}}{=} \mathbf{L}^{-1}\mathbf{x}$  is a *standardized random variable*.

- QR factorization of data matrix as  $\mathbf{X} = \mathbf{L} \tilde{\mathbf{X}}$  yields same  $\mathbf{L}$  as Cholesky factorization of sample covariance, but more accurate.
- Limitation:  $\mathbf{L}$  may not be invertible if the covariance  $\mathbf{\Gamma}_x$  is not full rank.

# Spatial whitening (2)

## Standardization via PCA

### Definition

PCA is based on second order statistics

- Observed random variable  $\mathbf{x}$  of dimension  $K$ . Then  $\exists(\mathbf{U}, \mathbf{z})$ :

$$\mathbf{x} = \mathbf{U}\mathbf{z}, \mathbf{U} \text{ unitary}$$

where *Principal Components*  $z_i$  are uncorrelated  
 $i$ th column  $\mathbf{u}_i$  of  $\mathbf{U}$  is called  *$i$ th PC Loading vector*

- Two possible calculations:
  - EVD of Covariance  $\mathbf{R}_x$ :  $\mathbf{R}_x = \mathbf{U}\mathbf{\Sigma}^2\mathbf{U}^H$
  - Sample estimate by SVD:  $\mathbf{X} = \mathbf{U}\mathbf{\Sigma}\mathbf{V}^H$



# Spatial whitening (3)

## Summary

Find a linear transform  $\mathbf{L}$  such that vector  $\tilde{\mathbf{x}} \stackrel{\text{def}}{=} \mathbf{L}\mathbf{x}$  has unit covariance. Many possibilities, including:

- PCA yields  $\tilde{\mathbf{x}} = \mathbf{\Sigma}^{-1} \mathbf{U}^H \mathbf{x}$
- Cholesky  $\mathbf{R}_x = \mathbf{L} \mathbf{L}^H$  yields  $\tilde{\mathbf{x}} = \mathbf{L}^{-1} \mathbf{x}$

## Remarks

- Infinitely many possibilities:  $\mathbf{L}$  is as good as  $\mathbf{L} \mathbf{Q}$ , for any unitary  $\mathbf{Q}$ .
- If  $\mathbf{R}_x$  not invertible, then  $\mathbf{L}$  not invertible (ill-posed). One may use pseudo-inverse of  $\mathbf{\Sigma}$  in PCA to compute  $\mathbf{L}$ , or regularize  $\mathbf{R}_x$ .

## Plane rotations

Application of a Givens rotation on both sides of a matrix allows to set a pair of zeros in a symmetric matrix:

$$\begin{pmatrix} c & . & s & . \\ . & 1 & . & . \\ -s & . & c & . \\ . & . & . & 1 \end{pmatrix} \mathbf{A} \begin{pmatrix} c & . & -s & . \\ . & 1 & . & . \\ s & . & c & . \\ . & . & . & 1 \end{pmatrix} = \begin{pmatrix} * & x & 0 & x \\ x & . & x & . \\ 0 & x & * & x \\ x & . & x & . \end{pmatrix}$$

Same result obtained:

- either by setting 0
- or by maximizing \*'s

# Jacobi sweeping for PCA

Cyclic by rows/columns algorithm for a  $4 \times 4$  real symmetric matrix

$$\begin{pmatrix} . & . & . & . \\ . & . & . & . \\ . & . & . & . \\ . & . & . & . \end{pmatrix} \rightarrow \begin{pmatrix} * & 0 & x & x \\ 0 & * & x & x \\ x & x & . & . \\ x & x & . & . \end{pmatrix} \rightarrow \begin{pmatrix} * & x & 0 & x \\ x & . & x & . \\ 0 & x & * & x \\ x & . & x & . \end{pmatrix} \rightarrow \begin{pmatrix} * & x & x & 0 \\ x & . & . & x \\ x & . & . & x \\ 0 & x & x & * \end{pmatrix} \\
 \begin{pmatrix} . & x & x & 0 \\ x & * & 0 & x \\ x & 0 & * & x \\ 0 & x & x & . \end{pmatrix} \rightarrow \begin{pmatrix} . & x & . & x \\ x & * & x & 0 \\ . & x & . & x \\ x & 0 & x & * \end{pmatrix} \rightarrow \begin{pmatrix} . & . & x & x \\ . & . & x & x \\ x & x & * & 0 \\ x & x & 0 & * \end{pmatrix}$$

\*: maximized, x: minimized, 0: canceled, .: unchanged

# Scalar Filter Decomposition

- Any rational scalar filter  $g[z]$  can be decomposed into:

$$\gamma[z] = u[z] \ell[z], \quad u[1/z^*] u[z] = 1, \quad \forall z \quad (4)$$

- $\ell[z]$  is *minimum phase*: all its roots inside the unit circle
- $u[z]$  is *all-pass*, and hence *lossless*: flat frequency response (only phase varies).

# Multivariate Filter Decomposition

- Any rational filter with Impulse Response matrix  $\mathbf{F}[k]$  and complex gain  $\check{\mathbf{F}}[z]$ , can be decomposed into:

$$\check{\mathbf{F}}[z] = \check{\mathbf{U}}[z] \check{\mathbf{L}}[z], \quad \check{\mathbf{U}}[1/z^*]^H \check{\mathbf{U}}[z] = \mathbf{I}, \quad \forall z \quad (5)$$

- $\check{\mathbf{L}}[k]$  is *triangular minimum phase* filter: roots of  $\det(\check{\mathbf{L}}[z])$  inside unit circle
- $\check{\mathbf{U}}[k]$  *para-unitary* filter
- In static MIMO case, one gets QR:

$$\mathbf{F} = \mathbf{U} \mathbf{L}, \quad \mathbf{U}^H \mathbf{U} = \mathbf{I} \quad (6)$$

where  $\mathbf{L}$  is triangular and  $\mathbf{U}$  unitary.

- Decomposition not unique.

# Time Whitening

Let  $x[k]$  be a scalar second order stationary process,  $\check{x}[z]$  its  $z$ -transform, and its power spectrum given by:

$$\gamma_x[z] \stackrel{\text{def}}{=} E\{\check{x}[z] \check{x}[1/z^*]^*\}$$

From (4), the power spectrum can be decomposed as:

$$\exists \ell[z] / \ell[z] \ell[1/z^*]^* = \gamma_x[z]$$

where filter  $\ell[z]$  is not unique, and defined up to an all-pass filter.  $1/\ell[z]$  is a *whitening filter*, if it exists.

# Space-time Whitening

- Let  $\mathbf{x}[k]$  be a multivariate second order stationary random process,  $\mathbf{x}[z]$  its  $z$ -transform, and power spectral matrix:

$$\mathbf{\Gamma}_x[z] \stackrel{\text{def}}{=} \mathbb{E}\{\mathbf{x}[z] \mathbf{x}[1/z^*]^H\}$$

Then, from (5)

$$\exists \check{\mathbf{L}}[z] / \check{\mathbf{L}}[z] \check{\mathbf{L}}[1/z^*]^H = \mathbf{\Gamma}_x[z]$$

- If  $\check{\mathbf{L}}[z]$  admits an inverse, then we may take  $\check{\mathbf{G}}[z] = \check{\mathbf{L}}[z]^{-1}$  as *whitening filter*, i.e.  $\tilde{\mathbf{x}}[k] = \mathbf{G}[k] \star \mathbf{x}[k]$ .

# Spatial Matched Filter

If  $\mathbf{x} = \mathbf{H} \mathbf{s} + \mathbf{v}$ , where  $\mathbf{H}$  is known, one can estimate  $\mathbf{s}$  by spatial filtering as

$$\hat{\mathbf{s}} = \mathbf{W} \mathbf{x}$$

- Spatial Matched Filter:  $\mathbf{W} = \mathbf{H}^H \mathbf{R}_x^{-1}$
- Least Squares:  $\mathbf{W} = (\mathbf{H}^H \mathbf{H})^{-1} \mathbf{H}^H$
- Weighted Least Squares:  $\mathbf{W} = (\mathbf{H}^H \mathbf{B}^{-1} \mathbf{H})^{-1} \mathbf{H}^H \mathbf{B}^{-1}$   
where  $\mathbf{B}$  is the noise spatial coherence.



# Statistical Tools

- Statistical Independence
- Mutual vs Pairwise Independence
- Cumulants of a scalar r.v.
- Multivariate Cumulants
- Complex variables, circularity
- Central limit, Edgeworth expansion
- Mutual Information, approximation

# Statistical Independence

## Definition

Components  $s_k$  of a  $K$ -dimensional r.v.  $\mathbf{s}$  are *mutually independent*



The *joint* pdf equals the *product of marginal* pdf's:

$$p_{\mathbf{s}}(\mathbf{u}) = \prod_k p_{s_k}(u_k) \quad (7)$$

## Definition

Components  $s_k$  of  $\mathbf{s}$  are *pairwise independent*  $\Leftrightarrow$  Any pair of components  $(s_k, s_\ell)$  are mutually independent.

# Mutual vs Pairwise independence (1)

## Example 3: Pairwise but not Mutual independence

- Bag containing 4 Bowls denoted  $\{RB, YB, GB, RYG\}$ :  
1 Red, 1 Yellow, 1 Green, 1 with the 3 colors.
- Equal drawing probabilities:  
 $P(RB) = P(YB) = P(GB) = P(RYG) = 1/4$
- Event “R”  $\stackrel{\text{def}}{=}$  draw a bowl containing Red  $\Rightarrow$   
 $P(R) = P(RB) + P(RYG) = 1/2$
- Then  $P(R \cap Y) = P(RYG) = 1/4$   
equal to  $P(R) * P(Y) \Rightarrow$  *Pairwise independent* Events
- But  $P(R \cap Y \cap G) = P(RYG) = 1/4$   
not equal to  $P(R) * P(Y) * P(G) = 1/8 \Rightarrow$   
Events are *not Mutually independent*

## Mutual vs Pairwise independence (2)

### Example 4: Pairwise but not Mutual independence

- 3 mutually independent BPSK sources,  $x_i \in \{-1, 1\}$ ,  $1 \leq i \leq 3$
- Define  $x_4 = x_1 x_2 x_3$ . Then  $x_4$  is also BPSK, *dependent on  $x_i$*
- $x_k$  are *pairwise independent*:  

$$p(x_1 = a, x_4 = b) = p(x_4 = b | x_1 = a) \cdot p(x_1 = a) =$$

$$p(x_2 x_3 = b/a) \cdot p(x_1 = a)$$

But  $x_1$  and  $x_2 x_3$  are BPSK  $\Rightarrow$

$$p(x_2 x_3 = b/a) \cdot p(x_1 = a) = \frac{1}{2} \cdot \frac{1}{2}$$
- But  $x_k$  obviously not mutually independent,  $1 \leq k \leq 4$   
 In particular,  $\text{Cum}\{x_1, x_2, x_3, x_4\} = 1 \neq 0$

# Mutual vs Pairwise independence (3)

## Darmois's Theorem (1953)

Let two random variables be defined as linear combinations of independent random variables  $x_i$ :

$$X_1 = \sum_{i=1}^N a_i x_i, \quad X_2 = \sum_{i=1}^N b_i x_i$$

Then, if  $X_1$  and  $X_2$  are independent, those  $x_j$  for which  $a_j b_j \neq 0$  are Gaussian.

# Mutual vs Pairwise independence (4)

## Corollary

If  $\mathbf{z} = \mathbf{C}\mathbf{s}$ , where  $s_i$  are independent r.v., with at most one of them being Gaussian, then the following properties are equivalent:

- 1 Components  $z_i$  are pairwise independent
- 2 Components  $z_i$  are mutually independent
- 3  $\mathbf{C} = \mathbf{\Lambda}\mathbf{P}$ , with  $\mathbf{\Lambda}$  diagonal and  $\mathbf{P}$  permutation

# Characteristic functions

## First c.f.

- Real Scalar:  $\Phi_x(t) \stackrel{\text{def}}{=} E\{e^{jt^T x}\} = \int_{\mathbf{u}} e^{jt^T \mathbf{u}} dF_x(\mathbf{u})$
- Real Multivariate:  $\Phi_x(\mathbf{t}) \stackrel{\text{def}}{=} E\{e^{j\mathbf{t}^T \mathbf{x}}\} = \int_{\mathbf{u}} e^{j\mathbf{t}^T \mathbf{u}} dF_x(\mathbf{u})$

## Second c.f.

- $\Psi(\mathbf{t}) \stackrel{\text{def}}{=} \log \Phi(\mathbf{t})$
- Properties:
  - Always exists in the neighborhood of 0
  - Uniquely defined as long as  $\Phi(\mathbf{t}) \neq 0$

# Definition of Cumulants

## ■ Moments:

$$\mu'_r \stackrel{\text{def}}{=} \mathbb{E}\{x^r\} = (-j)^r \left. \frac{\partial^r \Phi(t)}{\partial t^r} \right|_{t=0} \quad (8)$$

## ■ Cumulants:

$$\mathcal{C}_{x(r)} \stackrel{\text{def}}{=} \text{Cum}\underbrace{\{x, \dots, x\}}_{r \text{ times}} = (-j)^r \left. \frac{\partial^r \Psi(t)}{\partial t^r} \right|_{t=0} \quad (9)$$

## ■ Needs the existence of the expansion. Counter example: Cauchy

$$p_x(u) = \frac{1}{\pi(1+u^2)}$$

## ■ Relationship between Moments and Cumulants obtained by expanding both sides in Taylor series:

$$\text{Log } \Phi_x(t) = \Psi_x(t)$$



# First Cumulants

- $\mathcal{C}_{(2)}$  is the variance:
- For zero-mean r.v.:  $\mathcal{C}_{(3)} = \mu_{(3)}$ , and  $\mathcal{C}_{(4)} = \mu_{(4)} - 3\mu_{(2)}^2$
- Warning: it is not true that  $\mathcal{C}_{(r)}$  is the moment of a variable  $x - x_g$ ,  $x_g$  Gaussian
- Standardized cumulants:

$$\mathcal{K}_{(r)} = \text{Cum}_{(r)} \left\{ \frac{x - \mu'_{(1)}}{\sqrt{\mu_{(2)}}} \right\}$$

e.g. *Skewness*  $\mathcal{K}_3$ , and *Kurtosis*  $\mathcal{K}_4$ .

# Examples of cumulants (1)

## Example 5: Zero-mean Gaussian

- Moments

$$\mu_{(2r)} = \mu_{(2)}^r \frac{(2r)!}{r! 2^r}$$

In particular:

$$\mu_{(4)} = 3\mu_{(2)}^2, \quad \mu_{(6)} = 15\mu_{(2)}^3$$

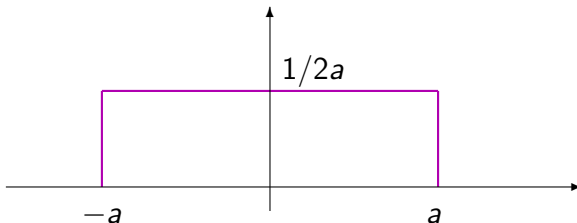
- $\mathcal{C}_{(4)} = 0, \quad \mathcal{K}_{(4)} = 0.$

- All Cumulants of order  $r > 2$  are null

# Examples of Cumulants (2)

## Example 6: Uniform

- uniformly distributed in  $[-a, +a]$  with probability  $\frac{1}{2a}$
- Moments:  $\mu_{(2k)} = \frac{a^{2k}}{2k+1}$
- 4th order Cumulant:  $\mathcal{C}_{(4)} = \frac{a^4}{5} - 3 \frac{a^4}{9} = -2 \frac{a^4}{15}$
- Kurtosis:  $\mathcal{K}_{(4)} = -\frac{6}{5}$ .



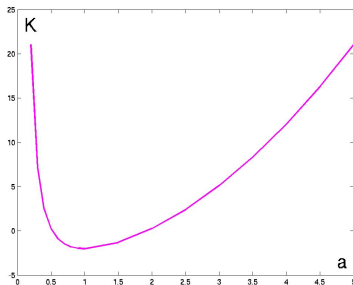
Proof...

# Examples of Cumulants (3)

## Example 7: Zero-mean standardized binary

- $x$  takes two values  $x_1 = -a$  and  $x_2 = 1/a$  with probabilities  $P_1 = \frac{1}{1+a^2}$ ,  $P_2 = \frac{a^2}{1+a^2}$
- Skewness is  $\mathcal{K}_{(3)} = \frac{1}{a} - a$
- Kurtosis is  $\mathcal{K}_{(4)} = \frac{1}{a^2} + a^2$
- Extreme values

Minimum Kurtosis  
for  $a = 1$  (symmetric):  
 $\mathcal{K}_{(4)} = -2$



# Sub- and Super-Gaussian r.v.

## Warning:

The concept of Sub/Super Gaussian is not uniquely defined in the literature. For instance, definitions below are *not equivalent*:

- Monotonicity of [BGR80]:  $f(u) = -\frac{1}{u} \frac{d \log p_x(u)}{du}$ .
- Tails of the standardized pdf are below/above those of Gaussian [ZIV95]
- Based on kurtosis [KS77]:
  - *Leptokurtic*: positive kurtosis
  - *mesokurtic*: zero kurtosis
  - *platykurtic*: negative kurtosis

# Definition of Multivariate cumulants

- Notation:  $C_{ij...l} \stackrel{\text{def}}{=} \text{Cum}\{X_i, X_j, \dots X_l\}$
- First cumulants:

$$\begin{aligned}\mu'_i &= C_i \\ \mu'_{ij} &= C_{ij} + C_i C_j \\ \mu'_{ijk} &= C_{ijk} + [3] C_i C_{jk} + C_i C_j C_k\end{aligned}$$

with  $[n]$ : Mccullagh's *bracket notation*.

- Next, for zero-mean variables:

$$\begin{aligned}\mu_{ijkl} &= C_{ijkl} + [3] C_{ij} C_{kl} \\ \mu_{ijklm} &= C_{ijklm} + [10] C_{ij} C_{klm}\end{aligned}$$

- General formula of Leonov Shirayev obtained by Taylor expansion of both sides of  $\Psi(\mathbf{t}) = \log \Phi(\mathbf{t}) \dots$

# Arrays and Tensors

**Definitions** Table  $\mathbf{T} = \{T_{ij..k}\}$

- *Order* of  $\mathbf{T} \stackrel{\text{def}}{=} \#$  of its ways =  $\#$  of its indices
- *Dimension*  $K_\ell \stackrel{\text{def}}{=} \text{range of the } \ell\text{th index}$
- $\mathbf{T}$  is *Cubic* when all dimensions  $K_\ell = K$  are equal
- $\mathbf{T}$  is *Symmetric* when it is square and when its entries do not change by *any* permutation of indices

**NB:** cf. course III for definitions and properties

# Definition of Complex Cumulants

## Definition

Let  $\mathbf{z} = \mathbf{x} + j\mathbf{y}$ . Then pdf  $p_{\mathbf{z}}$  = joint pdf  $p_{\mathbf{x},\mathbf{y}}$

## Notation

- Characteristic function:

$$\Phi_{\mathbf{z}}(\mathbf{w}) = \mathbb{E}\{\exp[j(\mathbf{x}^T \mathbf{u} + \mathbf{y}^T \mathbf{v})]\} = \mathbb{E}\{\exp[j\Re(\mathbf{z}^H \mathbf{w})]\}$$

where  $\mathbf{w} \stackrel{\text{def}}{=} \mathbf{u} + j\mathbf{v}$ .

- Generates Moments & Cumulants, e.g.:

Variance:  $\text{Var}\{\mathbf{z}\}_{ij} = C_{\mathbf{z}i}^j$

Higher orders:  $\text{Cum}\{z_i, \dots, z_j, z_k^*, \dots, z_\ell^*\} = C_{\mathbf{z}ij}^{kl}$

where *conjugated* r.v. are labeled *in superscript*.



# Circularity (1)

- $z$  is *circular in the strict sense* if its distribution does not depend on the phase of  $z$ . For a multivariate complex random variable  $\mathbf{z}$ , it means that:

$$\mathbf{z} \text{ and } \mathbf{z} e^{j\theta}, \forall \theta \in \mathbb{R}$$

have the same joint distribution.

- **Example 8: scalar circular complex Gaussian r.v.**

$$p_z(w) = \frac{1}{\pi \sigma^2} \exp -\frac{|w|^2}{\sigma^2}$$

defines a circular r.v.: only modulus appears.

## Circularity (2)

- There exist up to  $2^r$  distinct definitions of complex multivariate cumulants.
- At even order  $2r$ , cumulants having exacting  $r$  complex conjugations are termed *circular cumulants*.
- For instance, the cumulant below is circular

$$C_{\mathbf{z}ij}^{k\ell} = \text{Cum}\{z_i, z_j, z_k^*, z_\ell^*\}$$

whereas these ones are non circular

$$C_{\mathbf{z}ijl}^{\ell} = \text{Cum}\{z_i, z_j, z_k, z_\ell^*\}$$

$$C_{\mathbf{z}ijkl} = \text{Cum}\{z_i, z_j, z_k, z_\ell\}$$

- $\mathbf{z}$  is said to be *circular at order  $r$*  if its non circular cumulants of order  $r$  are all null:

$$p \neq r - p \Rightarrow \text{Cum}\{z_1, \dots, z_p, z_{p+1}^*, \dots, z_r^*\} = 0 \quad (10)$$

## Example of complex r.v.

**Example 9: PSK random variables** For a PSK-4 random variable,  $ZZ^* = 1$  and consequently:

$$\mathcal{C}_{(2)} = \mathbb{E}\{Z^2\} = 0, \mathcal{C}_{(2)}^{(2)} = -1, \mu_{(4)}^{(0)} = 1, \mathcal{C}_{(4)}^{(0)} = 1$$

It is thus circular up to order 3, but *non circular at order 4*.

# Properties of Cumulants

- **Multi-linearity** (also enjoyed by moments):

$$\text{Cum}\{\alpha X, Y, \dots, Z\} = \alpha \text{Cum}\{X, Y, \dots, Z\} \quad (11)$$

$$\text{Cum}\{X_1 + X_2, Y, \dots, Z\} = \text{Cum}\{X_1, Y, \dots, Z\} + \text{Cum}\{X_2, Y, \dots, Z\}$$

- **Cancellation:** If  $\{X_i\}$  can be partitioned into 2 groups of independent r.v., then

$$\text{Cum}\{X_1, X_2, \dots, X_r\} = 0 \quad (12)$$

- **Additivity:** If  $\mathbf{X}$  and  $\mathbf{Y}$  are independent, then

$$\begin{aligned} \text{Cum}\{X_1 + Y_1, X_2 + Y_2, \dots, X_r + Y_r\} &= \text{Cum}\{X_1, X_2, \dots, X_r\} \\ &+ \text{Cum}\{Y_1, Y_2, \dots, Y_r\} \end{aligned}$$

- **Inequalities**, e.g.:

$$\mathcal{K}_{(3)}^2 \leq \mathcal{K}_{(4)} + 2$$

Proof...

# Central Limit Theorem

Let  $N$  independent scalar r.v.,  $x(n)$ ,  $1 \leq n \leq N$  each with finite  $r$ th order Cumulant,  $\kappa_{(r)}(n)$ .

Define:

$$\bar{\kappa}_{(r)} = \frac{1}{N} \sum_{n=1}^N \kappa_{(r)}(n) \text{ and } y = \frac{1}{\sqrt{N}} \sum_{n=1}^N (x(n) - \bar{\kappa}_{(1)}).$$

As  $N \rightarrow \infty$ , the pdf  $f_y$  tends to a Gaussian.

## Proof.

Thanks to multi-linearity and additivity,  $C_{y(r)} = \frac{\bar{\kappa}_{(r)}}{N^{r/2-1}}$ ,  $\forall r \geq 2$ , tends to zero.

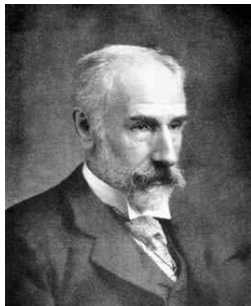
# Edgeworth expansion of a pdf (1)

The pdf  $p_{\mathbf{x}}(\mathbf{u})$  of a r.v.  $\mathbf{x}$  can be expanded about the Gaussian density  $g_{\mathbf{x}}(\mathbf{u})$  of same mean and variance, in terms of a combination of Hermite polynomials, ordered by decreasing significance in the sense of the Central Limit Theorem (CLT).

| Order      |            |                    |                      |                      |              |                    |                      |  |
|------------|------------|--------------------|----------------------|----------------------|--------------|--------------------|----------------------|--|
| $m^{-1/2}$ | $\kappa_3$ |                    |                      |                      |              |                    |                      |  |
| $m^{-1}$   | $\kappa_4$ | $\kappa_3^2$       |                      |                      |              |                    |                      |  |
| $m^{-3/2}$ | $\kappa_5$ | $\kappa_3\kappa_4$ | $\kappa_3^3$         |                      |              |                    |                      |  |
| $m^{-2}$   | $\kappa_6$ | $\kappa_3\kappa_5$ | $\kappa_3^2\kappa_4$ | $\kappa_4^2$         | $\kappa_3^4$ |                    |                      |  |
| $m^{-5/2}$ | $\kappa_7$ | $\kappa_3\kappa_6$ | $\kappa_3^2\kappa_5$ | $\kappa_4^2\kappa_3$ | $\kappa_3^5$ | $\kappa_4\kappa_5$ | $\kappa_3^3\kappa_4$ |  |

From slide 53,  $r$ th order Cumulants  $\sim O(m^{1-r/2})$ .

# Edgeworth expansion of a pdf (2)



Francis Edgeworth (1845-1926).

$$\begin{aligned} \frac{p_x(u)}{g_x(u)} = & 1 + \frac{1}{3!} \kappa_3 h_3(v) + \frac{1}{4!} \kappa_4 h_4(v) + \frac{10}{6!} \kappa_3^2 h_6(v) \\ & + \frac{1}{5!} \kappa_5 h_5(v) + \frac{35}{7!} \kappa_3 \kappa_4 h_7(v) + \frac{280}{9!} \kappa_3^3 h_9(v) + \dots \end{aligned}$$

## Mutual Information: definition

- According to the definition of page 34, one should measure a divergence:

$$\delta \left( p_{\mathbf{x}}, \prod_{i=1}^N p_{x_i} \right)$$

- If the *Kullback divergence* is used:

$$K(p_{\mathbf{x}}, p_{\mathbf{y}}) \stackrel{\text{def}}{=} \int p_{\mathbf{x}}(\mathbf{u}) \log \frac{p_{\mathbf{x}}(\mathbf{u})}{p_{\mathbf{y}}(\mathbf{u})} d\mathbf{u},$$

then we get the *Mutual Information* as an independence measure:

$$I(p_{\mathbf{x}}) = \int p_{\mathbf{x}}(\mathbf{u}) \log \frac{p_{\mathbf{x}}(\mathbf{u})}{\prod_{i=1}^N p_{x_i}(u_i)} d\mathbf{u}. \quad (13)$$



# Mutual Information: properties

- MI always positive
- Cancels if r.v. are mutually independent
- MI is invariant by scale change

Proof...

## ■ Example 10: Gaussian case

$$I(g_{\mathbf{x}}) = \frac{1}{2} \log \frac{\prod V_{ii}}{\det \mathbf{V}}$$

# Mutual Information: decomposition

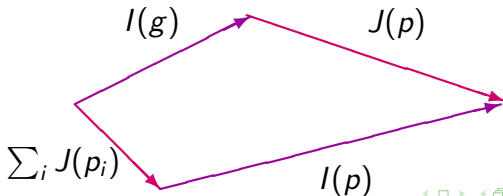
- Define the Negentropy as the divergence:

$$J(p_x) = K(p_x, g_x) = \int p_x(\mathbf{u}) \log \frac{p_x(\mathbf{u})}{g_x(\mathbf{u})} d\mathbf{u}. \quad (14)$$

Negentropy is invariant by invertible transforms

- Then MI can be decomposed into:

$$I(p_x) = I(g_x) + J(p_x) - \sum_i J(p_{x_i}). \quad (15)$$



# Sample Measures of Statistical Independence

## Independence at order $r$

- Definition:  
Components  $x_j$  of  $\mathbf{x}$  are independent at order  $r$  if all *cross cumulants* of order  $r$  are null
- In other words: the *Cumulant tensor*  $C_{ij\dots\ell}$  is diagonal.

### Example 11: Uncorrelated but not independent

$\mathbf{s}$  non Gaussian,  $s_i$  independent, then  $\mathbf{x} = \mathbf{Q}\mathbf{s}$  has uncorrelated components *at order 2* if  $\mathbf{Q}$  unitary  $\rightarrow$  cf. example slide 12.

# Edgeworth expansion of the MI

This yields for standardized random variables  $\mathbf{x}$ , after lengthy calculations:

$$I(p_{\mathbf{x}}) = J(p_{\mathbf{x}}) - \frac{1}{48} \sum_i 4C_{iii}^2 + C_{iiii}^2 + 7C_{iii}^4 - 6C_{iii}^2 C_{iiii} + o(m^{-2}). \quad (16)$$

- If 3rd order  $\neq 0$ , then  $I(p_{\mathbf{x}}) \approx J(p_{\mathbf{x}}) - \frac{1}{12} \sum_i C_{iii}^2$
- If 3rd order  $\approx 0$ , then  $I(p_{\mathbf{x}}) \approx J(p_{\mathbf{x}}) - \frac{1}{48} \sum_i C_{iiii}^2$

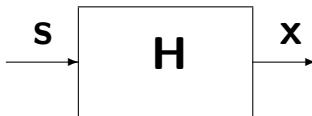
# Optimization Criteria

- Cumulant matching
- Contrast criteria
- Mutual Information
- Maximum Likelihood vs MI
- CoM family
- Other criteria

# Identification by Cumulant matching

## Principle

- Estimate the mixture by solving the I/O Multi-linear equations
- Apply a separating filter based on the latter estimate



# Noiseless mixture of 2 sources

**Example 12:  $2 \times 2$  by Cumulant matching** (cf. demo p.13)

- After standardization, the mixture takes the form

$$\mathbf{x} = \begin{pmatrix} \cos \alpha & -\sin \alpha e^{j\varphi} \\ \sin \alpha e^{-j\varphi} & \cos \alpha \end{pmatrix} \mathbf{s} \quad (17)$$

- Denote  $\gamma_{ij}^{k\ell} = \text{Cum}\{x_i, x_j, x_k^*, x_\ell^*\}$  and  $\kappa_i = \text{Cum}\{s_i, s_i, s_i^*, s_i^*\}$ .

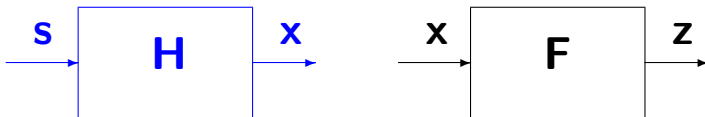
Then by *Multi-linearity*:

$$\begin{aligned} \gamma_{12}^{12} &= \cos^2 \alpha \sin^2 \alpha (\kappa_1 + \kappa_2) \\ \gamma_{11}^{12} &= \cos^3 \alpha \sin \alpha e^{j\varphi} \kappa_1 - \cos \alpha \sin^3 \alpha e^{j\varphi} \kappa_2 \\ \gamma_{12}^{22} &= \cos \alpha \sin^3 \alpha e^{j\varphi} \kappa_1 - \cos^3 \alpha \sin \alpha e^{j\varphi} \kappa_2 \end{aligned}$$

- Compact solution:  $\frac{\gamma_{12}^{22} - \gamma_{11}^{12}}{\gamma_{12}^{12}} = -2 \cot 2\alpha e^{j\varphi}$

# Now the inverse approach

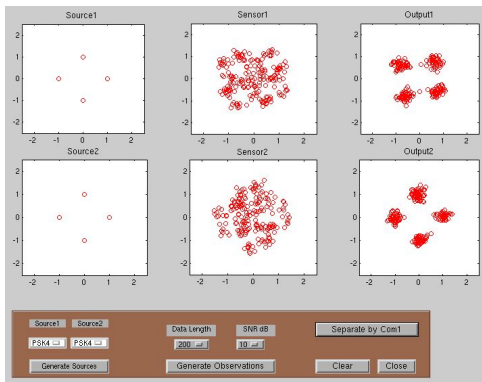
- Cumulant matching (direct approach: identification)
- Contrast Criteria (inverse approach: equalization):





# Noisy Mixtures of 2 sources

## Example 13: Separation of 2 non Gaussian sources by contrast maximization



# Source additional hypotheses

- **H1.** Each sources  $s_j[k]$  is an i.i.d. sequence, for any fixed  $j$
- H2.** Sources  $s_j$  are mutually statistically independent
- H3.** At most one source is Gaussian
- H4.** At most one source has a null marginal cumulant
- **H5.** Sources are Discrete, and belong to some known alphabet (but may be stat. dependent)
- H6.** Sources  $s_j[k]$  are sufficiently exciting
- **H7.** Sources are colored, and the set of source spectra forms a family of linearly independent functions
- **H8.** Sources are non stationary, and have different time profiles

# Trivial Filters

- They account for Inherent indeterminacies, remaining after assuming Source additional hypotheses

For instance:

- For dynamic (convolutive) mixtures, under **H1, H2, H3**,  $\check{\mathbf{T}}[z] = \mathbf{P} \check{\mathbf{D}}[z]$ , where  $\mathbf{P}$  is a permutation, and  $\check{\mathbf{D}}[z]$  a diagonal filter, with entries of the form  $\check{D}_{pp}[z] = \lambda_p z^{\delta_p}$ , where  $\delta_p$  is an integer.
- For static mixtures, under **H2, H3**,  $\mathbf{T} = \mathbf{P}\mathbf{D}$ , where  $\mathbf{P}$  permutation and  $\mathbf{D}$  diagonal invertible.
- In other words, if  $\mathbf{s}$  satisfies **Hi**, then so does  $\mathbf{T}\mathbf{s}$

# Contrast criteria: definition

## Axiomatic definition

A *Contrast* optimization criterion  $\Upsilon$  should enjoy 3 properties:

- *Invariance*:  $\Upsilon$  should not change under the action of trivial filters (as defined in slide 67)
- *Domination*: If sources are already separated, any filter should decrease (or leave unchanged)  $\Upsilon$
- *Discrimination*: The maximum achievable value should be reached only when sources are separated (i.e. all absolute maxima are related to each other by trivial filters)

# Mutual Information

$\gamma \stackrel{\text{def}}{=} -I(p_z)$  is a contrast

- Invariant by scale change and permutation
- Always negative
- Null if and only if components are independent

Proof... cf slide 57

# Likelihood

Given the source pdf's:  $p_s(\mathbf{u}) = \prod_i p_{s_i}(u_i)$ , and a sample  $\mathbf{x}_t$ , the ML approach consists of maximizing one of the criteria below w.r.t.  $\mathbf{H}$ :

## ■ Noiseless case

$$p_{\mathbf{x}|\mathbf{H}}(\mathbf{x}_T|\mathbf{H}) = \frac{1}{|\det \mathbf{H}|} p_s(\mathbf{H}^{-1}\mathbf{x})$$

## ■ Noisy case

$$p_{\mathbf{x},\mathbf{s}|\mathbf{H}}(\mathbf{x}_T, \mathbf{s}|\mathbf{H}) = g(\mathbf{x}_T - \mathbf{H}\mathbf{s}) p_s(\mathbf{s})$$

- And the *Joint MAP-ML* criterion for a joint estimation of sources:

$$\begin{aligned} (\mathbf{s}_{MAP}, \mathbf{H}_{MV}) &= \underset{\mathbf{s}, \mathbf{H}}{\text{Arg Max}} p_{\mathbf{x},\mathbf{s}|\mathbf{H}}(\mathbf{x}_T, \mathbf{s}|\mathbf{H}) \\ &= \underset{\mathbf{s}, \mathbf{H}}{\text{Arg Max}} p(\mathbf{x}_T|\mathbf{s}, \mathbf{H}) p_s(\mathbf{s}) \end{aligned}$$

# Noiseless Maximum Likelihood (1)

- For an increasing number of independent observations, the average log-likelihood converges to

$$\mathcal{L}_T \stackrel{\text{def}}{=} \frac{1}{T} \log p(\mathbf{x}_1 \dots \mathbf{x}_T | \mathbf{H}) \longrightarrow \mathcal{L}_\infty = \int p_{\mathbf{x}}(\mathbf{u}) \log p_{\mathbf{x}|\mathbf{H}}(\mathbf{u}) d\mathbf{u}$$

which can be seen to be, by making the change  $\mathbf{v} = \mathbf{H}^{-1}\mathbf{u}$ , and up to a constant:

$$\Upsilon_{ML} \stackrel{\text{def}}{=} \mathcal{L}_\infty - S(p_{\mathbf{x}}) = -K(p_{\mathbf{z}}, p_{\mathbf{s}}) \quad (18)$$

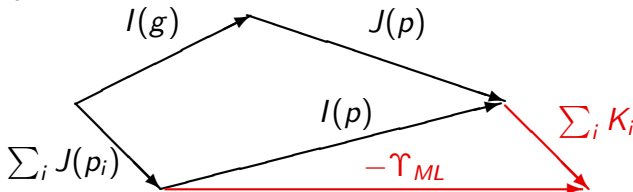
Proof...

## Noiseless maximum Likelihood (2)

- Yet, since  $s_i$  are independent, it can be shown that

$$K(p_z, p_s) = \underbrace{K(p_z, \prod_i p_{z_i})}_{MI} + \underbrace{\sum_i K(p_{z_i}, p_{s_i})}_{pdfdeviation}$$

This allows to take into account the source pdf's, if they are known



- **But** ML is not adequate if source pdf's are unknown  
 $\Rightarrow$  just use MI



# CoM Family of contrast functionss

When observations are standardized, and when only *unitary transforms* are considered, then the following are contrasts:

- If at most 1 source has a null skewness [COM94b]:

$$\Upsilon_{2,3} = \sum_{i=1}^P (\kappa_{iii})^2, \quad \kappa_{iii} \stackrel{\text{def}}{=} \mathcal{C}_{\mathbf{z} \, iii}$$

- If at most 1 source has a null kurtosis [COM94a]:

$$\Upsilon_{2,4} = \sum_{i=1}^P (\kappa_{ii}^{ii})^2, \quad \kappa_{ii}^{ii} \stackrel{\text{def}}{=} \mathcal{C}_{\mathbf{z} \, ii}^{ii}$$

- If at most 1 source has a null standardized Cumulant of order  $r \stackrel{\text{def}}{=} p + q > 2$ , and for any  $\alpha \geq 1$ :

$$\Upsilon_{\alpha,r} = \sum_{i=1}^P |\kappa_{i(p)}^{(q)}|^\alpha, \quad \kappa_{i(p)}^{(q)} \stackrel{\text{def}}{=} \text{Cum}\left\{ \underbrace{z_i, \dots, z_i}_{p \text{ times}}, \underbrace{z_i^*, \dots, z_i^*}_{q \text{ times}} \right\}$$

# General Family of contrasts

- **Theorem** All CoM contrasts belong to the larger family :

$$\Upsilon_g(\mathbf{z}) = \sum_i g(|\kappa_{i(p)}^{(q)}|) \quad (19)$$

where  $g(\cdot)$  is convex strictly increasing, and  $p + q > 2$ .

Proof...

# Contrast CoM(1, 4)

## Example 14: Kurtosis-based contrast without squaring

- In particular, if all source kurtosis have the same sign,  $\varepsilon$ , one can avoid the absolute value:

$$\Upsilon_{1,4} = \varepsilon \sum_{p=1}^P \kappa_{ii}^{ii}$$

Proof...

## Other criteria

- Contrasts based on Matrix slices of Cumulant tensor
- Contrasts dedicated to Discrete source alphabets
- Contrasts for convolutive mixtures - basically the same!

## Part II

# Applications

# Applications

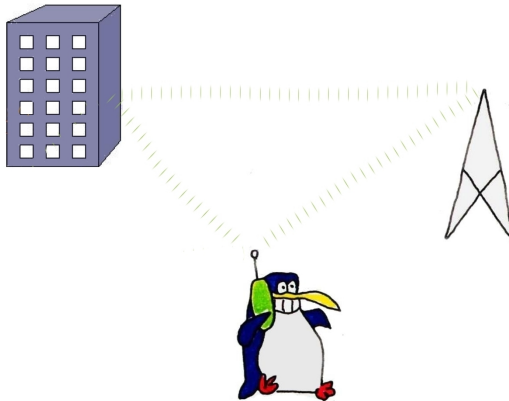
- Sensor Array Processing
- Telecommunications
- Speech
- Biomedical
- Machine Learning
- Exploratory Analysis...

# Application Areas (1)

## ■ Sensor Array Processing

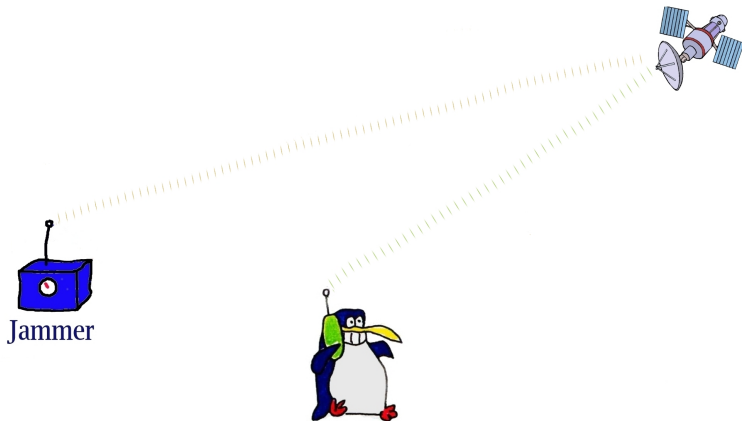
- Speech
- Localization with ill calibrated antennas
- Detection and/or extraction with unknown antennas  
(eg. sonar buoys, biomedical, audio, nuclear plants...)
- Blind extraction (eg. COMINT: interception, surveillance)
- Localization with reduced diversity (eg. Air traffic control)

# Telecommunications: SISO equalization

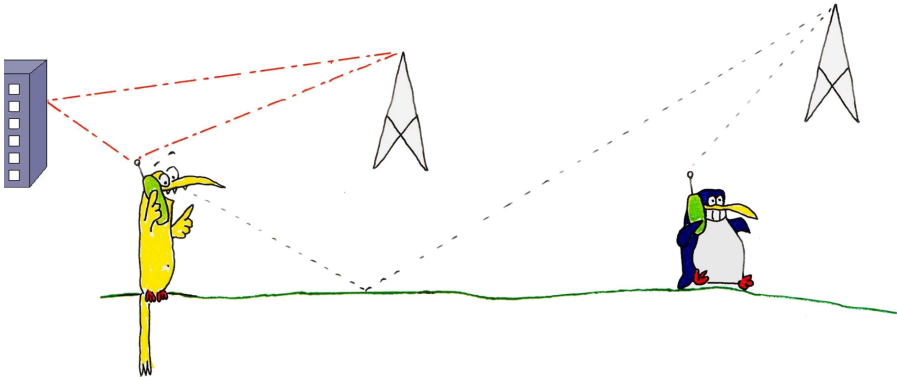




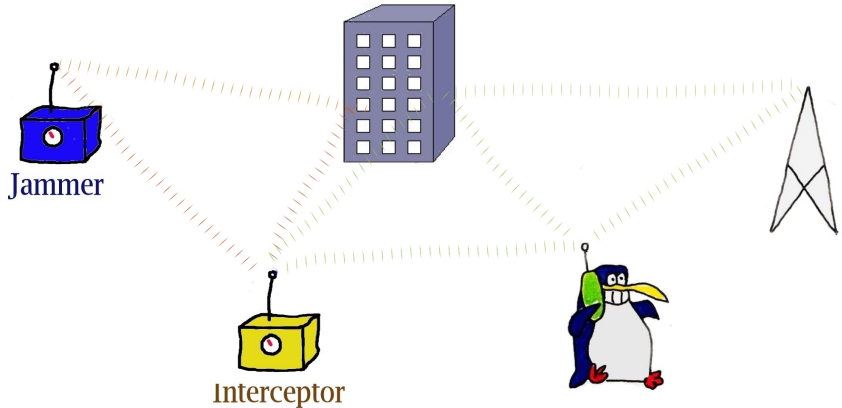
# Telecommunications: MISO equalization



# Telecommunications: MIMO equalization



# ComInt: MIMO equalization



# Speech

## The Cocktail Party problem



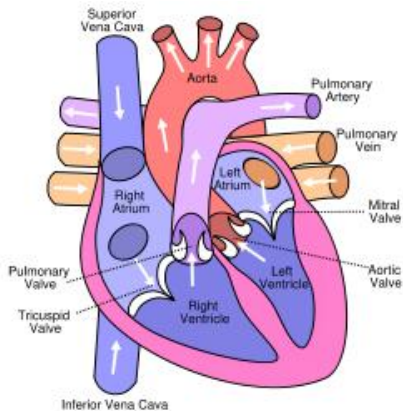
# Deconvolution

III focussing is a 2-D convolution: mixture with neighboring pixels.



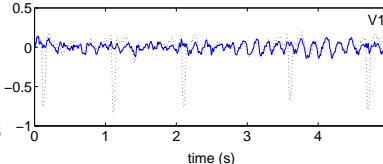
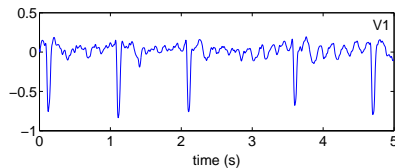
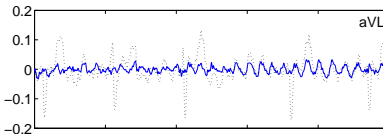
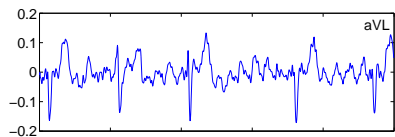
# Electrocardiography (1)

## Anatomy



# Electrocardiography (2)

## Atrial fibrillation [RCS<sup>+</sup>04]



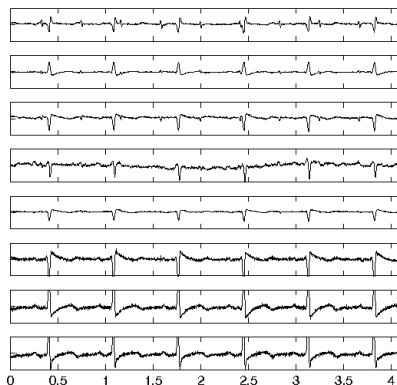
Atrial Fibrillation episode

Atrial activities

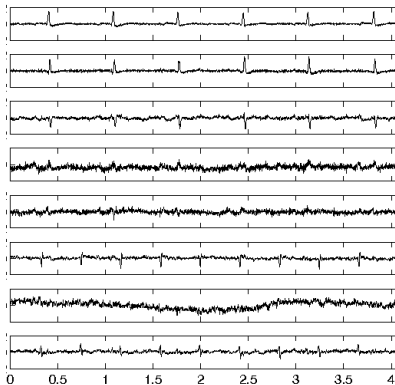
Extracted

# Electrocardiography (3)

## Mother-Phoetus separation [dLdMV00a]



Data



by ICA

Separati



# Application Areas (2)

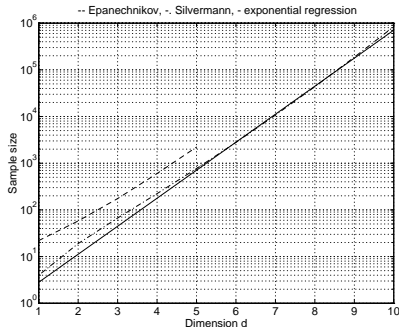
- Factor Analysis
  - Chemometrics
  - Econometrics
  - Psychology
- Compression
- Arithmetic Complexity
- Machine Learning
- Exploratory Analysis

# Machine Learning

## Curse of dimensionality



Number of samples required to reach a given relative error in pd.f. estimate,  $O(\epsilon)$ , is of order  $O(\epsilon^{-1-d/4})$  [SIL86]  $\Rightarrow$  exponential in  $d$



- Split of space into two lower dimensional subspaces allows the approximation of the p.d.f. [COM95]:

$$p_{\mathbf{x}}(\mathbf{u}) \approx p_{x_1}(\mathbf{u}_1) \cdot p_{x_2}(\mathbf{u}_2)$$

# Factor Analysis

Food Sciences:  
one of the numerous application areas



*judges × products × sensory properties*

## Part III

# Tensors

# Contents

- Introduction
- Canonical Decomposition (CanD), Tensor rank
- Symmetric tensors, Quantics, Topology
- Other tensors
- Tucker3, HOSVD
- Other decompositions

References

# Introduction

- Striking facts
- Order, dimensions, outer & inner products
- Contraction
- Multi-linearity property
- Unfoldings & storage
- Symmetry

## Striking facts

1. The rank of a matrix cannot be larger than its dimensions  $\rightarrow$  possible for a tensor
2. Matrices with entries drawn randomly have maximal rank  $\rightarrow$  not true for a tensor
3. The set of matrices of rank at most  $r$  is closed,  $\forall r$ ,  $\rightarrow$  not true for a tensor. Hence the approximation problem is generally ill-posed.
4. Worse: the maximal achievable rank of a tensor is generally still unknown.
5. There are several ways to extend the SVD to tensors
6. The computation of the rank of a given tensor still raises unsolved difficulties.
7. Rank and symmetric rank have not yet been proved to be the same
8. Subtraction of best rank-1 approximate does not necessarily decrease the rank

# Tensor product

- Let  $\mathcal{V}_\ell$  be vector spaces of dimension  $K_\ell$  on a field  $\mathbb{K}$ , and let  $\mathbf{v}_\ell \in \mathcal{V}_\ell$  be  $K_\ell$ -dimensional vectors.
- A *tensor*  $\mathbf{T}$  is an element of a tensor product  $\mathcal{V}_1 \otimes \mathcal{V}_2 \otimes \dots \otimes \mathcal{V}_P$ . For instance

$$\mathbf{v}_1 \otimes \mathbf{v}_2 \otimes \dots \otimes \mathbf{v}_P$$

is a tensor of *dimensions*  $K_1 \times K_2 \times \dots \times K_P$ .



# Arrays

- If coordinates of  $\mathbf{u} \in \mathcal{U}$ ,  $\mathbf{v} \in \mathcal{V}$  and  $\mathbf{w} \in \mathcal{W}$  are  $u_i$ ,  $v_j$ , and  $w_k$  in canonical bases of  $\mathcal{U}$ ,  $\mathcal{V}$  and  $\mathcal{W}$  respectively, then coordinates of tensor  $\mathbf{T} = \mathbf{u} \otimes \mathbf{v} \otimes \mathbf{w}$  are given by the array

$$T_{ijk} = u_i v_j w_k$$

- Given canonical bases, one often assimilates a *tensor* and its associated *array* of coordinates.

# Order & Dimensions

**Definitions** Let the array  $\mathbf{T} = \{T_{ij..k}\}$

- *Order* of  $\mathbf{T} \stackrel{\text{def}}{=} \#$  of its ways  $= \#$  of its indices
- *Dimension*  $K_\ell \stackrel{\text{def}}{=} \text{range of the } \ell\text{th index}$
- $\mathbf{T}$  is *Cubic* when all dimensions  $K_\ell = K$  are equal
- $\mathbf{T}$  is *Symmetric* when it is cubic and when its entries do not change by *any* permutation of indices

# Notation

- Let  $\mathbf{A}$  and  $\mathbf{B}$  be matrices of dimensions  $m_A \times n_A$  and  $m_B \times n_B$ , respectively
- Notation  $\mathbf{A} \circ \mathbf{B}$  will be preferred to  $\mathbf{A} \otimes \mathbf{B}$ , to avoid possible confusion with the Kronecker product  $\mathbf{A} \otimes \mathbf{B}$  between matrices. In fact:
  - 👉  $\mathbf{A} \otimes \mathbf{B}$  is a *matrix* of size  $m_A m_B \times n_A n_B$
  - 👉  $\mathbf{A} \otimes \mathbf{B}$  is a *tensor* of size  $m_A \times n_A \times m_B \times n_B$

# Matrix products

Again let  $\mathbf{A}$  and  $\mathbf{B}$  be matrices of dimensions  $m_A \times n_A$  and  $m_B \times n_B$ , with entries  $\{a_{ij}\}$  and  $\{b_{ij}\}$ , respectively

- **Kronecker product:**  $\mathbf{A} \otimes \mathbf{B}$  is  $m_A m_B \times n_A n_B$

$$\mathbf{A} \otimes \mathbf{B} \stackrel{\text{def}}{=} \begin{pmatrix} a_{11}\mathbf{B} & a_{12}\mathbf{B} & \cdots \\ a_{21}\mathbf{B} & a_{22}\mathbf{B} & \cdots \\ \vdots & \vdots & \ddots \end{pmatrix}$$

- **Khatri-Rao product** of matrices with same number of columns,  $n$ :

$$\mathbf{A} \odot \mathbf{B} \stackrel{\text{def}}{=} \begin{pmatrix} \mathbf{a}_1 \otimes \mathbf{b}_1 & \mathbf{a}_2 \otimes \mathbf{b}_2 & \cdots \end{pmatrix}$$

This is a *column-wise* Kronecker product.  $\mathbf{A} \odot \mathbf{B}$  is  $m_A m_B \times n$ .

# Outer product

- Outer or “tensor” product between two arrays,  $\mathbf{C} = \mathbf{A} \circ \mathbf{B}$ :

$$C_{ij..l\ ab..d} = A_{ij..l} B_{ab..d}$$

The *orders* add up

- **Example 15: Outer product between 2 vectors** The tensor

$$\mathbf{u} \circ \mathbf{v} = \mathbf{u} \mathbf{v}^T$$

has coordinates  $u_i v_j$  and is of order 2, and is hence a matrix.

# Arrays (cont'd)

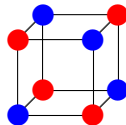
**Example 16:** Take

$$\mathbf{v} = \begin{pmatrix} 1 \\ -1 \end{pmatrix}$$

Then

$$\mathbf{v}^{\circ 3} \stackrel{\text{def}}{=} \mathbf{v} \circ \mathbf{v} \circ \mathbf{v} = \left( \begin{array}{cc|cc} 1 & -1 & -1 & 1 \\ -1 & 1 & 1 & -1 \end{array} \right)$$

This is a “rank-1” symmetric tensor



blue bullets = 1, red bullets = -1.

# Inner Product (1)

- **Mode-1 inner product:**  $\mathbf{A} \bullet_1 \mathbf{B}$ :

$$\{\mathbf{A} \bullet_1 \mathbf{B}\}_{i_2 \dots i_M j_2 \dots j_K} = \sum_k A_{ki_2 \dots i_M} B_{kj_2 \dots j_K}$$

This is a *contraction* on the 1st index

- **Mode- $p$  inner product:** similarly  $\mathbf{A} \bullet_p \mathbf{B}$  is obtained by summing up (i.e. contracting) on the  $p$ th index
- **Example 17: Matrix-vector product**  $\mathbf{A} \mathbf{u} = \mathbf{A}^\top \bullet_1 \mathbf{u}$
- **NB:**  
there exists a (less convenient & less used) other notation:  
 $\mathbf{A} \times_p \mathbf{B}$

## Inner Product (2)

- **Example 18: Matrix products** are *contractions*

$$\mathbf{A} \cdot \mathbf{B} = \mathbf{A} \bullet_2 \mathbf{B}^T = \mathbf{A}^T \bullet_1 \mathbf{B}$$

- **Example 19: Frobenius norm** of a  $P$ th order tensor in  $\mathbb{C}$ :

$$\|\mathbf{T}\|^2 = \sum_{i_1 i_2 \dots i_P} |T_{i_1 i_2 \dots i_P}|^2 = \mathbf{T} \bullet_1 \bullet_2 \dots \bullet_P \mathbf{T}^*$$

One contracts on all indices



## Inner Product (3)

- The *Contraction* is not associative

$$\mathbf{A} \bullet_1 (\mathbf{B} \bullet_1 \mathbf{C}) \neq (\mathbf{A} \bullet_1 \mathbf{B}) \bullet_1 \mathbf{C}$$

even for 2nd order tensors (matrices):  $\mathbf{A}^T \mathbf{B}^T \mathbf{C} \neq \mathbf{B}^T \mathbf{A} \mathbf{C}$

- **A convention** exists when a *single* tensor is contracted on several matrices, to avoid parentheses: the summation is always performed on the *second* matrix index.

**Example 20:** If  $\mathbf{A}$ ,  $\mathbf{B}$ ,  $\mathbf{C}$  are matrices, and  $\mathbf{T}$  a 3rd order tensor,

$$\mathbf{T}' = \mathbf{T} \bullet_1 \mathbf{A} \bullet_2 \mathbf{B} \bullet_3 \mathbf{C} \Rightarrow T'_{pqr} = \sum_{ijk} A_{pi} B_{qj} C_{rk} T_{ijk} \quad (20)$$

# Change of basis

Assume a change of basis is performed in every linear space  $\mathcal{V}_\ell$ , e.g. defined by matrix **A** in  $\mathcal{V}_1$ , **B** in  $\mathcal{V}_2$ , ... and **C** in  $\mathcal{V}_P$ .

- Multilinearity. An order- $P$  tensor **T** is transformed by the multi-linear map  $\{\mathbf{A}, \mathbf{B}, \dots, \mathbf{C}\}$  into a tensor **T'**:

$$T'_{ij\dots k} = \sum_{ab\dots c} A_{ia} B_{jb} \dots C_{kc} T_{ab\dots c}$$

- Compact writing (with convention of slide 105):

$$\mathbf{T}' = \mathbf{T} \bullet_1 \mathbf{A} \bullet_2 \mathbf{B} \dots \bullet_P \mathbf{C}$$

# Unfoldings (1)

- **Storage of a matrix in a vector** Let  $\mathbf{A}$  be a  $p \times q$  matrix, with columns  $A_{:j}$ . Then:

$$\text{vec}\{\mathbf{A}\} \stackrel{\text{def}}{=} \begin{bmatrix} A_{:1} \\ A_{:2} \\ \vdots \\ A_{:q} \end{bmatrix} \quad (21)$$

- Conversely,  $\mathbf{A} = \text{Unvec}_q(\text{vec}\{\mathbf{A}\})$ , if  $q$  denotes the # of columns
- **Storage of a tensor in a vector**  
Similarly, the linear operator  $\text{vec}\{\cdot\}$  maps a  $\alpha \times \beta \times \cdots \times \gamma$  tensor onto a vector ( $\alpha\beta \dots \gamma \times 1$  array)

## Unfoldings (2)

### ■ Storage of a tensor in a matrix

For a 3rd order tensor  $\mathbf{T}$ , one defines 3 *unfolding matrices*:

$$\mathbf{T}_{KI \times J} = \begin{bmatrix} \mathbf{T}_{::1} \\ \vdots \\ \mathbf{T}_{::k} \\ \vdots \\ \mathbf{T}_{::K} \end{bmatrix}, \quad \mathbf{T}_{IJ \times K} = \begin{bmatrix} \mathbf{T}_{1::} \\ \vdots \\ \mathbf{T}_{i::} \\ \vdots \\ \mathbf{T}_{I::} \end{bmatrix}, \quad \mathbf{T}_{JK \times I} = \begin{bmatrix} \mathbf{T}_{:1:}^T \\ \vdots \\ \mathbf{T}_{:j:}^T \\ \vdots \\ \mathbf{T}_{:J:}^T \end{bmatrix},$$

### ■ Conversely,

**Reshape** <sub>$I,K,J$</sub> ( $\mathbf{T}_{KI \times J}$ ), **Reshape** <sub>$J,I,K$</sub> ( $\mathbf{T}_{IJ \times K}$ ) or

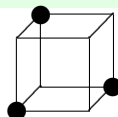
**Reshape** <sub>$K,J,I$</sub> ( $\mathbf{T}_{JK \times I}$ )

yield back  $\mathbf{T}$  up to a permutation of the modes.

### ■ Similar tools for higher orders...

$\ell$ -mode rank

- **Example 21:**  $2 \times 2 \times 2$ . Let  $\mathbf{T} =$



where bullets indicate nonzero entries, equal to 1 (see also slide 117). Then matrix unfoldings are

$$\mathbf{T}_{I \times JK} = \begin{pmatrix} 1 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 \end{pmatrix}$$

$$\mathbf{T}_{J \times KI} = \begin{pmatrix} 1 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 \end{pmatrix}$$

$$\mathbf{T}_{K \times IJ} = \begin{pmatrix} 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 \end{pmatrix}$$

- Note that  $\ell$  *mode ranks can be different*:  
 $\text{rank}_1 = \text{rank}_2 = 2 \neq \text{rank}_3 = 1$

# Canonical Decomposition

- Tensor rank
- Properties of the CanD
- Normalized CanD
- Matrix writings of the CanD
- Rank can exceed dimensions
- Field can change rank

# Tensor rank

- Any tensor or array  $\mathbf{T}$ , of dimensions  $I \times J \times \dots \times K$  can always be decomposed as

$$\mathbf{T} = \sum_q \mathbf{u}^{(q)} \circ \mathbf{v}^{(q)} \circ \dots \circ \mathbf{w}^{(q)}$$

- The *tensor rank* of  $\mathbf{T}$  is the minimal value of  $P$  such that equality holds  
This yields the *Canonical Decomposition* (CanD), sometimes referred to as *Parafac* decomposition:

$$\mathbf{T} = \sum_{q=1}^{\text{rank}\{\mathbf{T}\}} \mathbf{u}^{(q)} \circ \mathbf{v}^{(q)} \circ \dots \circ \mathbf{w}^{(q)} \quad (22)$$

- Tensor rank is always larger than or equal to all  $\ell$ -mode ranks:

$$\text{rank}_{\ell}\{\mathbf{T}\} \leq \text{rank}\{\mathbf{T}\}, \quad \forall \ell$$

# Other writings (1)

- Vectors can be normalized to unit norm, yielding a normalized version:

$$\mathbf{T} = \sum_{q=1}^{\text{rank}\{\mathbf{T}\}} \lambda_q \mathbf{u}^{(q)} \circ \mathbf{v}^{(q)} \circ \dots \circ \mathbf{w}^{(q)} \quad (23)$$

👉 Will be useful for symmetric tensors in the real field




## Other writings (2)

Let  $\mathbf{T}$  be a 3rd order tensor, and denote  $\mathbf{U}$ ,  $\mathbf{V}$ ,  $\mathbf{W}$  the matrices containing  $\mathbf{u}^{(p)}$ ,  $\mathbf{v}^{(p)}$ ,  $\mathbf{w}^{(p)}$  as columns.

- Assuming  $\mathbf{\Lambda}$  is a diagonal tensor of same order  $P$  as  $\mathbf{T}$ , with entries  $\lambda_q$ , the normalized CanD (23) admits a writing by contractions, with convention (20) of slide 105:

$$\mathbf{T} = \mathbf{\Lambda} \underset{1}{\bullet} \mathbf{U} \underset{2}{\bullet} \mathbf{V} \dots \underset{P}{\bullet} \mathbf{W}$$

In other words, the CanD is a means to model a tensor as a transformation from a *diagonal* one.

 **Warning:** matrices  $\mathbf{U}$ ,  $\mathbf{V}$ , ...  $\mathbf{W}$  may not be invertible nor even square!

## Other writings (3)

- The CanD (22) can be written in matrix form:

$$\mathbf{T}_{I \times JK} = \mathbf{U} (\mathbf{W} \odot \mathbf{V})^T \quad (24)$$

- Alternatively, each matrix slice of  $\mathbf{T}$  can be written as

$$\mathbf{T}_{::k} = \mathbf{U} \text{Diag}\{\mathbf{W}(k, :)\} \mathbf{V}^T \quad (25)$$

**NB:** This extends to any order. In particular at order 4, with appropriate notations:

$$\mathbf{T}_{::k\ell} = \mathbf{A} \text{Diag}\{\mathbf{C}(k, :)\} \text{Diag}\{\mathbf{D}(\ell, :)\} \mathbf{B}^T$$

# Properties

- The CanD of a *multilinear transform* is the *transformed CanD*:  
 If  $\mathbf{T} \stackrel{\text{def}}{=} \mathbf{\Lambda} \bullet_1 \mathbf{U} \bullet_2 \mathbf{V} \bullet_3 \mathbf{W}$  is transformed into  
 $\mathbf{T}' = \mathbf{T} \bullet_1 \mathbf{A} \bullet_2 \mathbf{B} \bullet_3 \mathbf{C}$ ,  
 then  $\mathbf{T}'$  admits the CanD:

$$\mathbf{T}' = \mathbf{\Lambda} \bullet_1 (\mathbf{A} \mathbf{U}) \bullet_2 (\mathbf{B} \mathbf{V}) \bullet_3 (\mathbf{C} \mathbf{W})$$

- The CanD is valid in a ring (only multiplies)

# Examples (1)

## ■ Example 22: $2 \times 2 \times 2$ tensor of rank 2

$$\mathbf{T} = \left( \begin{array}{cc|cc} 1 & 2 & 2 & 4 \\ 3 & 4 & 6 & 8 \end{array} \right) = \begin{pmatrix} 1 \\ 1 \end{pmatrix} \circ \begin{pmatrix} 1 \\ 2 \end{pmatrix} \circ \begin{pmatrix} 1 \\ 2 \end{pmatrix} + \begin{pmatrix} 0 \\ 2 \end{pmatrix} \circ \begin{pmatrix} 1 \\ 1 \end{pmatrix} \circ \begin{pmatrix} 1 \\ 2 \end{pmatrix}$$

here matrix slices are proportional

## ■ Example 23: $2 \times 2 \times 2$ of rank 2

$$\mathbf{T} = \left( \begin{array}{cc|cc} 1 & 2 & 2 & 4 \\ 3 & 4 & 4 & 6 \end{array} \right) = \begin{pmatrix} 1 \\ 1 \end{pmatrix} \circ \begin{pmatrix} 1 \\ 2 \end{pmatrix} \circ \begin{pmatrix} 1 \\ 2 \end{pmatrix} + \begin{pmatrix} 0 \\ 2 \end{pmatrix} \circ \begin{pmatrix} 1 \\ 1 \end{pmatrix} \circ \begin{pmatrix} 1 \\ 1 \end{pmatrix}$$

even if matrix slices are not proportional

## Examples (2)

### ■ Example 24: $2 \times 2 \times 2$ tensor of rank 3 [COM02b]

$$\mathbf{T} = \left( \begin{array}{cc|cc} 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 0 \end{array} \right)$$

and

$$2\mathbf{T} = \begin{pmatrix} 1 \\ 1 \end{pmatrix}^{\circ 3} + \begin{pmatrix} -1 \\ 1 \end{pmatrix}^{\circ 3} + 2 \begin{pmatrix} 0 \\ -1 \end{pmatrix}^{\circ 3}$$

👉 This is the maximal rank in dimension  $2 \times 2 \times 2$

👉 Here we have  $\text{rank}_3 = 1 < \text{rank}_1 = \text{rank}_2 = 2 < \text{rank}\{\mathbf{T}\}$  (cf. slide 109).

**NB:** Other writing:  $6x^2y = (x+y)^3 + (-x+y)^3 - 2y^3$

# Field can change rank

- We have for any real tensor  $\mathbf{T}$

$$\text{rank}\{\mathbf{T}\}_{\mathbb{C}} \leq \text{rank}\{\mathbf{T}\}_{\mathbb{R}}$$

**Example 25:** A  $2 \times 2 \times 2$  tensor of rank 3 in  $\mathbb{R}$ , but 2 in  $\mathbb{C}$  [CMLG06]

$$\mathbf{T} = \left( \begin{array}{cc|cc} -1 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{array} \right)$$

In fact

$$\mathbf{T} = \frac{1}{2} \begin{pmatrix} 1 \\ 1 \end{pmatrix}^{\circ 3} + \frac{1}{2} \begin{pmatrix} 1 \\ -1 \end{pmatrix}^{\circ 3} + 2 \begin{pmatrix} -1 \\ 0 \end{pmatrix}^{\circ 3} = \frac{j}{2} \begin{pmatrix} -j \\ 1 \end{pmatrix}^{\circ 3} - \frac{j}{2} \begin{pmatrix} j \\ 1 \end{pmatrix}^{\circ 3}$$

# Symmetric tensors

- Usefulness
- Symmetric rank
- Link with quantics
- Why rank can exceed dimension
- Generic & typical ranks
- Clebsh's statement
- Topology
- Hirschowitz theorem

# Usefulness of symmetric tensors

- They occur as derivatives of a multivariate function
  - Moments
  - Cumulants
  - Hessian



# Space of symmetric tensors

- $\mathcal{S}_K$ : symmetric tensors of dimensions  $K$  and order  $d$   
👉 space of dimension  $D_S(K, d) = \binom{K+d-1}{d}$

| $K \backslash d$ | quadratic<br>2 | cubic<br>3 | quartic<br>4 | quintic<br>5 | sextic<br>6 |
|------------------|----------------|------------|--------------|--------------|-------------|
| 2                | 3              | 4          | 5            | 6            | 7           |
| 3                | 6              | 10         | 15           | 21           | 28          |
| 4                | 10             | 20         | 35           | 56           | 84          |
| 5                | 15             | 35         | 70           | 126          | 210         |
| 6                | 21             | 56         | 126          | 252          | 462         |

Number of free parameters in a symmetric tensor of order  $d$  and dimension  $K$

- $\mathcal{A}_K$ : general tensors of dimensions  $K_\ell = K$ ,  $1 \leq \ell \leq d$   
👉 space of dimension  $D_A(K, d) = K^d$

# Symmetric rank

- **Definition** For decomposing a *symmetric* tensor, one can impose symmetry of each rank-1 term. Hence the *symmetric rank*:

$$\mathbf{T} = \sum_{q=1}^{\text{rank}_s(\mathbf{T})} [\mathbf{u}^{(q)}]^{\circ P}$$

- **Property** We have that

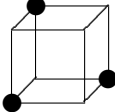
$$\text{rank}\{\mathbf{T}\} \leq \text{rank}_s \mathbf{T}, \quad \forall \mathbf{T} \text{ symmetric}$$

- It is not yet proved that both coincide for all values of order and dimensions:  
this is a conjecture [CGLM08].

# Link with quantics (1)

- A *quantic* is a homogeneous polynomial in several variables.  
For instance: quadric, cubic, quartic...
- **Example 26: Binary cubic**  $(d, K) = (3, 2)$   
Take again example in slide 117:

$$p(x_1, x_2) = \sum_{i,j,k=1}^2 T_{ijk} x_i x_j x_k$$

$$\mathbf{T} = \left( \begin{array}{cc|cc} 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 0 \end{array} \right) = \text{cube diagram}$$


$$\Rightarrow p(\mathbf{x}) = 3x_1^2x_2 = 3\mathbf{x}^{[2,1]}$$

## Link with quantics (2)

- **Bijection:** Symmetric tensor of order  $d$  and dimension  $K \leftrightarrow$  quantic of degree  $d$  in  $K$  variables:

$$p(\mathbf{x}) = \sum_{\mathbf{j}} T_{\mathbf{j}} \mathbf{x}^{\mathbf{f}(\mathbf{j})} \quad (26)$$

- integer vector  $\mathbf{j}$  of dimension  $d \leftrightarrow$  integer vector  $\mathbf{f}(\mathbf{j})$  of dimension  $K$
- entry  $f_k$  of  $\mathbf{f}(\mathbf{j})$  being  $\stackrel{\text{def}}{=} \#$  of times index  $k$  appears in  $\mathbf{j}$
- We have in particular  $|\mathbf{f}(\mathbf{j})| = d$ .
- Standard conventions:  $\mathbf{x}^{\mathbf{j}} \stackrel{\text{def}}{=} \prod_{k=1}^K x_k^{j_k}$  and  $|\mathbf{f}| \stackrel{\text{def}}{=} \sum_{k=1}^K f_k$ , where  $\mathbf{j}$  and  $\mathbf{f}$  are integer vectors.

# Literature

Gauss'1825  
Sylvester'1851  
Cayley'1854  
Clebsch'1861  
Salmon'1874  
Poincaré'1890  
Hilbert'1900  
Wakeford'1918  
Grothendieck'1966  
Dieudonné'1970  
Shafarevich'1975

Ehrenborg, Kogan...

## Why rank can exceed dimension

**Theorem** Let  $\mathbf{v}_{(1)}, \mathbf{v}_{(2)}, \dots, \mathbf{v}_{(r)}$ , be  $r$  *pairwise* linearly independent vectors, then for all  $k \geq r - 1$ , the rank-1 symmetric tensors are *linearly independent*:

$$\mathbf{v}_{(1)}^{\circ k}, \mathbf{v}_{(2)}^{\circ k}, \dots, \mathbf{v}_{(r)}^{\circ k}$$

### Example 27: 3 vectors in dimension 2

$$\mathbf{v}_{(1)} = \begin{pmatrix} 1 \\ 0 \end{pmatrix}, \mathbf{v}_{(2)} = \begin{pmatrix} 0 \\ 1 \end{pmatrix}, \mathbf{v}_{(3)} = \begin{pmatrix} 1 \\ 1 \end{pmatrix}$$

are pairwise linearly independent, but matrix of  $\{\mathbf{v}_{(q)}^{\circ 2}\}$  is full rank:

$$\begin{pmatrix} 1 & 0 & 1 \\ 0 & 0 & 1 \\ 0 & 0 & 1 \\ 0 & 1 & 1 \end{pmatrix}$$

# Orbits (1)

- General Linear group  $\mathcal{GL}$ : group of invertible matrices
- Orbit of a polynomial  $p$ : all polynomials  $q$  that can be transformed into  $p$  by  $\mathbf{A} \in \mathcal{GL}$ :  $q(\mathbf{x}) = p(\mathbf{Ax})$ .
- Allows to classify polynomials

# Orbits (2)

## Example 28: Quadrics

- Binary quadrics are associated with  $2 \times 2$  symmetric matrices (tensors of order 2)
  - Orbits in  $\mathbb{R}$ :  $\{0, x^2, x^2 + y^2, x^2 - y^2\}$ 
    - ☞  $2xy \in \mathcal{O}(x^2 - y^2)$  in  $\mathbb{R}[x, y]$
  - Orbits in  $\mathbb{C}$ :  $\{0, x^2, x^2 + y^2\}$ 
    - ☞  $2xy \in \mathcal{O}(x^2 + y^2)$  in  $\mathbb{C}[x, y]$
- Set of singular matrices is closed
- Set  $\mathcal{Y}_r$  of matrices of at most rank  $r$  is closed



$$3 \times 3$$

Classification of ternary quadrics

**Orbits in  $\mathbb{C}$ :**

| $\mathcal{GI}$ -orbit | $\omega(p)$          |
|-----------------------|----------------------|
| 0                     | 0                    |
| $x^2$                 | 1                    |
| $x^2 + y^2$           | 2                    |
| $x^2 + y^2 + z^2$     | 3 ( <b>generic</b> ) |

**Question:** what is the answer in  $\mathbb{R}$ ?

# CanD of polynomials

By using bijection (26), decomposing a  $d$ th order symmetric tensor into a sum of rank-1 tensors means

$$p(\mathbf{x}) = \sum_{q=1}^{r(p)} (\mathbf{v}_{(q)}^T \mathbf{x})^d \quad (27)$$

- This is a sum of powers of linear forms.
- $r(p)$  coincides with the rank of associated tensor
- $r(p)$  is sometimes called the *width* of  $p$  [REZ92].

# Generic & Typical Ranks

- **Informal definition** A property is *typical* if it holds true on a non-zero-volume set
- **Informal definition** A property is *generic* if it is true almost everywhere.
- There can be *several* typical ranks, but only *a single* generic rank.

# Bounds on generic rank (1)

For quantics of degree  $d$  in  $K$  variables

- Lower bound

$$\left\lceil \frac{\binom{K+d-1}{d}}{K} \right\rceil \leq \overline{R}$$

- Upper bound [Reznick'92]

$$\overline{R} \leq \binom{K+d-2}{d-1}$$

# Bounds on generic rank (2)

- Tensors of order  $d$  and dimensions  $(K_1, ..K_d)$  without symmetry:

- Upper bound

$$\left\lceil \frac{\prod_{i=1}^d K_i}{1 + \sum_{i=1}^d (K_i - 1)} \right\rceil \leq \overline{R}$$

- Square case  $K_i = K$ :

$$K^d / (dK - d + 1) \leq \overline{R}$$

- Lower bound (Square case):

$$K^d / (dK - d + 1) \leq \overline{R}$$

# Topology of quantics

- Every elementary closed set  $\stackrel{\text{def}}{=}$  varieties, defined by  $p(\mathbf{x}) = 0$
- Closed sets = finite union of varieties
- Closure of a set  $\mathcal{E}$ : smallest closed set  $\overline{\mathcal{E}}$  containing  $\mathcal{E}$

➡ is called the *Zariski* topology in  $\mathbb{C}$  [CLO92]

➡ this is not Euclidian topology, but results still apply [CGLM08]:  
Tensors with *entries randomly drawn* according to a continuous pdf  
are generic

# Clebsch's statement



Alfred Clebsch (1833-1872)

The generic ternary quartic cannot be written as the sum of 5 fourth powers

- $D(3, 4) = 15$
- $3r$  free parameters in the CAND
- But  $r = 5$  is not enough  $\rightarrow r = 6$  is generic !

# Tensor subsets

- Set of tensors of rank *at most*  $r$  with values in  $\mathbb{C}$ :

$$\mathcal{Y}_r = \{\mathbf{T} \in \mathcal{T} : r(\mathbf{T}) \leq r\}$$

- Set of tensors of rank *exactly*  $r$ :  $\mathcal{Z}_r = \{\mathbf{T} \in \mathcal{T} : r(\mathbf{T}) = r\}$

$$\mathcal{Z} = \mathcal{Y}_r - \mathcal{Y}_{r-1}, \quad r > 1$$

- Zariski closures:  $\overline{\mathcal{Y}}_r, \overline{\mathcal{Z}}_r$  ?



# Lack of closeness of $\mathcal{Z}_r$

## ■ PROPOSITION

$\mathcal{Z}_1$  is closed, *but not*  $\mathcal{Z}_r$  for any  $r > 1$

[Burgisser'97] [Strassen'83]

## ■ Proof

If  $\text{rank}\{\mathbf{T}\} > 1$ , there exist  $\mathbf{T}_0 \in \mathcal{Z}_{r-1}$  and  $\mathbf{y} \neq 0$  such that

$$\mathbf{T} = \mathbf{T}_0 + \mathbf{y}^{\circ d}$$

Then define  $\mathbf{T}_\varepsilon = \mathbf{T}_0 + \varepsilon \mathbf{y}^{\circ d}$ . This series converges to  $\mathbf{T}_0 \notin \mathcal{Z}_r$  as  $\varepsilon \rightarrow 0$

# Lack of closeness of $\mathcal{Y}_r$ (1)

## ■ PROPOSITION

If  $d > 2$ ,  $\mathcal{Y}_r$  is not closed for  $1 < r < R$ .

## ■ Example 29: Sequence of rank-2 tensors converging towards a rank-4:

$$\mathbf{T}_\varepsilon = \frac{1}{\varepsilon} [(\mathbf{u} + \varepsilon \mathbf{v})^{\circ 4} - \mathbf{u}^{\circ 4}]$$

In fact, as  $\varepsilon \rightarrow 0$ , it tends to:

$$\mathbf{T}_0 = \mathbf{u} \circ \mathbf{u} \circ \mathbf{u} \circ \mathbf{v} + \mathbf{u} \circ \mathbf{u} \circ \mathbf{v} \circ \mathbf{u} + \mathbf{u} \circ \mathbf{v} \circ \mathbf{u} \circ \mathbf{u} + \mathbf{v} \circ \mathbf{u} \circ \mathbf{u} \circ \mathbf{u}$$

which can be shown to be proportional to the rank-4 tensor:

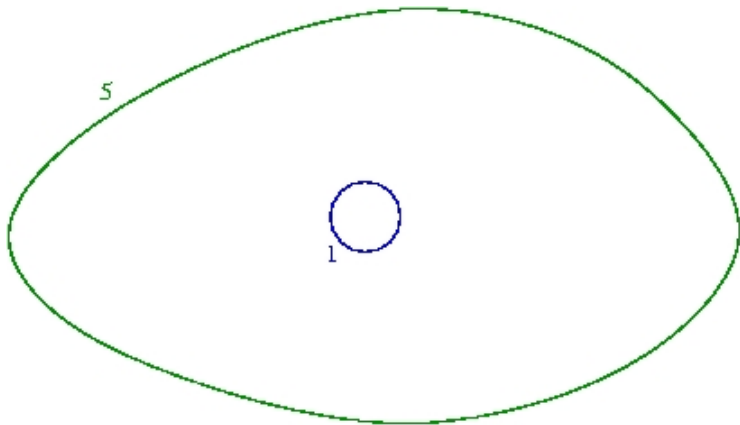
$$3\mathbf{T}_0 = 8(\mathbf{u} + \mathbf{v})^{\circ 4} - 8(\mathbf{u} - \mathbf{v})^{\circ 4} - (\mathbf{u} + 2\mathbf{v})^{\circ 4} + (\mathbf{u} - 2\mathbf{v})^{\circ 4} \quad (28)$$

where  $\mathbf{u}$  and  $\mathbf{v}$  are not collinear.

👉 This is the *maximal rank* of 4th order tensors of dimension 2.

# Lack of closeness of $\mathcal{Y}_r$ (2)

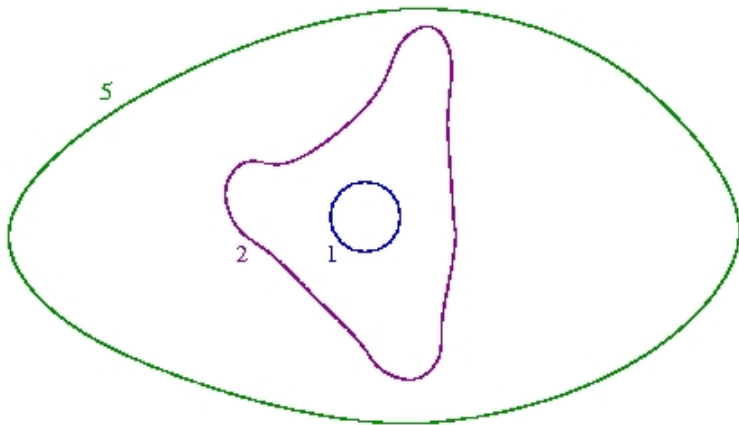
Successive sets  $\mathcal{Y}_r = \{\mathbf{T} : \text{rank}(\mathbf{T}) \leq r\}$



A tensor sequence in  $\mathcal{Y}_r$  can converge to a limit in  $\mathcal{Y}_{r+h}$

# Lack of closeness of $\mathcal{Y}_r$ (2)

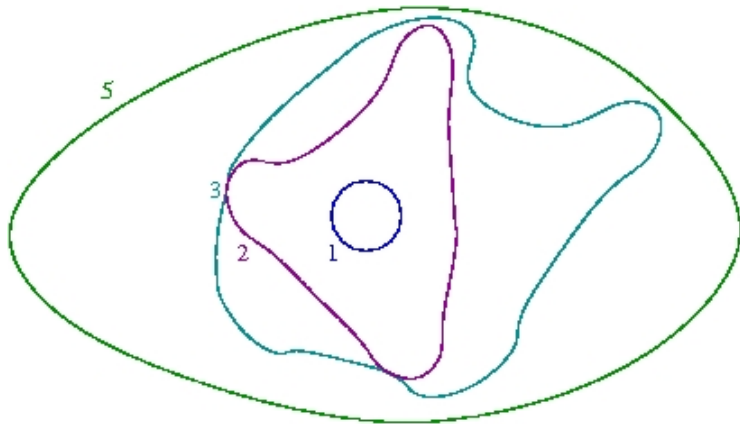
Successive sets  $\mathcal{Y}_r = \{\mathbf{T} : \text{rank}(\mathbf{T}) \leq r\}$



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# Lack of closeness of $\mathcal{Y}_r$ (2)

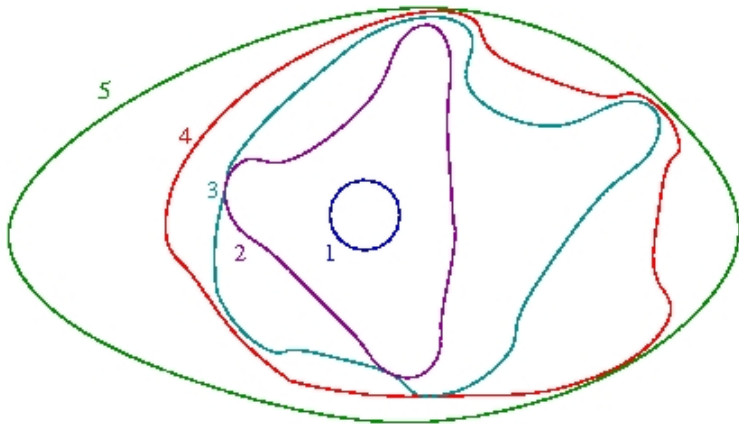
Successive sets  $\mathcal{Y}_r = \{\mathbf{T} : \text{rank}(\mathbf{T}) \leq r\}$



A tensor sequence in  $\mathcal{Y}_r$  can converge to a limit in  $\mathcal{Y}_{r+h}$

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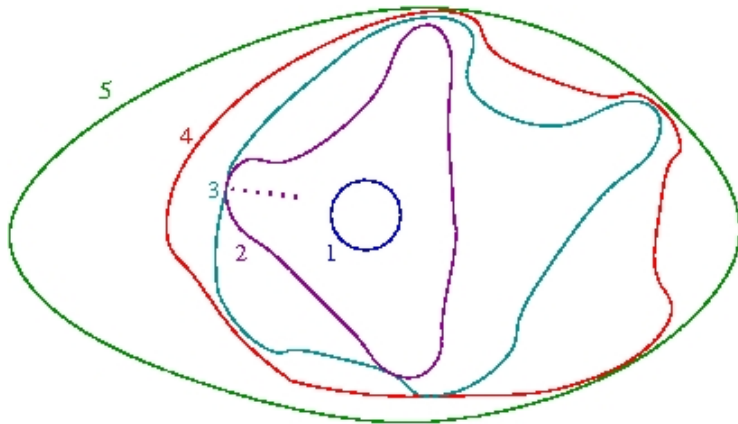
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# Lack of closeness of $\mathcal{Y}_r$ (2)

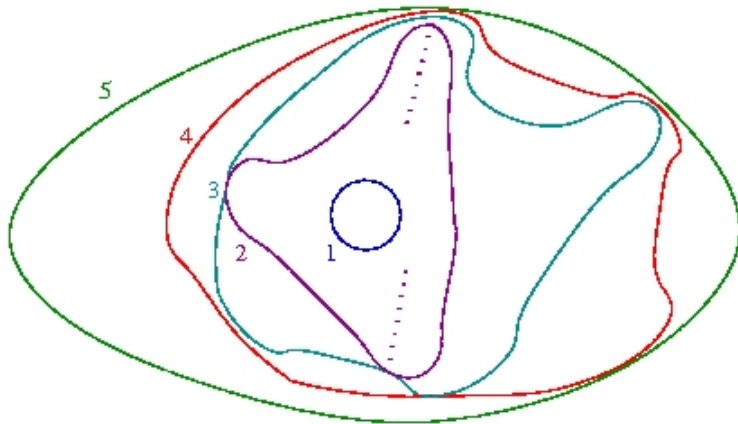
Successive sets  $\mathcal{Y}_r = \{\mathbf{T} : \text{rank}(\mathbf{T}) \leq r\}$



A tensor sequence in  $\mathcal{Y}_r$  can converge to a limit in  $\mathcal{Y}_{r+h}$

# Lack of closeness of $\mathcal{Y}_r$ (2)

Successive sets  $\mathcal{Y}_r = \{\mathbf{T} : \text{rank}(\mathbf{T}) \leq r\}$



A tensor sequence in  $\mathcal{Y}_r$  can converge to a limit in  $\mathcal{Y}_{r+h}$

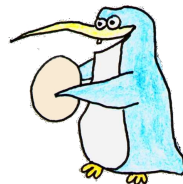


# Genericity

- **Formal definition**  $r$  is a typical rank if (density argument with Zariski):

$\overline{\mathcal{Z}}_r$  is the whole space

- **Formal definition** Generic rank is *the typical rank when unique*
- In  $\mathbb{C}$  a typical rank is unique, and hence generic
- For given values of order  $d$  and dimension  $K$ , the smallest typical rank in  $\mathbb{R}$  coincides with the generic rank in  $\mathbb{C}$



# Existence of the generic rank in $\mathbb{C}$

- **LEMMA** The series of  $\overline{\mathcal{Y}}_k$  is strictly increasing for  $k \leq \overline{R}$  and then constant:

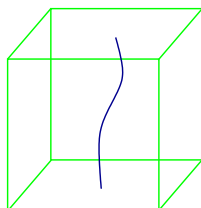
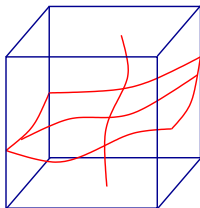
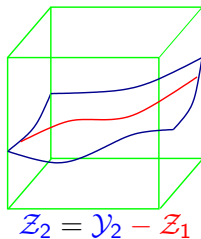
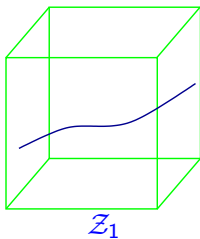
$$\overline{\mathcal{Y}}_1 \subsetneq \overline{\mathcal{Y}}_2 \subsetneq \dots \subsetneq \overline{\mathcal{Y}}_{\overline{R}} = \overline{\mathcal{Y}}_{\overline{R}+1} = \dots \mathcal{T}$$

which guarantees the existence of a unique  $\overline{R}$

- **PROPOSITION** For tensors in  $\mathbb{C}$   
If  $r_1 < r_2 < \overline{R} < r_3 \leq R$ , then

$$\overline{\mathcal{Z}}_{r_1} \subset \overline{\mathcal{Z}}_{r_2} \subset \overline{\mathcal{Z}}_{\overline{R}} \supset \overline{\mathcal{Z}}_{r_3} \supseteq \overline{\mathcal{Z}}_R \quad (29)$$

➡ Proves that  $\overline{R}$  is the generic rank in  $\mathbb{C}$

Generic rank in  $\mathbb{C}$ 

$$\begin{aligned}\mathcal{Z}_3 &= \mathcal{Y}_3 - \mathcal{Z}_1 - \mathcal{Z}_2 \\ &= \mathcal{T} - \mathcal{Z}_1 - \mathcal{Z}_2 - \mathcal{Z}_4\end{aligned}$$

$$\mathcal{Z}_4 = \mathcal{Y}_4 - \mathcal{Y}_3$$

# Numerical computation of the Generic Rank (1)

## Mapping

$$\begin{aligned} \{\mathbf{u}(\ell), 1 \leq \ell \leq r\} &\xrightarrow{\varphi} \sum_{\ell=1}^r \mathbf{u}(\ell)^{\circ d} \\ \{\mathbb{C}^K\}^r &\xrightarrow{\varphi} \mathcal{S} \end{aligned}$$

## Rank

The rank of the Jacobian of  $\varphi$  equals  $\dim(\bar{\mathcal{Z}}_r)$ , and hence  $D$  for large enough  $r$ .

➡ The **smallest**  $r$  for which  $\text{rank}(\text{Jacobian}(\varphi)) = D$  is  $\bar{R}$ .

# Numerical computation of the Generic Rank (2)

## Example 30: 3rd order symmetric tensors

$$\{\mathbf{u}(\ell), 1 \leq \ell \leq r\} \xrightarrow{\varphi} \mathbf{T} = \sum_{\ell=1}^r \mathbf{u}(\ell)^{\circ 3}$$

$\mathbf{T}$  has coordinate vector:  $\sum_{\ell=1}^r \mathbf{u}(\ell) \otimes \mathbf{u}(\ell) \otimes \mathbf{u}(\ell)$ . Hence the Jacobian of  $\varphi$  is the  $r n \times n^3$  matrix:

$$\mathbf{J} = \begin{bmatrix} \mathbf{I}_n \otimes \mathbf{u}^T(1) \otimes \mathbf{u}^T(1) + \mathbf{u}(1)^T \otimes \mathbf{I}_n \otimes \mathbf{u}^T(1) + \mathbf{u}(1)^T \otimes \mathbf{u}(1)^T \otimes \mathbf{I}_n \\ \mathbf{I}_n \otimes \mathbf{u}^T(2) \otimes \mathbf{u}^T(2) + \mathbf{u}(2)^T \otimes \mathbf{I}_n \otimes \mathbf{u}^T(2) + \mathbf{u}(2)^T \otimes \mathbf{u}(2)^T \otimes \mathbf{I}_n \\ \dots + \dots + \dots \\ \mathbf{I}_n \otimes \mathbf{u}^T(r) \otimes \mathbf{u}^T(r) + \mathbf{u}(r)^T \otimes \mathbf{I}_n \otimes \mathbf{u}^T(r) + \mathbf{u}(r)^T \otimes \mathbf{u}(r)^T \otimes \mathbf{I}_n \end{bmatrix}$$

and

$$\begin{cases} \text{rank}\{\mathbf{J}\} = \dim(\text{Im}(\varphi)) \\ \bar{R} = \text{Min}\{r : \text{Im}\{\varphi\} = \mathcal{S}\} \end{cases}$$

# Numerical computation of the Generic Rank (3)

The symmetric rank is generically:

| $d \backslash K$ | 2 | 3        | 4         | 5         | 6  | 7  | 8  |
|------------------|---|----------|-----------|-----------|----|----|----|
| 3                | 2 | 4        | 5         | <b>8</b>  | 10 | 12 | 15 |
| 4                | 3 | <b>6</b> | <b>10</b> | <b>15</b> | 21 | 30 | 42 |

$$\bar{R}_s \geq \frac{1}{K} \binom{K+d-1}{d}$$

**Bold:** exceptions to the ceil rule:  $\bar{R}_s \geq \lceil \frac{1}{K} \binom{K+d-1}{d} \rceil$   
[CM96]

# Uniqueness of CanD

Number of solutions

- The fiber of solutions has dimension

$$F(n) = K \bar{R} - \binom{K + d - 1}{d}$$

| $d \quad K$ | 2 | 3        | 4        | 5        | 6 | 7 | 8 |
|-------------|---|----------|----------|----------|---|---|---|
| 3           | 0 | 2        | 0        | <b>5</b> | 4 | 0 | 0 |
| 4           | 1 | <b>3</b> | <b>5</b> | <b>5</b> | 0 | 0 | 6 |

➡ 0 means a finite number of solutions

# Hirschowitz theorem

From Alexander-Hirschowitz theorem (cf. appendix), one can deduce [CGLM08]:

**THEOREM** For  $d > 2$ , the generic rank of a  $d$ th order symmetric tensor of dimension  $K$  is **always** equal to the lower bound

$$\bar{R}_s = \left\lceil \frac{\binom{K+d-1}{d}}{K} \right\rceil \quad (30)$$

**except** for the following cases:

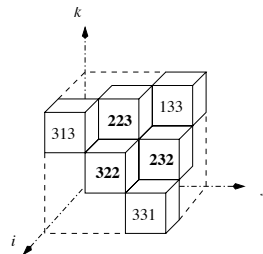
$(d, K) \in \{(3, 5), (4, 3), (4, 4), (4, 5)\}$ , for which it should be increased by 1.

➡ Only a *finite number* of exceptions !



# Classification of ternary cubics

| $\mathcal{GI}$ -orbit        | $\omega(p)$ |
|------------------------------|-------------|
| $x^3$                        | 1           |
| $x^3 + y^3$                  | 2           |
| $x^2y$                       | 3           |
| $x^3 + 3y^2z$                | 4           |
| $x^3 + y^3 + 6xyz$           | 4           |
| $x^3 + 6xyz$                 | 4           |
| $a(x^3 + y^3 + z^3) + 6bxyz$ | 4 (generic) |
| $xz^2 + y^2z$                | 5           |



George Salmon (1819-1904)

# Other tensors

- Numerical computation of the Generic Rank
- Uniqueness of the CanD
- Tensors with particular symmetries
- Link with polynomials

# Numerical computation of the Generic Rank (1)

## Mapping

$$\{\mathbf{u}(\ell), \mathbf{v}(\ell), \dots, \mathbf{w}(\ell), 1 \leq \ell \leq r\} \xrightarrow{\varphi} \sum_{\ell=1}^r \mathbf{u}(\ell) \circ \mathbf{v}(\ell) \circ \dots \circ \mathbf{w}(\ell)$$

$$\{\mathbb{C}^{n_1} \circ \dots \circ \mathbb{C}^{n_d}\}^r \xrightarrow{\varphi} \mathcal{A}$$

➡ The **smallest**  $r$  for which  $\text{rank}(\text{Jacobian}(\varphi)) = D$  is the generic rank,  $\bar{R}$ .

# Numerical computation of the Generic Rank (2)

## Example 31: 3rd order non symmetric tensors

$$\{\mathbf{a}(\ell), \mathbf{b}(\ell), \mathbf{c}(\ell)\} \xrightarrow{\varphi} \mathbf{T} = \sum_{\ell=1}^r \mathbf{a}(\ell) \circ \mathbf{b}(\ell) \circ \mathbf{c}(\ell)$$

$\mathbf{T}$  has coordinate vector:  $\sum_{\ell=1}^r \mathbf{a}(\ell) \otimes \mathbf{b}(\ell) \otimes \mathbf{c}(\ell)$ . Hence the Jacobian of  $\varphi$  is the  $r(n_1 + n_2 + n_3) \times n_1 n_2 n_3$  matrix:

$$\mathbf{J} = \begin{bmatrix} \mathbf{I}_{n_1} & \otimes & \mathbf{b}^T(1) & \otimes & \mathbf{c}^T(1) \\ \mathbf{I}_{n_1} & \otimes & \dots & \otimes & \dots \\ \mathbf{I}_{n_1} & \otimes & \mathbf{b}^T(r) & \otimes & \mathbf{c}^T(r) \\ \mathbf{a}(1)^T & \otimes & \mathbf{I}_{n_2} & \otimes & \mathbf{c}^T(1) \\ \dots & \otimes & \mathbf{I}_{n_2} & \otimes & \dots \\ \mathbf{a}(r)^T & \otimes & \mathbf{I}_{n_2} & \otimes & \mathbf{c}^T(r) \\ \mathbf{a}(1)^T & \otimes & \mathbf{b}(1)^T & \otimes & \mathbf{I}_{n_3} \\ \dots & \otimes & \dots & \otimes & \mathbf{I}_{n_3} \\ \mathbf{a}(r)^T & \otimes & \mathbf{b}(r)^T & \otimes & \mathbf{I}_{n_3} \end{bmatrix} \quad \text{and} \quad \begin{cases} \text{rank}\{\mathbf{J}\} = \dim(\text{Im}\{\varphi\}) \\ \bar{R} = \text{Min}\{r : \text{Im}\{\varphi\} \end{cases}$$

# Numerical computation of the Generic Rank (3)

**Example 32:** Tensors of order  $d$  with dimensions all equal to  $K$

| $d \quad K$ | 2 | 3        | 4  | 5  | 6  | 7  |
|-------------|---|----------|----|----|----|----|
| 3           | 2 | <b>5</b> | 7  | 10 | 14 | 19 |
| 4           | 4 | 9        | 20 | 37 | 62 | 97 |

$$\bar{R} \geq \frac{K^d}{Kd - d + 1}$$

**Bold:** exceptions to the ceil rule  $\bar{R} = \lceil \frac{K^d}{Kd - d + 1} \rceil$

# Uniqueness of CanD

Number of solutions

- **Example 33: 3rd order with dimensions  $n_\ell$**

$$F(n_1, n_2, n_3) = (n_1 + n_2 + n_3 - 2) \bar{R} - n_1 n_2 n_3$$

- **Example 34:  $d$ th order with equal dimensions,  $K$**

$$F(n) = (Kd - d + 1) \bar{R} - K^d$$

| $d \quad K$ | 2 | 3        | 4 | 5 | 6 | 7  |
|-------------|---|----------|---|---|---|----|
| 3           | 0 | <b>8</b> | 6 | 5 | 8 | 18 |
| 4           | 4 | 0        | 4 | 4 | 6 | 24 |

➡ For generic/typical values, almost always infinitely many CanD's

# Numerical computation of the Generic Rank (4)

**Example 35: 3rd order tensors with unequal dim.  $N_\ell$  [CtB08] [CtB06]**

| $N_3$ |   | 2   |     |     |     | 3   |          |          | 4         |           |
|-------|---|-----|-----|-----|-----|-----|----------|----------|-----------|-----------|
| $N_2$ |   | 2   | 3   | 4   | 5   | 3   | 4        | 5        | 4         | 5         |
| $N_1$ | 2 | 2,3 | 3   | 4   | 4   | 3,4 | 4        | 5        | 4,5       | 5         |
|       | 3 | 3   | 3,4 | 4   | 5   | 5   | <b>5</b> | 5,6      | <b>6</b>  | <b>6</b>  |
|       | 4 | 4   | 4   | 4,5 | 5   | 5   | <b>6</b> | <b>6</b> | <b>7</b>  | <b>8</b>  |
|       | 5 | 4   | 5   | 5   | 5,6 | 5,6 | <b>6</b> | <b>8</b> | <b>8</b>  | <b>9</b>  |
|       | 6 | 4   | 6   | 6   | 6   | 6   | <b>7</b> | <b>8</b> | <b>8</b>  | <b>10</b> |
|       | 7 | 4   | 6   | 7   | 7   | 7   | <b>7</b> | <b>9</b> | <b>9</b>  | <b>10</b> |
|       | 8 | 4   | 6   | 8   | 8   | 8   | 8,9      | <b>9</b> | <b>10</b> | <b>11</b> |
|       | 9 | 4   | 6   | 8   | 9   | 9   | 9        | <b>9</b> | <b>10</b> | <b>12</b> |

- There are exceptions to the ceil rule  $\bar{R} = \lceil \frac{\prod_\ell N_\ell}{\sum_\ell (N_\ell - 1) + 1} \rceil$
- **Bold:** values that have not yet been proved theoretically

# Third order tensors with symmetric slices

**Example 36:** Typical ranks for  $N_1 \times N_2 \times N_2$  arrays, with  $N_2 \times N_2$  symmetric slices.

| $N_1$ | $N_2$ | 2   | 3   | 4        | 5         |
|-------|-------|-----|-----|----------|-----------|
| 2     |       | 2,3 | 3,4 | 4,5      | 5,6       |
| 3     |       | 3   | 4   | <b>6</b> | <b>7</b>  |
| 4     |       | 3   | 4,5 | <b>6</b> | <b>8</b>  |
| 5     |       | 3   | 5,6 | <b>7</b> | <b>9</b>  |
| 6     |       | 3   | 6   | <b>7</b> | <b>9</b>  |
| 7     |       | 3   | 6   | <b>7</b> | <b>10</b> |
| 8     |       | 3   | 6   | <b>8</b> | <b>10</b> |
| 9     |       | 3   | 6   | 9,10     | <b>11</b> |
| 10    |       | 3   | 6   | 10       | <b>11</b> |

**Bold:** smallest typical ranks computed numerically.

Plain: known typical ranks; in  $\mathbb{C}$ , the smallest value is generic.



# Tucker 3

- Definitions
- Properties
- Usefulness

# Definition (1)

According to Ledyard R. Tucker (1910-2004), any  $d$ th order  $I_1 \times I_2 \times \cdots \times I_d$  tensor  $\mathbf{T}$  can be decomposed as [TUC66]:

$$\mathbf{T} = \mathbf{S} \underset{1}{\bullet} \mathbf{U}^{(1)} \underset{1}{\bullet} \mathbf{U}^{(2)} \cdots \underset{1}{\bullet} \mathbf{U}^{(d)}$$

where  $\mathbf{S}$  has *smaller dimensions* than  $\mathbf{T}$  (or equal to), and  $\mathbf{U}^{(\ell)}$  are semi-unitary, i.e.  $\mathbf{U}^{(\ell)\top} \mathbf{U}^{(\ell)} = \mathbf{I}_{n_\ell}$ ,  $n_\ell \leq I_\ell$ .

➡  $\mathbf{S}$  is called the *core tensor*.

➡ This decomposition is referred to as *Tucker3* or as *HOSVD* [dLdMV00b] [SBG04].

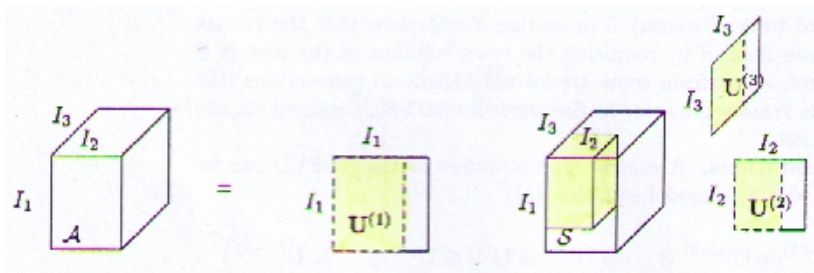
# Two equivalent optimization problems

$$\text{Max}_{\mathbf{U}^{(1)}, \mathbf{U}^{(2)} \dots \mathbf{U}^{(d)}} \left\| \mathbf{T} \bullet_1 \mathbf{U}^{(1)\top} \bullet_1 \mathbf{U}^{(2)\top} \dots \bullet_1 \mathbf{U}^{(d)\top} \right\|^2$$

$$\text{Min}_{\mathbf{U}^{(1)}, \mathbf{U}^{(2)} \dots \mathbf{U}^{(d)}} \left\| \mathbf{T} - \mathbf{S} \bullet_1 \mathbf{U}^{(1)} \bullet_1 \mathbf{U}^{(2)} \dots \bullet_1 \mathbf{U}^{(d)} \right\|^2$$

## Definition (2)

**Other writing** with unitary matrices  $\mathbf{U}^{(\ell)}$ . In that case, the core tensor  $\mathbf{S}$  has same dimensions as  $\mathbf{T}$  but is padded with zeros [dLdMV00b]:



# Properties (1)

The  $d$ th order core tensor can be imposed to be quite particular

- All its  $d - 1$ st order subtensors obtained by fixing one index are *all orthogonal* (w.r.t. scalar product induced by Frobenius norm); there are  $d$  of them.
- Entries of the core tensor can be sorted in such a way that for every mode  $\ell$ :

$$\|\mathbf{S}_{i_\ell=1}\| \geq \|\mathbf{S}_{i_\ell=2}\| \geq \dots \|\mathbf{S}_{i_\ell=I_\ell}\|$$

- These norms may be viewed as  $\ell$ -mode singular values.
- When  $\mathbf{T}$  is a matrix, so is  $\mathbf{S}$ , and all-orthogonality can be satisfied only when  $\mathbf{S}$  is diagonal. The sequence of norms  $\sigma_i = \|\mathbf{S}_{:i}\| = \|\mathbf{S}_{i:}\|$  are then the singular values.

## Properties (2)

- $\ell$ -mode singular vectors can be computed as singular vectors of the  $\ell$ -mode unfolding matrix; hence an easy computation
- The  $\ell$ -mode singular values are uniquely defined
- When  $\ell$ -mode singular values are different, corresponding  $\ell$ -mode singular vectors are unique up to a unit-modulus scale factor
- For any fixed mode  $\ell$ , the sum of all mode- $\ell$  squared singular values yields  $\|\mathbf{T}\|^2$

# Usefulness

- The nesting of  $\ell$ -mode singular values & vectors allows to easily find the best approximate of a tensor of lower  $\ell$ -mode rank by truncation of the HOSVD [dLdMV00c].
- May be applied to noise reduction
- May reduce subsequent computational complexity (dimension reduction)
- May be used as a pre-processing before the CanD calculation

# Other decompositions

- Exact decompositions (if not truncated):
  - CanD
  - Tucker3 – HOSVD
- Approximate decompositions:
  - Diagonalization by orthogonal transform
  - Diagonalization by invertible transform



# Conclusions on Tensors

- Still open problems
- Efficient numerical algorithms lacking
- Several ways of extending SVD to tensors
- Very powerful, and numerous application areas

# Polynomial interpolation

**Alexander-Hirschowitz Theorem** [AH92] [AH95] Let  $\mathcal{L}(d, m)$  be the space of hypersurfaces of degree at most  $d$  in  $m$  variables. This space is of dimension  $D(m, d) \stackrel{\text{def}}{=} \binom{m+d}{d} - 1$ .

**THEOREM** Denote  $\{p_i\}$   $K$  given distinct points in the complex projective space  $\mathbb{P}^m$ . The dimension of the linear subspace of hypersurfaces of  $\mathcal{L}(d, m)$  having multiplicity at least 2 at every point  $p_i$  is:

$$D(m, d) - K(m + 1)$$

except for the following cases:

- $d = 2$  and  $2 \leq K \leq m$
- $d \geq 3$  and  $(m, d, K) \in \{(2, 4, 5), (3, 4, 9), (4, 1, 14), (4, 3, 7)\}$

In other words, there are a *finite number* of exceptions.

## Part IV

### Algorithms for static mixtures

# Contents of part IV

## Overview

- Introduction
- Algorithms based on pair sweeping (CoM1, CoM2)  
Link with tensor diagonalization
- Algorithms based on matrix slices (JADE, STOTD)
- Algorithms based on Deflation (FastICA, RobustICA, SAUD)
- Finite alphabet inputs (APF, MAP, ILSP...)

References

# What we have seen so far

- Cumulants can measure independence at a given order
- Cumulants form a (symmetric) tensor object
- Tensors may have a rank larger than dimensions, even generically
- We have well-founded optimization criteria. Some of them amount to *approximately* diagonalizing a tensor.

# Hypotheses

- Mixture is over-determined
- The rank of the signal cumulant tensor is equal (at most) to its dimension
- The mixture may be given by the CanD of the signal cumulant tensor
- Noise & measurement errors yield a measured cumulant tensor that has generic rank

# Performance measure

How to test performances of algorithms in computer simulations?

- Difficulty because of the  $\mathbf{A P}$  indeterminacy
- **Identification:** Gap between  $\mathbf{F H}$  and matrix of the form  $\mathbf{A P}$
- **Source extraction:** SINR (Signal to Interference plus Noise): needs exhaustive search for best  $\mathbf{A P}$

## Example of Gap

This gap does not need a combinatorial search, because it is  $\Lambda \mathbf{P}$ -invariant [COM94a]:

$$\begin{aligned} \varepsilon(\mathbf{A}, \hat{\mathbf{A}}) &= \sum_i \left| \sum_j |\mathbf{D}_{ij}| - 1 \right|^2 + \left| \sum_j |\mathbf{D}_{ij}|^2 - 1 \right| \\ &+ \sum_j \left| \sum_i |\mathbf{D}_{ij}| - 1 \right|^2 + \left| \sum_i |\mathbf{D}_{ij}|^2 - 1 \right| \end{aligned}$$

where  $\mathbf{D} = \mathbf{A}^{-1} \hat{\mathbf{A}}$

### Properties

- $\varepsilon\{\mathbf{A} \Lambda \mathbf{P}, \hat{\mathbf{A}}\} = \varepsilon\{\mathbf{A}, \hat{\mathbf{A}}\} = \varepsilon\{\mathbf{A}, \hat{\mathbf{A}} \Lambda^{-1} \mathbf{P}\}$
- $\varepsilon\{\mathbf{A}, \hat{\mathbf{A}}\} = 0 \Leftrightarrow \hat{\mathbf{A}} = \mathbf{A} \Lambda \mathbf{P}$



# Algorithms based on pair sweeping

- Block vs Adaptive
- Closed-form solutions in dimension 2, for various contrasts
- Sweeping of all pairs
- Complexity and convergence

# Numerical Algorithms

## What problem are they supposed to solve?

- Are we given a single block of data?
- Are we observing a sequence of blocks, or a series of samples?
- Must we update the solution at every block, or at every sample?

## What kind of algorithms?

- Gradient ascent: the simplest
- Gradient-based ascents (Newton, quasi-Newton, conjugate gradient..)
- Quasi-algebraic algorithms: try to avoid *local maxima*
- Algebraic algorithms: find all absolute maxima in *closed-form*

# Block vs Adaptive

- Increase power of DSP
- Limitations of time-recursive Adaptive Algorithms
  - Convergence time of optimization algorithm
  - Convergence time of moment estimators
  - Local extrema harder to handle
- Coherence time sometimes limited  
(e.g. GSM: 900MHz, 190km/h,  $T_c \approx 2ms \approx 300$  symbols)
- Well matched to block transmission (TDMA)
- Better exploitation of data  
(uniform weight, resistance to loss in synchro, time reversal)

## Solution of the 2-dimensional problem

- Assume data  $x$  have been standardized into  $\tilde{x}$ .
- Then one looks for an estimate  $z$  of the source vector  $s$  as:

$$z = Q \tilde{x}$$

where  $Q$  is unitary, and may be assumed of the form:

$$Q = \begin{pmatrix} \cos \beta & \sin \beta e^{j\varphi} \\ -\sin \beta e^{-j\varphi} & \cos \beta \end{pmatrix} = \frac{1}{\sqrt{1 + \theta\theta^*}} \begin{pmatrix} 1 & \theta \\ -\theta^* & 1 \end{pmatrix} \quad (31)$$

where  $\theta \stackrel{\text{def}}{=} \tan \beta e^{j\varphi}$  denotes the complex tangent, and  $\beta \in ]-\pi/2, \pi/2]$ .

# Invariance & Indeterminacy (1)

- There is a whole class of equivalent absolute maxima, which can be deduced from each other by trivial filtering
- In the  $2 \times 2$  real case, there are 8 equivalent absolute maxima, generated by two **P** **A** transformations:

$$\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \quad \text{and} \quad \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$$

- In the complex case, there are infinitely many, when  $\varphi \in \mathbb{R}$ .
- Expression (31) fixes this indeterminacy, so that only 2 solutions remain

## What is the problem in dimension 2 ?

- $\Upsilon_{\alpha,r}$  is a homogeneous trigonometric polynomial in  $(\cos \beta, \sin \beta)$  of *degree  $\alpha r$* .
- And we want a closed-form (algebraic) solution
- But only polynomials of a single variable of *degree at most 4* can generally be rooted algebraically
- **Our problem:** check out whether  $\Upsilon_{\alpha,r}$  could be transformed into a particular function that can be algebraically maximized

# Invariance & Indeterminacy (2)

- Remark that  $\mathbf{Q}[\theta]$  and  $\mathbf{Q}[-1/\theta^*]$  are  $\mathbf{P}\mathbf{\Lambda}$ -related:

$$\mathbf{Q}[-\frac{1}{\theta^*}] = \mathbf{Q}[\theta] \begin{pmatrix} 0 & -e^{j\varphi} \\ e^{-j\varphi} & 0 \end{pmatrix}$$

- Thus, rational function  $\Upsilon$  satisfies  $\Upsilon[-\frac{1}{\theta^*}] = \Upsilon[\theta]$ .
- Consequently if  $\theta_o$  is stationary point of  $\Upsilon$ , so is  $-1/\theta_o^* \Rightarrow$   
Stationary points are roots of a polynomial  $\omega(\xi)$  in  
 $\xi \stackrel{\text{def}}{=} \theta - 1/\theta^*$
- Idea of algorithm:
  - Compute coefficients of  $\omega$  from cumulants of  $\tilde{\mathbf{x}}$
  - Compute roots  $\xi_o$  of  $\omega$
  - Root  $\theta^2 - \xi_o \theta - 1$  in order to get  $(\theta_o, -1/\theta_o^*)$ .

# Solution for contrast $\Upsilon_{2,3}$ in $\mathbb{R}$ (1)

- Contrast  $\Upsilon_{2,3}$  is defined as:

$$\Upsilon_{2,3} = \text{Cum}\{z_1, z_1, z_1\}^2 + \text{Cum}\{z_2, z_2, z_2\}^2 \stackrel{\text{def}}{=} (\kappa_{111})^2 + (\kappa_{222})^2$$

- Yet, by *multilinearity* of cumulants:

$$\kappa_{iii} = \sum_{jkl} Q_{ij} Q_{ik} Q_{il} \gamma_{jkl}, \quad \gamma_{jkl} \stackrel{\text{def}}{=} \text{Cum}\{\tilde{x}_j, \tilde{x}_k, \tilde{x}_l\}$$

- Then  $\Upsilon_{2,3}$  is a degree-6 polynomial in  $(\cos \beta, \sin \beta)$ , or a rational function in the tangent  $\theta$ :

$$\psi_3(\theta) = \left(\theta + \frac{1}{\theta}\right)^{-3} \sum_{i=1}^3 a_i (\theta^i - (-\theta)^{-i})$$



## Solution for contrast $\Upsilon_{2,3}$ in $\mathbb{R}$ (2)

- Denote  $\xi = \theta - 1/\theta$ .

Because of the invariance under transformation  $\theta \rightarrow -1/\theta$ , stationary points are roots of a *very simple* polynomial:

$$\omega_3(\xi) = d_2 \xi^2 + d_1 \xi - 4 d_2$$

where  $d_1 = a_1/3 - a_3$ , and  $d_2 = a_2/6$   
and:

$$a_3 = \gamma_{111}^2 + \gamma_{222}^2,$$

$$a_2 = 6(\gamma_{122} \gamma_{222} - \gamma_{111} \gamma_{112}),$$

$$a_1 = 9(\gamma_{122}^2 + \gamma_{112}^2) + 6(\gamma_{112} \gamma_{222} + \gamma_{111} \gamma_{122})$$

- **Conclusion:** solution obtainable *algebraically* from estimates of cumulants  $\gamma_{jkl} \stackrel{\text{def}}{=} \text{Cum}\{\tilde{x}_j, \tilde{x}_k, \tilde{x}_l\}$  [COM94b].

## Another solution for contrast $\Upsilon_{2,3}$ in $\mathbb{R}$ (2)

$\Upsilon_{2,3} = \kappa_{111}^2 + \kappa_{222}^2$  can be proved to be a quadratic form  $\mathbf{u}^T \mathbf{B} \mathbf{u}$  where

$$\mathbf{u} \stackrel{\text{def}}{=} [\cos 2\beta, \sin 2\beta]^T \quad (32)$$

and

$$\mathbf{B} \stackrel{\text{def}}{=} \begin{pmatrix} a_1 & 3a_4/2 \\ 3a_4/2 & 9a_2/4 + 3a_3/2 + a_1/4 \end{pmatrix}$$

with [dLdMV01]:

$$a_1 = \gamma_{111}^2 + \gamma_{222}^2$$

$$a_2 = \gamma_{112}^2 + \gamma_{122}^2$$

$$a_3 = \gamma_{111} \gamma_{122} + \gamma_{112} \gamma_{222}$$

$$a_4 = \gamma_{122} \gamma_{222} - \gamma_{111} \gamma_{112}$$

## Solution for contrast $\Upsilon_{2,4}$ in $\mathbb{R}$

- Now take  $\Upsilon_{2,4} \stackrel{\text{def}}{=} (\kappa_{1111})^2 + (\kappa_{2222})^2$
- This contrast is a degree-8 polynomial  $(\cos \beta, \sin \beta)$ .  
Denote again  $\xi = \theta - 1/\theta$ . Then it is a rational function in  $\xi$ :

$$\psi_4(\xi) = (\xi^2 + 4)^{-2} \sum_{i=0}^4 b_i \xi^i$$

- Then its stationary points are roots of a polynomial of degree 4:

$$\omega_4(\xi) = \sum_{i=0}^4 c_i \xi^i$$

whose roots are thus obtainable *algebraically*  
(e.g. via *Ferrari*'s technique).

- Coefficients  $b_i$  and  $c_i$  are given in [COM94b] as functions of  $\gamma_{ijkl}$

## Solution for contrast $\Upsilon_{1,4}$ in $\mathbb{R}$

Same approach feasible, but easier because absence of squares  
 $\Rightarrow$  Here another easier-accessible approach

### ■ Input-Output relations

$$\begin{aligned}\kappa_1 &= \gamma_1 \cos^4 \beta + 4\gamma_{1112} \cos^3 \beta \sin \beta + 6\gamma_{1122} \cos^2 \beta \sin^2 \beta \\ &\quad + 4\gamma_{1222} \cos \beta \sin^3 \beta + \gamma_2 \sin^4 \beta \\ \kappa_2 &= \gamma_1 \sin^4 \beta - 4\gamma_{1112} \cos \beta \sin^3 \beta + 6\gamma_{1122} \cos^2 \beta \sin^2 \beta \\ &\quad - 4\gamma_{1222} \cos^3 \beta \sin \beta + \gamma_2 \cos^4 \beta\end{aligned}$$

### ■ Then $\varepsilon \Upsilon_{1,4} = \kappa_1 + \kappa_2 =$

$$[\cos 2\beta \quad \sin 2\beta] \begin{pmatrix} \gamma_1 + \gamma_2 & \gamma_{1112} - \gamma_{1222} \\ \gamma_{1112} - \gamma_{1222} & \frac{\gamma_1 + \gamma_2}{2} + 3\gamma_{1122} \end{pmatrix} \begin{bmatrix} \cos 2\beta \\ \sin 2\beta \end{bmatrix}$$

### ■ Conclusion: again entirely *algebraic* since dominant eigenvector of a matrix of size $< 4$ .

# Solution for contrast $\Upsilon_{1,4}$ in $\mathbb{C}$

- Define  $\kappa_i = \text{Cum}\{z_i, z_i, z_i^*, z_i^*\}$ ,  $\gamma_{ij}^{k\ell} = \text{Cum}\{\tilde{x}_i, \tilde{x}_j, \tilde{x}_k^*, \tilde{x}_\ell^*\}$
- **Then...** again a quadratic form

$$\varepsilon \Upsilon_{1,4} = \kappa_1 + \kappa_2 = \mathbf{u}^T \mathbf{B} \mathbf{u}$$

with

$$\mathbf{u}^T = [\cos 2\beta \quad \sin 2\beta \cos \varphi \quad \sin 2\beta \sin \varphi]$$

and

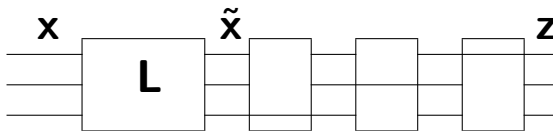
$$\mathbf{B} = \begin{pmatrix} \gamma_{1111} + \gamma_{2222} & \Re\{\delta\} & -\Im\{\delta\} \\ \Re\{\delta\} & 2\gamma_{12}^{12} + \Re\{\gamma_{22}^{11}\} & \Im\{\gamma_{22}^{11}\} \\ -\Im\{\delta\} & \Im\{\gamma_{22}^{11}\} & 2\gamma_{12}^{12} - \Re\{\gamma_{22}^{11}\} \end{pmatrix};$$

$$\delta = \gamma_{12}^{11} - \gamma_{22}^{12}$$

**Conclusion:** unexpectedly *entirely algebraic*! [COM01]

# Jacobi Sweeping

Cyclic sweeping with fixed ordering: Example in dimension  $P = 3$



Carl Jacobi, 1804-1851

# Jacobi Sweeping for tensors

**Question:** Why not select another ordering, e.g. process pairs having cross cumulants of largest magnitude?

**Response:** the computational complexity would be dominated by the computation of the tensor entries themselves!

How do we compute tensor entries then?

## Jacobi Sweeping for tensors

Sweeping a  $3 \times 3 \times 3$  tensor [COM89]

$$\begin{pmatrix} \textcolor{red}{X} & x & x \\ x & x & x \\ x & x & . \end{pmatrix} \begin{pmatrix} x & x & x \\ x & \textcolor{red}{X} & x \\ x & x & . \end{pmatrix} \begin{pmatrix} x & x & x \\ x & x & x \\ x & x & . \end{pmatrix} \rightarrow \begin{pmatrix} \textcolor{red}{X} & x & x \\ x & . & x \\ x & x & x \end{pmatrix} \begin{pmatrix} x & x & x \\ x & . & x \\ x & x & x \end{pmatrix} \begin{pmatrix} x & x & x \\ x & . & x \\ x & x & \textcolor{red}{X} \end{pmatrix} \rightarrow \begin{pmatrix} . & x & x \\ x & x & x \\ x & x & x \end{pmatrix} \begin{pmatrix} . & x & x \\ x & \textcolor{red}{X} & x \\ x & x & x \end{pmatrix} \begin{pmatrix} . & x & x \\ x & x & x \\ x & x & \textcolor{red}{X} \end{pmatrix}$$



$\textcolor{red}{X}$  : maximized  
 $x$  : minimized  
 $.$  : unchanged

} by last Givens rotation



## Two possible updates of $\mathbf{T}$

After processing *every* pair, one can:

- Update based on *multilinearity*:

$$T_{ij..k} \leftarrow \sum_{pq..r} Q_{ip} Q_{jq} .. Q_{kr} T_{pq..r}$$

requires an initial computation of  $\mathbf{T}$

- Update of *observations* themselves

$$\mathbf{X} \leftarrow \mathbf{Q} \mathbf{X}$$

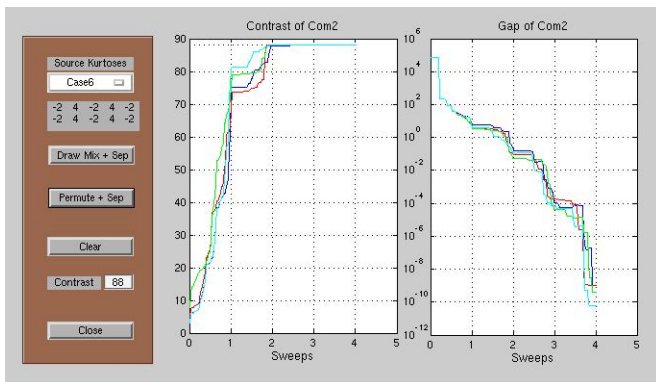
and then

$$T_{ij..k} = \text{Cum}\{x_i, x_j, ..x_k\}$$

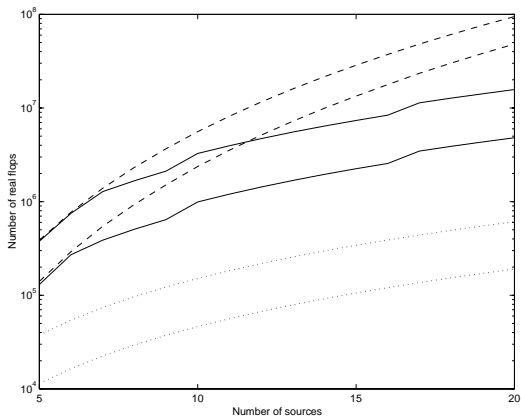
The best choice (i.e. least costly) depends on data length and dimensions.

# Influence of ordering

With update based on multilinearity.



# Complexity



# Interpretation in terms of pairwise independence

- Pairs are processed in turns, so as to make outputs as independent as possible
- Ultimately: a set of *pairwise independent* outputs
- Legitimate because of corollary of Darmois's theorem (cf., slide 38)

# Interpretation in terms of tensor diagonalization

## Explanation for order 3 tensors

- Given a tensor  $g_{ijk}$ , find a matrix  $\mathbf{Q}$  transforming  $g$  into  $G_{pqr} = \sum_{ijk} Q_{pi} Q_{qj} Q_{rk} g_{ijk}$  such as to maximize:

$$\Psi_3(\mathbf{Q}) \stackrel{\text{def}}{=} \sum_i |G_{iii}|^2$$

- *Theorem:* if  $\mathbf{Q}$  is unitary, then  $\Omega \stackrel{\text{def}}{=} \sum_{ijk} |G_{ijk}|^2$  is constant independent of  $\mathbf{Q}$

*Proof:* uses  $\sum_p Q_{ip} Q_{jp} = \delta_{ij}$

- *Corollary:* Maximize  $\Upsilon_{3,2} \Leftrightarrow$  minimize all non diagonal entries  
Hence: “Tensor Diagonalization”

# Tensor diagonalization

**Warning:** Tensors cannot in general be diagonalized by congruent transforms, even non unitary!

**Why?**

...

# Stationary points

## Example of diagonalization of real symmetric matrices

- Given a matrix  $g$  with components  $g_{ij}$ , it is sought for an orthogonal matrix  $Q$  such that  $\psi_2$  is maximized:

$$\psi_2(G) = \sum_i G_{ii}^2; \quad G_{ij} = \sum_{p,q} Q_{ip} Q_{jq} g_{pq}.$$

- Stationary points of  $\psi_2$  satisfy for any pair of indices  $(q, r), q \neq r$ :

$$G_{qq} G_{qr} = G_{rr} G_{qr}$$

- Next,  $d^2\psi_2 < 0 \Leftrightarrow G_{qr}^2 < (G_{qq} - G_{rr})^2$ , which proves that
  - $G_{qr} = 0, \forall q \neq r$  yields a maximum
  - $G_{qq} = G_{rr}, \forall q, r$  yields a minimum
  - Other stationary points are saddle points

# Stationary points

## Procedure applied to real 3rd or 4th order tensors

- Similarly, one can look at relations characterizing local maxima of criteria  $\Psi_3$  and  $\Psi_4$  [COM94b]:

$$\begin{aligned}
 G_{qqq}G_{qqr} - G_{rrr}G_{qrr} &= 0, \\
 4G_{qqr}^2 + 4G_{qrr}^2 - (G_{qqq} - G_{qrr})^2 - (G_{rrr} - G_{qqr})^2 &< 0; \\
 G_{qqqq}G_{qqqr} - G_{rrrr}G_{qrrr} &= 0, \\
 4G_{qqqr}^2 + 4G_{qrrr}^2 - (G_{qqqq} - \frac{3}{2}G_{qqqr})^2 \\
 - (G_{rrrr} - \frac{3}{2}G_{qrrr})^2 &< 0.
 \end{aligned}$$

for any pair of indices  $(p, q), p \neq q$ . As a conclusion, contrary to  $\Psi_2$  in the matrix case,  $\Psi_r$  might have theoretically spurious local maxima in the tensor case,  $r > 2$



# Algorithms based on matrix slices

- JADE contrast
- JADE algorithm
- STOTD recursion on the order
- Other

# Tensors as Linear Operators

## Overview

- Linear Operator  $\Omega$  acting on square matrices:

$$\mathbf{M} \longrightarrow \Omega(\mathbf{M})_{ij} = \sum_{k\ell} \mathcal{C}_{ik}^{j\ell} M_{k\ell}$$

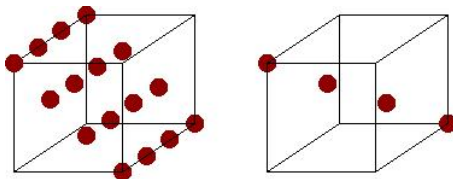
admits eigen-matrices  $\mathbf{N}(p)$ ,  $1 \leq p \leq P^2$ .

- In the absence of noise,  $P$  nonzero eigenvalues
- In practice, retain  $P$  dominant eigen-matrices  $\Rightarrow$  (i) reduced complexity  $P^2$ , and (ii) noise reduction

# Joint Approximate Diagonalization (JAD)

## Back to tensor diagonalization

Example of  $4 \times 4 \times 4$  tensors



Matrix slices diagonalization  $\neq$  Tensor diagonalization

# Real symmetric tensors

**Definition** (reminder)

$\mathbf{G}$  is real symmetric iff:

$$G_{ij..k} = G_{\sigma(ij..k)}$$

for all permutation  $\sigma$

## Two equivalent writings (order 3)

**Lemma 1** Let  $\mathbf{U}$  be an orthogonal real matrix, relating two 3rd order real symmetric tensors  $\mathbf{G}$  and  $\mathbf{g}$ , then

$$\sum_{ik} G_{iik}^2 = \sum_r \|\mathbf{Diag}(\mathbf{U}^T \mathbf{M}(r) \mathbf{U})\|^2$$

where  $\mathbf{M}(r)$  are symmetric matrix slices of  $\mathbf{g}$ :  $M_{pq}(r) \stackrel{\text{def}}{=} g_{pqr}$

*Proof...*

**Theorem** One can prove that  $\mathcal{J} \stackrel{\text{def}}{=} \sum_{ik} |G_{iik}|^2$  is a *contrast*. (at least 2 indices are equal)

# Hermitian tensors

## Definition

**G** is complex hermitian iff it is of even order, and enjoys the symmetries:

- $G_{ij..k}^{pq..r} = G_{\sigma(ij..k)}^{pq..r}$
- $G_{ij..k}^{pq..r} = G_{ij..k}^{\sigma(pq..r)}$
- $G_{ij..k}^{pq..r} = \left( G_{pq..r}^{ij..k} \right)^*$

for any permutation  $\sigma$ .

## Two equivalent writings (order 4)

**Lemma 2** Let  $\mathbf{U}$  be a unitary matrix relating two complex hermitian tensors of even order 4,  $\mathbf{G}$  and  $\mathbf{g}$ , then

$$\sum_{ik\ell} |G_{i\ell}^{ik}|^2 = \sum_{rs} \|\mathbf{Diag}(\mathbf{U}^H \mathbf{M}(r, s) \mathbf{U})\|^2$$

where  $\mathbf{M}(r, s)$  are hermitian matrix slices of  $\mathbf{g}$ :  $M_{pq}(r, s) \stackrel{\text{def}}{=} g_{ps}^{qr}$

*Proof...*

**Theorem** One can prove that  $\mathcal{J} \stackrel{\text{def}}{=} \sum_{ik\ell..mn} |G_{i\ell..n}^{ik..m}|^2$  is a *contrast*. (only 2 indices are equal)

## JADE as an approximation of $\Upsilon_{\alpha,4}$

**Lemma 3** denote the EVD  $\mathbf{g} = \sum_p \lambda_p \mathbf{N}(p) \mathbf{N}(p)^H$ , i.e.  $g_{jkr} = \sum_p \lambda_p N_{jk}(p) N_{rs}(p)$ , then 3rd writing:

$$\mathcal{J}_{2,4} = \sum_{p=1}^{P^2} \lambda_p^2 \|\mathbf{diag}(\mathbf{U}^H \mathbf{N}(p) \mathbf{U})\|^2$$

**Second approximation:** Keep only the most significant eigen-matrices,  $p \leq P$ , which amounts to maximizing:

$$\mathcal{J}_{\alpha,4}^E \stackrel{\text{def}}{=} \sum_{p=1}^P \lambda_p^\alpha \|\mathbf{diag}(\mathbf{U}^H \mathbf{N}(p) \mathbf{U})\|^2$$

- Hence the name of Joint Approximate Diagonalization of Eigenmatrices (JADE).
- $\mathcal{J}_{\alpha,4}$  can be seen as an *approximation* of  $\Upsilon_{\alpha,4}$ .



# Implementation of JADE with pair sweeping

## Algebraic solution in dim 2

- Goal is to maximize the diagonal terms of  $\mathbf{Q}^H \mathbf{N}(r) \mathbf{Q}$

- Denote  $\mathbf{N}(r) = \begin{pmatrix} a_r & b_r \\ c_r & d_r \end{pmatrix}$  and

$$\mathbf{Q} = \begin{pmatrix} \cos \theta & \sin \theta e^{j\varphi} \\ -\sin \theta e^{-j\varphi} & \cos \theta \end{pmatrix}$$

- Then this amounts to maximizing w.r.t.  $(\theta, \varphi)$ :  $\mathbf{v}^T \Re(\mathbf{G}^H \mathbf{G}) \mathbf{v}$  where

$$\mathbf{G}^H \mathbf{G} = \sum_r \begin{bmatrix} a_r - d_r \\ b_r + c_r \\ j(c_r - b_r) \end{bmatrix}^* [a_r - d_r, b_r + c_r, j(c_r - b_r)]$$

and  $\mathbf{v} = [\cos 2\theta, \sin 2\theta \cos \varphi, \sin 2\theta \sin \varphi]^T$

- Thus, solution is the dominant eigenvector of a (real) symmetric matrix

# Lower order simultaneous diagonalization (1)

## Extend the idea: Slicing decreases the order

- Similarly, one can try to diagonalize a 4th order tensor  $\mathbf{T} = [\gamma_{ijkl}]$  by jointly diagonalizing 3rd order slices  $\mathbf{T}(\ell)$  (STOTD) [dLdMV01]
- Algorithm: Each Givens rotation is obtained again by maximizing a quadratic form  $\mathbf{u}^T \mathbf{B} \mathbf{u}$
- Noise reduction possibility: replace slices by a family of 3rd order tensors forming a basis of the map  $\mathbb{C}^K \rightarrow \mathbb{C}^{K \times K \times K}$  (consider the 4th order tensor as a linear map; basis obtained by SVD)

## Lower order simultaneous diagonalization (2)

In the real case,  $\mathbf{B}$  is given as in slide 186 by:

$$\mathbf{B} = \begin{pmatrix} a_1 & 3 a_4/2 \\ 3 a_4/2 & 9 a_2/4 + 3 a_3/2 + a_1/4 \end{pmatrix}$$

with [dLdMV01]:

$$a_1 = \sum_{\ell} \gamma_{111\ell}^2 + \gamma_{222\ell}^2$$

$$a_2 = \sum_{\ell} \gamma_{112\ell}^2 + \gamma_{122\ell}^2$$

$$a_3 = \sum_{\ell} \gamma_{111\ell} \gamma_{122\ell} + \gamma_{112\ell} \gamma_{222\ell}$$

$$a_4 = \sum_{\ell} \gamma_{122\ell} \gamma_{222\ell} - \gamma_{111\ell} \gamma_{112\ell}$$

# Diagonalization algorithms

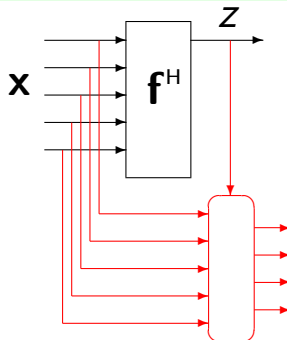
Obtain a diagonal tensor or diagonal slices:

- by orthogonal transforms [dL78] [CS93] [Com92]
- by invertible transforms [AFS07] [YER02] [?] [PHA01] [LAT06]
- by rectangular transforms [PAA99] [VO06] [COM04a] [NL06]

# Algorithms based on Deflation

- Principle: Joint extraction vs Deflation
- Unitary adaptive deflation
- A so-called fixed point: FastICA
- RobustICA
- Deflation without spatial prewhitening, algebraic deflation
- Discussion on MISO criteria

# Joint extraction vs Deflation



## Deflation:

- Advantage: (a) reduced complexity at each stage, (b) simpler to understand
- Drawbacks: (i) accumulation of regression errors, limitation of number of extracted sources, (ii) possibly larger final complexity

# Adaptive algorithms

Deflation by Kurtosis Gradient Ascent

## Again same idea

After standardization, it is equivalent to maximize 4th order moment criterion,  $\mathcal{M}_z(4) = \mathbb{E}\{|z|^4\}$ , whose gradient is:

$$\nabla \mathcal{M} = 4 \mathbb{E}\{\mathbf{x} (\mathbf{f}^H \mathbf{x})(\mathbf{x}^H \mathbf{f})^2\}$$

## Overview

- Fixed step gradient on angular parameters: [DL95]
- Locally optimal step gradient on filter taps: FastICA [HYV97]
- Globally optimal step gradient on filter taps: RobustICA [COM02a]
- Semi-Algebraic Unitary Deflation (SAUD) [COM05]

# Adaptive algorithms

## Adaptive implementation

- Fully adaptive solutions (update at every sample arrival)  
nowadays little useful
- Always easy to devise fully adaptive, or block-adaptive  
solutions from block semi-algebraic algorithms (but  
reverse is not true!)



# Unitary adaptive deflation (1)

## ■ Extraction

- To extract the first source, find a unitary matrix  $\mathbf{U}$  so as to maximize the kurtosis of the first output
- Matrix  $\mathbf{U}$  can be iteratively determined by a sequence of Givens rotations
- At each step, determine the best angle of the Givens rotation, e.g. by a gradient ascent [DL95]

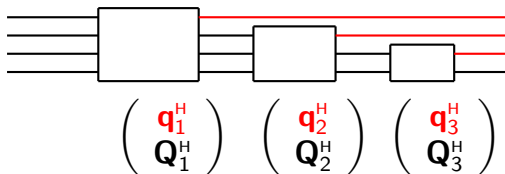
NB: only  $P - 1$  Givens rotations are involved

## ■ Deflation

- After convergence, the first output is extracted, and the  $P - 1$  remaining outputs of  $\mathbf{U}$  can be processed in the same way

## Unitary adaptive deflation (2)

At stage  $k$ ,  $\mathbf{Q} = \begin{pmatrix} \mathbf{q}_k^H \\ \mathbf{Q}_k^H \end{pmatrix}$  is unitary of size  $P - k + 1$ , and only its first row is used to extract source  $k$ ,  $1 \leq k \leq P - 1$



# A so-called fixed point: FastICA (1)

- Any gradient ascent of a function  $\mathcal{M}_\rho = \mathbb{E}\{\rho(\mathbf{f}^H \mathbf{x})\}$  under unit-norm constraint  $\|\mathbf{f}\|^2 = 1$  admits the Lagrangian formulation:

$$\mathbb{E}\{\mathbf{x} \dot{\rho}(\mathbf{f}^H \mathbf{x})\} = \lambda \mathbf{f}$$

- Convergence:** when  $\nabla \mathcal{C}$  and  $\mathbf{f}$  collinear (and *not* when gradient is null, because of constraint  $\|\mathbf{f}\| = 1$ ).
- Remark:** It is *not* a fixed point algorithm, contrary to what had been claimed in [HYV97], because  $\lambda$  is not known!
- One can take  $\rho(z) = |z|^4$

## A so-called fixed point: FastICA (2)

Details of the algorithm proposed in [HYV99] in the real field; only difference compared to [TUG97] is fixed step size.

- **Gradient:**  $\nabla \mathcal{M} = 4 \mathbb{E}\{\mathbf{x} (\mathbf{f}^T \mathbf{x})^3\}$
- **Hessian:**  $12 \mathbb{E}\{\mathbf{x} \mathbf{x}^T (\mathbf{f}^T \mathbf{x})^2\}$
- **Heavy approximation** of Hyvarinen [HYV99]:

$$\mathbb{E}\{\mathbf{x} \mathbf{x}^T (\mathbf{f}^T \mathbf{x})^2\} \approx \mathbb{E}\{\mathbf{x} \mathbf{x}^T\} \mathbb{E}\{(\mathbf{f}^T \mathbf{x})^2\}$$

- If  $\mathbf{x}$  standardized and  $\mathbf{f}$  unit norm, then Hessian equals Identity.
- This yields an approximate Newton iteration: *a mere fixed step gradient!*

$$\begin{aligned} \mathbf{f} &\leftarrow \mathbf{f} - \frac{1}{3} \mathbb{E}\{\mathbf{x} (\mathbf{f}^T \mathbf{x})^3\} & \text{or} & & \mathbf{f} &\leftarrow \mathbb{E}\{\mathbf{x} (\mathbf{f}^T \mathbf{x})^3\} - 3 \mathbf{f} \\ \mathbf{f} &\leftarrow \mathbf{f} / \|\mathbf{f}\| \end{aligned}$$

## FastICA: weaknesses

This is a mere fixed step-size projected gradient algorithm, inheriting problems such as:

- Saddle points (slow/ill convergence)
- Flat areas (slow convergence)
- Local maxima (ill convergence)

NB: slow convergence may mean high complexity to reach the solution, or stopping iterations before reaching convergence (depends on stopping criterion).

# Polynomial rooting

**Theorem (1830).** A polynomial of degree higher than 4 cannot in general be rooted algebraically in terms of a finite number of additions, subtractions, multiplications, divisions, and radicals (root extractions).



**Niels Abel, 1802-1829**



**Evariste Galois 1811-1832**

# How to fix most drawbacks: RobustICA

**Principle:** Cheap exhaustive Line Search of a criterion  $\mathcal{J}$

- Look for *absolute maximum* in the gradient direction (1-dim search)
- *Not costly* when criteria are polynomials or rational functions of low degree (same as AMiSRoF: polynomial to root, but here *at most of degree 4*)
- Applies to Kurtosis Maximization (KMA), Constant-Modulus (CMA), Constant-Power (CPA) Algorithms...

This yields corresponding Optimal-Step (OS) algorithms:  
OS-KMA, OS-CMA, OS-CPA... [ZC08] [ZC05]

# RobustICA

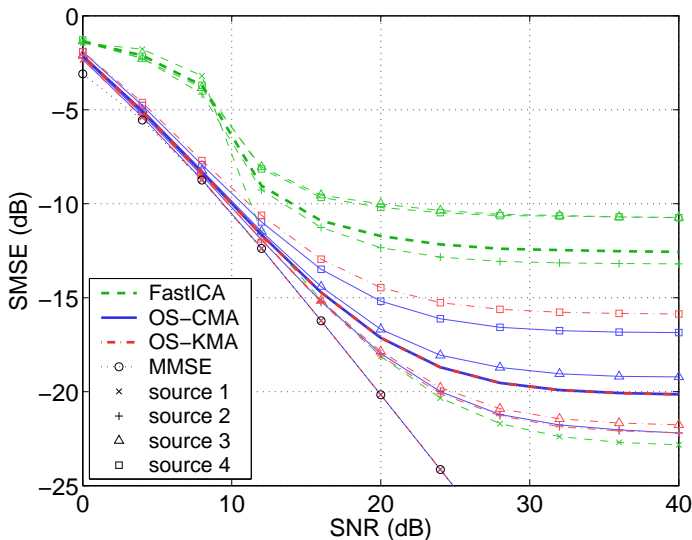
## Algorithm

- compute coefficients of polynomial  $\frac{\partial}{\partial \mu} \mathcal{J}(\mathbf{f} + \mu \nabla)$  for fixed  $\mathbf{f}$  and  $\nabla$
- compute all its roots  $\{\mu_i\}$
- select  $\mu_{opt}$  among those roots, which yields the absolute maximum
- set  $\mathbf{f} \leftarrow \mathbf{f} + \mu_{opt} \nabla$

[ZC07]



## RobustICA vs FastICA



# Semi-Algebraic Unitary Deflation

## CoM1 [COM01]

```

Loop on sweeps
  for  $i = 1$  to  $P - 1$ 
    for  $j = i$  to  $P$ 
      Algebraic  $2 \times 2$  separ.
    end
  end
end
Extraction

```

## SAUD [ACX07]

```

for  $i = 1$  to  $P - 1$ 
  Loop on sweeps
    for  $j = i$  to  $P$ 
      Algebraic  $2 \times 2$  separ.
    end
  end
  Extraction
end

```

# Equivalence between KMA and CMA

- Recall the 2 criteria:

$$\Upsilon_{KMA} = \frac{\text{Cum}\{z, z, z^*, z^*\}}{[\text{E}\{|z|^2\}]^2}, \quad \mathcal{J}_{CMA} = \text{E}\{[|z|^2 - R]^2\}$$

- Assume 2nd Order circular sources:  $\text{E}\{s^2\} = 0$
- Then KMA and CMA are equivalent

Proof.

# Discussion on Deflation (MISO) criteria

Let  $z \stackrel{\text{def}}{=} \mathbf{f}^H \mathbf{x}$ . Criteria below stationary iff differentials  $\dot{\mathbf{p}}$  and  $\dot{\mathbf{q}}$  are collinear:

■ **Ratio:**  $\text{Max}_{\mathbf{f}} \frac{p(\mathbf{f})}{q(\mathbf{f})}$

*Example: Kurtosis*, with  $p = E\{|z|^4\} - 2E\{|z|^2\}^2 - |E\{z^2\}|^2$  and  $q = E\{|z|^2\}^2$

■ **Difference:**  $\text{Min}_{\mathbf{f}} p(\mathbf{f}) - \alpha q(\mathbf{f})$

*Example: Constant Modulus*, with  $p = E\{|z|^4\}$  and  $q = 2a E\{|z|^2\} - a^2$  or **Constant Power**, with  $q = 2a \Re(E\{z^2\}) - a^2$

■ **Constrained:**  $\text{Max}_{q(\mathbf{f})=1} p(\mathbf{f})$

*Example: Cumulant*, with

$$p = E\{|z|^4\} - 2E\{|z|^2\}^2 - |E\{z^2\}|^2$$

*Example: Moment*, with  $p = E\{|z|^4\}$ , if standardized and with either  $q = \|\mathbf{f}\|^2$  or  $q = E\{|z|^2\}^2$

# Finite alphabets

- Back to contrast criteria: APF
- Approximation of the MAP estimate
- Semi-Algebraic Blind Extraction algorithm: AMiSRoF
- Blind Extraction by ILSP
- Convolutional model
- Presence of Carrier Offset (in Digital Communications)

## Contrast for discrete inputs (1)

- **Hypothesis H5** The sources take their value in a finite alphabet  $\mathcal{A}$  defined by the roots in  $\mathbb{C}$  of some polynomial  $q(z) = 0$
- **Theorem** [COM04b]  
Under **H5**, the following is a contrast over the set  $\mathcal{H}$  of invertible  $P \times P$  FIR filters.

$$\Upsilon(\mathbf{G}; \mathbf{z}) \stackrel{\text{def}}{=} - \sum_n \sum_i |q(z_i[n])|^2 \quad (33)$$

**APF:** Algebraic Polynomial Fitting

## Contrast for discrete inputs (2)

- For given alphabet  $\mathcal{A}$ , denote  $\mathcal{G}$  the set of numbers  $c$  such that  $c\mathcal{A} \subseteq \mathcal{A}$ .
- **Lemma 1** Trivial filters satisfying **H5** are of the form:

$$\mathbf{P} \mathbf{D}[z]$$

with  $\mathbf{D}[z]$  diagonal and  $D_{pp}[z] = c_p z^n$ , for some  $n \in \mathbb{Z}$  and some  $c_p \in \mathcal{G}$ .

- Because  $\mathcal{A}$  is finite, any  $c \in \mathcal{G}$  must be of unit modulus, and we must have  $c\mathcal{A} = \mathcal{A}, \forall c \in \mathcal{G}$ .  
Also any  $c \in \mathcal{G}$  has an inverse  $c^{-1}$  in  $\mathcal{G}$ .

# Contrast for discrete inputs (3)

Sketch of proof of the theorem. We prove the 3 properties of slide 68:

- $\forall \mathbf{T} \in \mathcal{T}, \Upsilon(\mathbf{T}; \mathbf{x}) = \Upsilon(\mathbf{I}; \mathbf{x})$
- $\forall \mathbf{G} \in \mathcal{H}, \forall \mathbf{s} \in \mathcal{S}$ , set of independent sources in  $\mathcal{A}$ ,  
 $\Upsilon(\mathbf{G}; \mathbf{s}) \leq \Upsilon(\mathbf{I}; \mathbf{s})$
- $\forall \mathbf{G} \in \mathcal{H}, \forall \mathbf{s} \in \mathcal{S}$ , equality  $\Upsilon(\mathbf{G}; \mathbf{s}) = \Upsilon(\mathbf{I}; \mathbf{s}) \Rightarrow \mathbf{G}$  trivial.

The proof needs the lemma

- **Lemma 2** Let  $\mathcal{A}$  be  $\{a_k, 1 \leq k \leq d\} \neq \{0\}$ . If  
 $\sum_{i=1}^L c_i a_{\sigma(i)} \in \mathcal{A}$  for all mappings  $\sigma$  from  $\{1, \dots, L\}$  to  $\{1, \dots, d\}$ , then only one  $c_i \neq 0$ .
- The proof of this lemma needs sources to be *sufficiently exciting*, e.g. that all binary states are present.



## Contrast for discrete inputs (4)

### Idea of the proof of Lemma 2

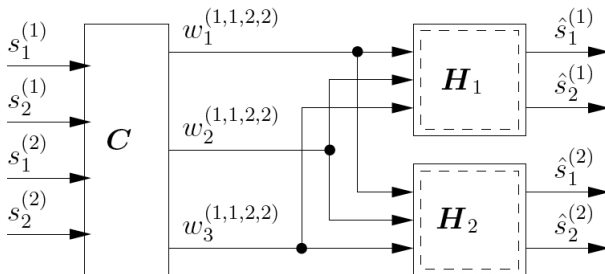
- Assume that for some  $\mathbf{c} \in \mathbb{C}^L$ , we have  $\mathbf{x}^T \mathbf{c} \in \mathcal{A}$  for all  $\mathbf{x} \in \mathcal{A}^L$ .
- Then  $\mathbf{c}$  must be trivial:  
Non trivial vectors  $\mathbf{c}$  may generate symbols that lie outside the convex hull of  $\mathcal{A}$ , or between the two closest symbols.

## Contrast for discrete inputs (4)

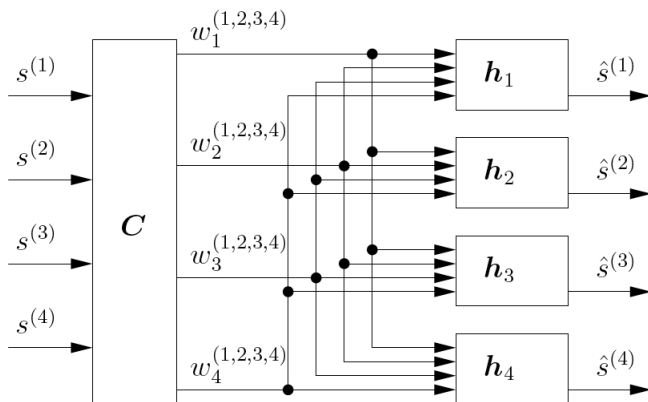
### Advantages

- The previous contrast allows to separate *correlated* sources
- But it needs all sources to have the same (known) alphabet
- If sources have *different* alphabets, one can extract sources in parallel with different criteria: *Parallel Extraction* [RZC05]
- By deflation with different criteria, one can extract more sources than sensors: *Parallel Deflation* [RZC05]

# Parallel Deflation

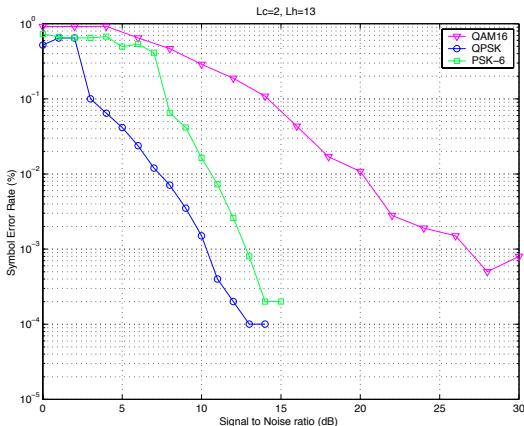


# Parallel Extraction



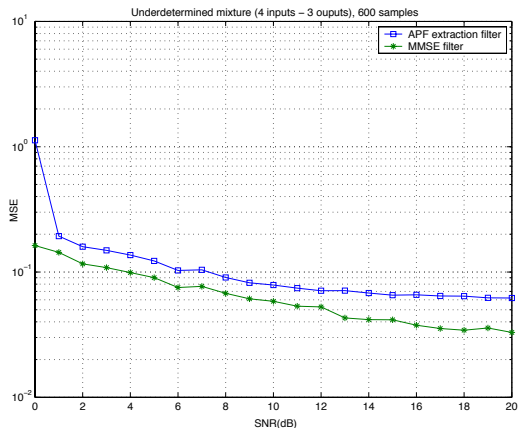
# Parallel extraction

Parallel extraction of 3 sources (QPSK, QAM16, PSK6), from a 3-sensor length-3 random Gaussian channel [RZC05]



# APF extraction

Parallel Deflation from a mixture of 4 sources (2 QPAK and 2 QAM16) received on 3 sensors. Extraction of a QPSK source in figure, compared to MMSE [RZC05]



# MAP estimate

- Optimal solution

$$(\hat{\mathbf{H}}, \hat{\mathbf{s}})_{MAP} = \underset{\mathbf{H}, \mathbf{s}}{\operatorname{Arg\,Max}} p_{s|x,H}(\mathbf{x}, \mathbf{s}, \mathbf{H})$$

- If  $s_p \in \mathcal{A}$ , and if noise is Gaussian, then

$$(\hat{\mathbf{H}}, \hat{\mathbf{s}})_{MAP} = \underset{\mathbf{H}, \mathbf{s} \in \mathcal{A}^P}{\operatorname{Arg\,Min}} \|\mathbf{x} - \mathbf{H} \mathbf{s}\|^2$$

- Less costly to search (inverse filter when it exists)

$$(\hat{\mathbf{F}}, \hat{\mathbf{s}})_{MAP} = \underset{\mathbf{F}, \mathbf{s} \in \mathcal{A}^P}{\operatorname{Arg\,Min}} \|\mathbf{F} \mathbf{x} - \mathbf{s}\|^2$$

- or by deflation:

$$(\hat{\mathbf{f}}, \hat{\mathbf{s}})_{MAP} = \underset{\mathbf{f}, \mathbf{s} \in \mathcal{A}^P}{\operatorname{Arg\,Min}} \|\mathbf{f}^H \mathbf{x} - \mathbf{s}\|^2 \quad (34)$$

# Approximation of the MAP estimate

For alphabet of constant modulus, MAP criterion (34) is asymptotically equivalent (for large samples of size  $T$ ) to [GC98]:

$$\Upsilon_T(\mathbf{f}) = \frac{1}{T} \sum_{t=1}^T \prod_{j=1}^{\text{card } \mathcal{A}} |\mathbf{f}^H \mathbf{x}[t] - a_j[t]|^2$$

where  $a_j[t] \in \mathcal{A}$

We have transformed an exhaustive search into a *polynomial* alphabet fit



# Algorithm AMiSRoF

## Absolute Minimum Search by Root Finding [GC98]

- Initialize  $\mathbf{f} = \mathbf{f}_o$
- For  $k = 1$  to  $k_{\max}$ , and while  $|\mu_k| > \text{threshold}$ , do
  - Compute gradient  $\mathbf{g}_k$  and Hessian  $\mathbf{H}_k$  at  $\mathbf{f}_{k-1}$
  - Compute a search direction  $\mathbf{v}_k$ , e.g.  $\mathbf{v}_k = \mathbf{H}_k^{-1} \mathbf{g}_k$
  - Normalize  $\mathbf{v}_k$  to  $\|\mathbf{v}_k\| = 1$
  - Compute the *absolute* minimum  $\mu_k$  of the rational function in  $\mu$ :

$$\Phi(\mu) \stackrel{\text{def}}{=} \Upsilon_T(\mathbf{f}_{k-1} + \mu \mathbf{v}_k)$$

- Set  $\mathbf{f}_k = \mathbf{f}_{k-1} + \mu_k \mathbf{v}_k$

# Algorithm ILSP

## Iterative Least-Squares with Projection [TVP96]

- Assumes that components  $s_i[n] \in \mathcal{A}$ , known alphabet
- Assumes columns of  $\mathbf{H}$  belong to a known *array manifold*
- Initialize  $\mathbf{H}$ , and start the loop
  - Compute LS estimate of matrix  $\mathbf{S}$  in equation  $\mathbf{X} = \mathbf{H}\mathbf{S}$
  - Project  $\mathbf{S}$  onto  $\mathcal{A}$
  - Compute LS estimate of  $\mathbf{H}$  in equation  $\mathbf{X} = \mathbf{H}\mathbf{S}$
  - Project  $\mathbf{H}$  onto the array manifold

## Part V

# Algorithms for convolutive mixtures

# Contents

Here limited to over-determined mixtures

- Blind equalization,
  - Modeling, Carrier offset
  - Contrast criteria
  - Algorithms (Pajod, subspace, linear prediction...)
- Blind identification
  - Cumulant matching
  - Algebraic approaches
  - Subspace techniques
  - ARMA mixtures

References

# Blind Equalization

- Modeling of Dynamic Mixtures
- Contrast-based
  - MISO Deflation
  - Para-unitary
- SIMO channel
  - subspace
  - mutually referenced
  - Linear prediction
- MIMO
- Matched Filter after Blind Identification

# SISO Modeling

- Sequence of symbols  $s[k]$  at a rate  $1/T_s$
- Overall channel  $h(t)$ , containing transmit&receive filters and propagation
- received process  $x(t) = \sum_{k \in \mathbb{Z}} h(t - k T_s) s[k]$
- If sampled at a rate  $1/T$ :

$$x[n] = \sum_{k \in \mathbb{Z}} h(n T - k T_s) s[k]$$

- If sampled exactly at symbol rate, we get a *discrete convolution*:

$$x[n] = \sum_{k \in \mathbb{Z}} h[n - k] s[k]$$

with  $h[m] \stackrel{\text{def}}{=} h(m T)$

# MIMO Modeling

In practice, one often assumes the approximation of discrete convolutive FIR:

$$\mathbf{x}[n] = \sum_{k=0}^L \mathbf{H}[k] \mathbf{s}[n - k] + \mathbf{v}[k] \quad (35)$$

Either:

- **Blind Identification**

Estimate the finite matrix sequence  $\mathbf{H}[k]$ , or

- **Blind Equalization**

Estimate a FIR filter  $\mathbf{F}[\ell]$ ,  $0 \leq \ell \leq L'$

## Carrier offset (1)

In practical contexts of Blind Techniques, carrier frequency might be unaccurately estimated

- In the SISO case, this yields

$$x[n] = \sum_k h[n-k] s[k] e^{j k \delta}$$

- An *equivalent writing* is

$$x[n] = e^{j n \delta} \sum_k h'[n-k] s[k]$$

where  $h'[m] \stackrel{\text{def}}{=} h[m] e^{j m \delta}$ .

- alphabet fitting at the output may be limited by the presence of this Carrier residual. But Blind Equalization is still feasible.



## Carrier offset (2)

- In the MIMO case, the carrier offset cannot be pulled into the channel anymore, unless all sources have *the same* carrier offset
- In fact on sensor  $k$ :

$$x_k[n] = \sum_{\ell} \sum_p H_{kp}[n - \ell] s_p[\ell] e^{j\ell\delta_p}$$

or

$$x_k[n] = \sum_{\ell} \sum_p e^{j n \delta_p} H'_{kp}[n - \ell] s_p[\ell]$$

with  $H'_{kp}[m] \stackrel{\text{def}}{=} H_{kp}[m] e^{j m \delta_p}$

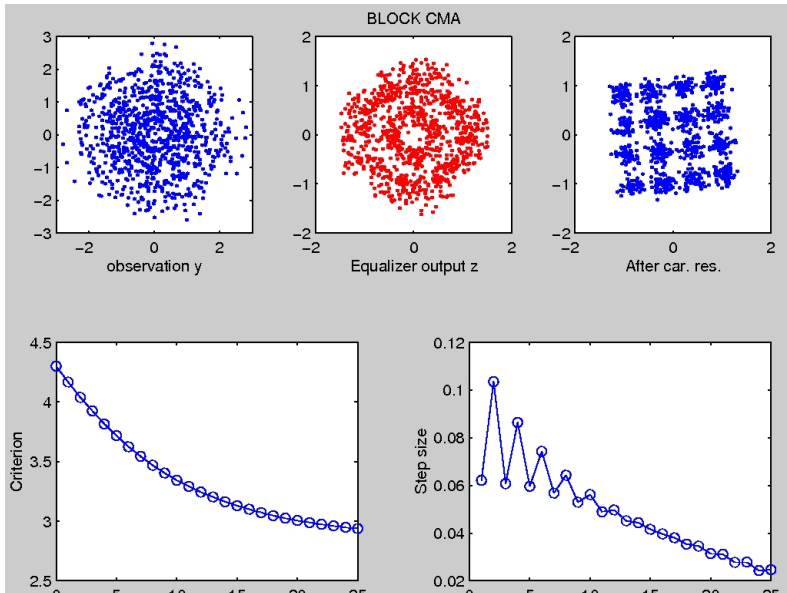
- Thus blind equalization is not possible anymore before carrier residual mitigation

# Carrier offset (3)

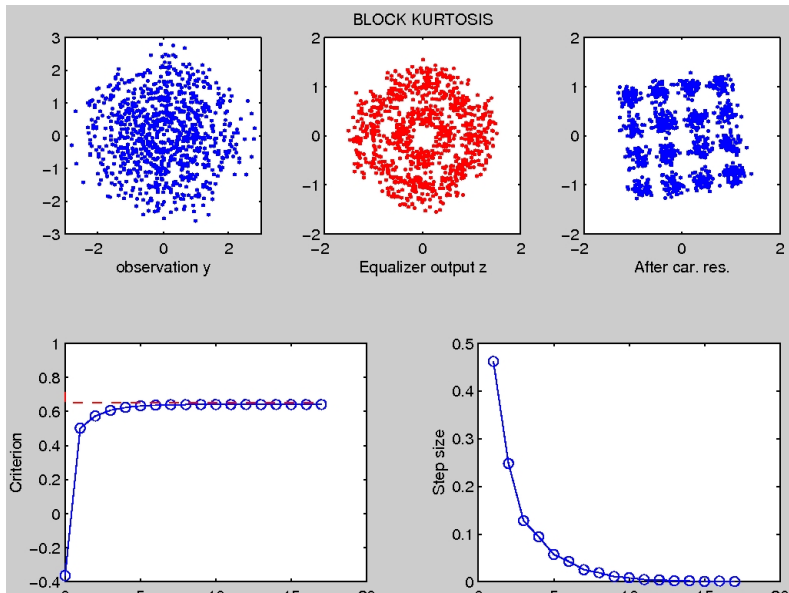
## Summary

- SISO case: BE and CO can be permuted
- MIMO case: BE and CO cannot generally be permuted

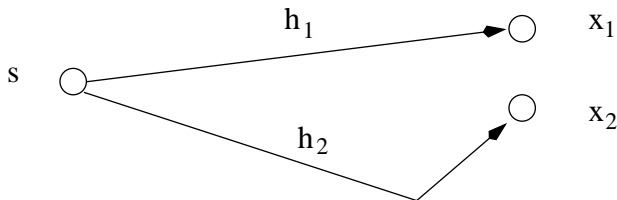
# Carrier offset (3)



# Carrier offset (4)



# SIMO mixture with diversity $K = 2$ (1)



**Disparity condition:**

$$h_1[z] \wedge h_2[z] = 1 \Rightarrow x_1[z] \wedge x_2[z] = s[z]$$

**Bézout:**

$$\begin{aligned} \exists v_1[z], v_2[z] / \quad v_1[z] h_1[z] + v_2[z] h_2[z] &= 1 \\ \Rightarrow v_1[z] x_1[z] + v_2[z] x_2[z] &= s[z] \end{aligned}$$

**Thus**

FIR filter  $\mathbf{h} = \begin{pmatrix} h_1 \\ h_2 \end{pmatrix}$  admits the FIR inverse  $\mathbf{v} = (v_1, v_2)$ .

# SIMO mixture with diversity $K = 2$ (2)

## Theorem

If two polynomials  $p(z) = \sum_{i=0}^m a_i z^i$  and  $q(z) = \sum_{i=0}^n b_i z^i$  are prime, then the resultant below is non zero:

$$\mathcal{R}(p, q) = \begin{vmatrix} a_0 & \dots & a_m & 0 & \dots \\ 0 & \ddots & & \ddots & 0 \\ 0 & 0 & a_0 & \dots & a_m \\ b_0 & \dots & b_n & 0 & \dots \\ 0 & \ddots & & \ddots & 0 \\ 0 & 0 & b_0 & \dots & b_n \end{vmatrix} \stackrel{\text{def}}{=} \det \begin{pmatrix} \mathbf{A} \\ \mathbf{B} \end{pmatrix}$$

# Use of time diversity

## Time diversity

If channel bandwidth exceeds symbol rate  $\frac{1}{T_s}$  (excess bandwidth), then a sampling faster than  $\frac{1}{T_s}$  brings extra information on channel. [?] [?]

## How to build a SIMO channel from a SISO?

- sample twice faster:  $x[k] = x(k T_s/2)$
- denote odd samples  $x_1[k] = x[2k + 1]$ , and even samples  $x_2[k] = x[2k]$
- then

$$\begin{pmatrix} x_1[k] \\ x_2[k] \end{pmatrix} = \begin{pmatrix} \mathbf{H}_1 \\ \mathbf{H}_2 \end{pmatrix} \mathbf{s}[k] \stackrel{\text{def}}{=} \mathbf{H} \mathbf{s}[k]$$

Matrix  $\mathbf{H}$  is full rank (well conditioned) if sufficient *excess bandwidth*

# Mutually Referenced Equalizers (1)

- Recall the compact modeling of equation (42) slide 279:

$$\mathbf{X}(n) = \mathcal{H}_T \mathbf{S}(n)$$

Then observe that if  $\mathcal{H}_T$  is column shaped and full rank (here  $T + L + 1$ ):

$$\exists \mathbf{V} : \quad \mathbf{V}^H \mathcal{H}_T = \mathbf{I}$$

- Each row of  $\mathbf{V}^H$  defines an equalizer  $\mathbf{v}_i^H$ , deduced from each other by a delay [?]:

$$\mathbf{v}_j^H \mathbf{X}(n - i) = \mathbf{v}_i^H \mathbf{X}(n - j) = s(n - i - j)$$



## Mutually Referenced Equalizers (2)

- The equations  $E\{|\mathbf{v}_k^H \mathbf{X}(n) - \mathbf{v}_{k+1}^H \mathbf{X}(n+1)|^2\} = 0$  for  $0 \leq k \leq T + L$  yield:

$$\mathcal{V}^H \mathcal{R} \mathcal{V} = 0 \quad \text{with}$$

$$\mathcal{V} \stackrel{\text{def}}{=} \begin{pmatrix} \mathbf{v}_0 \\ \mathbf{v}_1 \\ \vdots \\ \mathbf{v}_{T+L} \end{pmatrix}, \text{ and } \mathcal{R} \stackrel{\text{def}}{=} \begin{pmatrix} \mathbf{R}(0) & -\mathbf{R}(1)^H & 0 & \dots \\ -\mathbf{R}(1) & 2\mathbf{R}(0) & \ddots & 0 \\ 0 & \ddots & \ddots & \ddots \\ 0 & \dots & 0 & 2\mathbf{R}(0) & -\mathbf{R}(1) \\ \vdots & 0 & \ddots & -\mathbf{R}(1) & \mathbf{R}(0) \end{pmatrix}$$

and  $\mathbf{R}(k) \stackrel{\text{def}}{=} E\{\mathbf{X}(n+k)\mathbf{X}(n)^H\}$ .

- Thus, take  $\mathcal{V}$  as being the dominated eigenvector, and extract  $\mathbf{v}_k$  from it
- In practice, necessary to add a constraint to avoid  $\mathbf{v}_k \in \text{null}(\mathcal{H}^H)$

# Contrast criteria (1)

Proofs derived in the static case hold true in the convolutive case,  
e.g. family of contrasts of slides 73-74

Proofs...

## Contrast criteria (2)

- But also possible to devise new families of contrasts for para-unitary equalizers after prewhitening [?] [?]. For instance:

$$\Upsilon(\mathbf{y}) = \sum_i \sum_j \sum_p \sum_k \sum_q |\text{Cum}\{y_i[n], y_i[n], y_j[n-p], y_k[n-q]\}|^2 \quad (36)$$

- In the above, one can conjugate any of the variables  $y_\ell$ 's
- Holds true for almost any cumulants of order  $\geq 3$
- Only two indices need to be identical with same delay

Proof Based on the property that, for para-unitary filters  $\mathbf{G}$ :

$$\mathbf{y}[n] \stackrel{\text{def}}{=} \sum_t \mathbf{G}[n-t] s[t] \Rightarrow \Upsilon(\mathbf{y}) = \sum_i \sum_{\ell} \sum_t |G_{i\ell}[t]|^4 |\kappa_\ell|^2$$

# MISO Dynamic Extractor: Deflation

- Fixed step gradient Deflation [TUG97]
- Optimal Line search along a descent direction, OS-KMA [?]  
[ZC05]

# PAJOD (1)

- Technique applied *after* space-time prewhitening
- Then one looks for a *para-unitary* equalizer, by maximizing the contrast

$$\mathcal{J}_{2,r}(\mathbf{y}) = \sum_{\mathbf{b}} \sum_{\beta} \|\mathbf{Diag}\{\mathcal{H}^H \mathbf{M}(\mathbf{b}, \beta) \mathcal{H}\}\|^2$$

Matrix  $\mathcal{H}$  is now defined differently, and is semi-unitary.

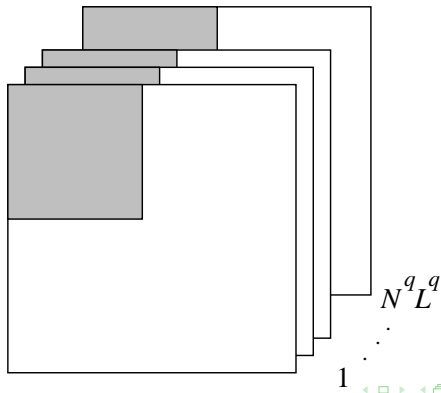
Matrices  $\mathbf{M}(\cdot)$  contain cumulants of whitened observations

- Contrast (36) is maximized again by a sweeping technique

## PAJOD (2)

PAJOD: Partial Approximate Joint Diagonalization of matrix slices

One actually attempts to diagonalize only a portion of the tensor



# MIMO Blind Equalization

- linear prediction after BI [?]
- linear prediction [?] [?] [?]
- subspace [?] [?] [?] [?] [?]
- identifiability issues by subspace techniques [?] [?]

# Equalization after prior Blind Identification

Assume channel  $\mathbf{H}[z]$  has been identified, with:

$$\mathbf{x}[n] = \mathbf{H}[z] \star \mathbf{s}[z] + \mathbf{v}[z]$$

An estimate of  $\mathbf{s}[z]$  is obtained with  $\mathbf{F}[z] \star \mathbf{x}[z]$ .

Possible equalizers  $\mathbf{F}[z]$ :

- Zero-Forcing:  $\mathbf{F}[z] = \mathbf{H}[z]^{-1}$
- Matched Filter:  $\mathbf{F}[z] = \mathbf{H}[1/z^*]^H$   
(used in MLSE; optimal if channel AWGN; maximizes output SNR)
- Minimum Mean Square Error (MSE):  
$$\mathbf{F}[z] = (\mathbf{H}[z]\mathbf{H}[1/z^*]^H + \mathbf{R}_v[z])^{-1}\mathbf{H}[1/z^*]^H$$

➡ One can insert soft or hard decision to stabilize the inverse, or to reduce noise, e.g. decision Feedback Equalizers (DFE).



# Overview

- Interest
- MA identifiability (second order vs hOS)
- SISO: Cumulant matching
- MIMO: Cumulant matching and linear prediction (non monic MA)
- Algebraic approaches, Quotient Ring
- SIMO: Subspace approaches
- MIMO: Subspace, IIR, ...

# BE vs BI

If sources  $s_p[k]$  are discrete, it is:

- rather easy to define a BE optimization criterion in order to match an output alphabet
- difficult to exploit a source alphabet in BI

**Example** the property of constant modulus of an alphabet is mainly used in Blind Equalization: CMA (Constant Modulus Algorithm)

# Interest of Blind Identification

- When the mixture does not have a stable inverse
  - ➡ When may want to control stability by soft/hard decision in a Feedback Equalizer
- When sources are not of interest (e.g. channel characteristics, localization only)

# SISO Cumulant matching (1)

- Consider first the **SISO** version of (35):

$$x[n] = \sum_{k=0}^L h[k] s[n-k] + v[k]$$

where  $v[k]$  is Gaussian stationary, and  $s[n]$  is 4th order white stationary.

- Then, by the multilinearity property of cumulants (slide 52):

$$C_x(i, j) \stackrel{\text{def}}{=} \text{Cum}\{x[t+i], x[t+j], x[t+L], x[t]\} = h[i] h[j] h[L] h[0] c_s$$

with  $c_s \stackrel{\text{def}}{=} \text{Cum}\{s[n], s[n], s[n], s[n]\}$ .

- By substitution of the unknown  $h[L] h[0] c_s$ , one gets a whole family of equations [?] [?]:

$$h[i] h[j] C_x(k, \ell) = h[k] h[\ell] C_x(i, j), \quad \forall i, j, k, \ell \quad (37)$$

## SISO Cumulant matching (2)

- A solution to the subset of (37) for which  $j = \ell$  can be easily obtained:

$$h[i] C_x(k, j) - h[k] C_x(i, j) = 0, \quad 0 \leq i < k \leq L, \quad 0 \leq j \leq L \quad (38)$$

- This is a *linear system* of  $L(L + 1)^2/2$  equations in  $L + 1$  unknowns  
 $\Rightarrow$  **Least Square (LS) solution**, up to a scale factor (e.g.  $h(0) = 1$ ).
- Since 4th order only, asymptotically (for large samples) insensitive to Gaussian noise.
- Total Least Squares (TLS) solution possible as well

# MIMO Cumulant matching (1)

## Inteterminacy

- Scale (scalar) factor for SISO, but  $\mathbf{\Lambda P}$  factor for MIMO

## Reduction to a monic model [COM94b]

- If  $\mathbf{H}[0]$  is invertible, (35) can be rewritten as

$$\mathbf{y}[n] = \mathbf{H}[0] \mathbf{s}[n], \quad (39)$$

$$\mathbf{x}[n] = \sum_{k=0}^L \mathbf{B}[k] \mathbf{y}[n-k] + \mathbf{w}[k] \quad (40)$$

where  $\mathbf{B}[k] \stackrel{\text{def}}{=} \mathbf{H}[k] \mathbf{H}[0]^{-1}$ .

- Because  $\mathbf{B}[0] = \mathbf{I}$ , MA model (40) is said to be *monic*.
- *Indeterminacy* is only in (39), which is solved by ICA if  $\mathbf{s}[n]$  is spatially white

# MIMO Cumulant matching (2)

## Kronecker notation

- Store 4th order cumulant tensors in vector form:

$$\mathbf{c}_{\mathbf{a},\mathbf{b},\mathbf{c},\mathbf{d}} \stackrel{\text{def}}{=} \mathbf{vec}\{\text{Cum}\{\mathbf{a}, \mathbf{b}, \mathbf{c}, \mathbf{d}\}\}$$

- Then, we have the property (where  $*$  denotes term-wise product):

$$\begin{aligned} \mathbf{c}_{\mathbf{a},\mathbf{b},\mathbf{c},\mathbf{d}} = & E\{\mathbf{a} \otimes \mathbf{b} \otimes \mathbf{c} \otimes \mathbf{d}\} - E\{\mathbf{a} \otimes \mathbf{b}\} \otimes E\{\mathbf{c} \otimes \mathbf{d}\} - E\{\mathbf{a} \otimes E\{\mathbf{b} \\ & - E\{\mathbf{a} \otimes \mathbf{1}_\beta \otimes \mathbf{c} \otimes \mathbf{1}_\delta\} * E\{\mathbf{1}_\alpha \otimes \mathbf{b} \otimes \mathbf{1}_\gamma \otimes \mathbf{d}\} \end{aligned}$$

## MIMO Cumulant matching (3)

Assume monic MA model (40) where  $\mathbf{s}[n]$  white in time and  $L$  fixed, and denote

$$\mathbf{c}_x(i, j) \stackrel{\text{def}}{=} \mathbf{vec}\{\text{Cum}\{\mathbf{x}[t+i], \mathbf{x}[t+j], \mathbf{x}[t+L], \mathbf{x}[t]\}\}$$

- Then we can prove [?]:

$$C_x(i, j) = C_x(0, j) \mathbf{B}[i]^T, \quad \forall j, 0 \leq j \leq L$$

where  $C_x(i, j) \stackrel{\text{def}}{=} \mathbf{Unvec}_P(\mathbf{c}_x)$  is  $P^3 \times P$

- For every fixed  $i$ ,  $\mathbf{B}[i]$  is obtained by solving the system of  $(L+1)P^4$  equations in  $P^2$  unknowns in LS sense:

$$[\mathbf{I}_P \otimes C_x(0, j)] \mathbf{vec}\{\mathbf{B}[i]\} = \mathbf{c}_x(i, j) \quad (41)$$



# MIMO Cumulant matching (4)

## ■ Summary of the algorithm

- Choose a maximum  $L$
- Estimate cumulants of observation,  $\mathbf{c}_x(i, j)$  for  $i, j \in \{0, \dots, L\}$
- Solve the  $(L + 1)$  systems (41) in  $\mathbf{B}[i]$
- Compute the residue  $\mathbf{y}[t]$  (*Linear Prediction*)
- Solve the ICA problem  $\mathbf{y}[t] = \mathbf{H}[0] \mathbf{s}[t]$

## ■ Weaknesses

- $\mathbf{H}[0]$  must be invertible
- FIR model (35) needs to have a stable inverse

# Algebraic Blind identification (1)

## Types of discrete source studied

- BPSK:  $b[k] \in \{-1, 1\}$ , *i.i.d.*
- MSK:  $m[k+1] = j m[k] b[k]$
- QPSK:  $p[k] \in \{-1, -j, 1, j\}$ , *i.i.d.*
- $\frac{\pi}{4}$ -DQPSK:  $d[k+1] = e^{j\pi/4} d[k] p[k]$
- 8-PSK:  $q[k] \in \{e^{jn\pi/4}, n \in \mathbb{Z}\}$ , *i.i.d.*
- etc...

# Algebraic Blind identification (2)

## Input/Output relations:

- For  $s[k]$  **BPSK**:  $E\{x[n] x[n - \ell]\} = s[0]^2 \sum_{m=0}^L h[m] h[m + \ell]$
- For  $s[k]$  **MSK**:  
 $E\{x[n] x[n - \ell]\} = s[0]^2 \sum_{m=0}^L (-1)^m h[m] h[m + \ell]$
- For  $s[k]$  **QPSK**:  
 $E\{x[n]^2 x[n - \ell]^2\} = s[0]^2 \sum_{m=0}^L h[m]^2 h[m + \ell]^2$
- etc..

# Algebraic Blind identification (3)

## Principle:

- Compute all roots of the polynomial system in  $h[n]$ .  
For instance for MSK sources and a channel of length 2 [?]:

$$h[0]^2 - h[1]^2 + h[2] = \beta_0$$

$$h[0] h[1] - h[1] h[2] = \beta_1$$

$$h[0] h[2] = \beta_2$$

- Choose among these roots the one that best matches the I/O correlation:

$$E\{x[n] x[n - \ell]^*\} = \sum_{m=0}^L h[m] h[m + \ell]^*$$

# Algebraic Blind identification (4)

**Theorem (Bezout)** A polynomial system of degree  $d$  in  $N$  variables has either:

- infinitely many solutions
- no solution
- exactly  $d^N$  solutions (distinct or not)



**Etienne Bézout, 1730-1783**

# Algebraic Blind identification (5)

## ■ Standard approaches

- Gröbner bases

## ■ Efficient solution of polynomial system: Normal Forms

There are two approaches, both working in the Quotient Ring modulo the Ideal defined by polynomial system:

- Eigenvectors of the transposed multiplication matrix  $\mathbf{M}_u^T$  in the Quotient Ring
- The Rational Univariate Representation (RUR) of eigenvalues of  $\mathbf{M}_u$

**Main advantage:** most (symbolic) calculations depend *only* on distribution of  $s[n]$ , and may thus be *stored in ROM* → Limited numerical computations left depending on measurements.

# SIMO mixture (1)

- FIR of length  $L$  and dimension  $K$ :

$$\mathbf{x}(n) = \sum_{i=1}^L \mathbf{h}(i)s(n-i) + \mathbf{b}(n)$$

with:  $\mathbb{E}\{\mathbf{b}(m)\mathbf{b}(n)^H\} = \sigma_b^2 \mathbf{I} \delta(m-n)$   
 and  $\mathbb{E}\{\mathbf{b}(m)s(n)^*\} = 0$

- For  $T$  successive values:

$$\begin{pmatrix} \mathbf{x}(n) \\ \mathbf{x}(n-1) \\ \vdots \\ \mathbf{x}(n-T) \end{pmatrix} = \begin{pmatrix} \mathbf{h}(0) & \mathbf{h}(1) & \dots & \mathbf{h}(L) & 0 & \dots & 0 \\ 0 & \mathbf{h}(0) & \dots & \dots & \mathbf{h}(L) & \dots & 0 \\ \vdots & \ddots & \ddots & \ddots & & \ddots & \vdots \\ 0 & 0 & \dots & \mathbf{h}(0) & \mathbf{h}(1) & \dots & \mathbf{h}(L) \end{pmatrix} \begin{pmatrix} s(n) \\ \vdots \\ s(n-T-L) \end{pmatrix}$$

Or in compact form:

$$\mathbf{X}(n:n-T) = \mathcal{H}_T \mathbf{S}(n:n-T-L) \quad (42)$$

Here,  $\mathcal{H}_T$  is of size  $(T+1)K \times (T+L+1)$

## SIMO mixture (2)

### ■ Condition of “column” matrix

$\mathcal{H}$  has strictly more rows than columns **iff**

$$(T + 1)K > T + L + 1$$

$$\Leftrightarrow T > L/(K - 1) - 1 \Leftarrow T \geq L$$

It suffices that  $T$  exceeds channel memory.

### ■ Disparity condition

Columns of  $\mathcal{H}$  are linearly independent **iff**

$$\mathbf{h}[z] \neq \mathbf{0}, \forall z$$

### ■ Noise subspace

Under these 2 conditions, there exists a “noise subspace”:

$$\exists \mathbf{v} / \mathbf{v}^H \mathcal{H}_T = \mathbf{0}$$



# SIMO mixture (3)

## Properties of vectors in the null space

- Since  $\mathbf{R}_x \stackrel{\text{def}}{=} \mathbb{E}\{\mathbf{X}\mathbf{X}^H\} = \mathcal{H}_T \mathcal{H}_T^H + \sigma_b^2 \mathbf{I}$ ,  
vectors  $\mathbf{v}^{(p)}$  of noise space can be computed from  $\mathbf{R}_x$ :

$$\mathbf{R}_x \mathbf{v}^{(p)} = \sigma_b^2 \mathbf{v}^{(p)}$$

- And since convolution is commutative:

$$\mathbf{v}^{(p)H} \mathcal{H}_T = \mathbf{h}^H \mathcal{V}^{(p)}$$

where  $\mathcal{V}^{(p)}$  block Töplitz, built on  $\mathbf{v}^{(p)}$ .

- Thus  $\mathbf{h}^H = [\mathbf{h}(0)^H, \mathbf{h}(1)^H, \dots, \mathbf{h}(L)^H]$  are obtained by computing the left singular vector common to  $\mathcal{V}^{(p)}$ .

# SIMO mixture (4)

## Summary of the SIMO Subspace Algorithm

- Choose  $T \geq L$
- Compute  $\mathbf{R}_x$ , correlation matrix of size  $(T + 1)K$
- Compute the  $d = T(K - 1) + K - L - 1$  vectors  $\mathbf{v}^{(p)}$  of the noise space
- Compute vector  $\mathbf{h}$  minimizing the quadratic form

$$\mathbf{h}^H \left[ \sum_{p=1}^d \mathbf{v}^{(p)} \mathbf{v}^{(p)H} \right] \mathbf{h}$$

- Cut  $\mathbf{h}$  into  $L + 1$  slices  $\mathbf{h}(i)$  of length  $K$

Under the assumed hypotheses, the solution is unique up to a scalar scale factor [?]

# SIMO mixture (5)

## Summary of the SIMO Subspace Algorithm when $K = 2$

- Choose  $T = L$ . There is a single vector  $\mathbf{v}$  in the noise space
- Compute  $\mathbf{R}_x$ , correlation matrix of size  $(T + 1)K$
- Compute the vector  $\mathbf{v}$  of the noise space
- Cut  $\mathbf{v}$  into  $L + 1$  slices  $\mathbf{v}(i)$  of length  $K = 2$

- Compute  $\mathbf{h}(i) = \begin{pmatrix} 0 & -1 \\ 1 & 0 \end{pmatrix} \mathbf{v}(i)$

In fact  $x_i = h_i \star s \Rightarrow h_2 \star x_1 - h_1 \star x_2 = 0$

Approach called **SRM** (*Subchannel Response Matching*) [?]  
[?]

# SISO Identifiability

## ■ Second order statistics

■  $\alpha_\ell = E\{x[n]x[n-\ell]^*\}$  ➡ allow to estimate  $|h[m]|$

■  $\beta_\ell = E\{x[n]x[n-\ell]\}$  ➡ allow to estimate  $h[m]$  if  $E\{s^2\} \neq 0$

## ■ Fourth order statistics ➡ many (polynomial) additional equations

■  $\gamma_{0jkl} = \text{Cum}\{x[n], x[n-j], x[n-k], x[n-\ell]\}$

■  $\gamma_{0j}^{k\ell} = \text{Cum}\{x[n], x[n-j], x[n-k]^*, x[n-\ell]^*\}$

If some sources are 2nd order circular, sample Statistics of order higher than 2 are mandatory, but *otherwise not* [?] !

# SIMO Identifiability

With a receive diversity, (deterministic) identifiability conditions are weaker [?] [?]

- **Definition** A length- $N$  input sequence  $s[n]$  has  $P$  modes iff the Hankel matrix below is full row rank:

$$\begin{pmatrix} s[1] & s[2] & \dots & s[N-p+1] \\ s[2] & s[3] & \ddots & s[N-p+2] \\ \vdots & \vdots & & \vdots \\ s[p] & s[p+1] & \dots & s[N] \end{pmatrix}$$

- **Theorem** A  $K \times L$  FIR channel  $\mathbf{h}$  is identifiable if:
  - Channels  $h_k[z]$  do not have common zeros
  - The observation length of each  $x_k[n]$  must be at least  $L+1$
  - The input sequence should have at least  $L+1$  modes (sufficiently exciting)

## Subspace algorithm for MIMO mixtures

- Similarly to the SIMO case, we have the compact form:

$$\mathbf{X}(n) = \mathcal{H}_T \mathbf{S}(n) + \mathbf{B}(n)$$

where  $\mathcal{H}$  is now built on matrices  $\mathbf{H}(k)$ ,  $1 \leq k \leq L$ , and is of size  $(T+1)K \times (T+L+1)P$ .

- For large enough  $T$ , this matrix is "column shaped"
- Again  $\mathbf{R}_x = \mathcal{H}_T \mathcal{H}_T^H + \sigma_b^2 \mathbf{I}$
- But now, vectors of the noise space characterize  $\mathbf{H}[z]$  only up to a constant post-multiplicative matrix  $\Rightarrow$  *ICA must be used afterwards*
- Foundations of the MIMO subspace algorithm are more complicated [Loubaton'99]

In the MIMO case, HOS are in general mandatory.

# SISO ARMA mixtures

## What are the tools when the channel is IIR?

- In general, just consider it as a FIR (truncation)  $\rightarrow$  already seen
- But also possible to assume presence of a recursive part
  - Define I/O relation:  $\sum_{i=0}^p a_i x[n-i] = \sum_{j=0}^q b_j w[n-j]$   
where  $w[\cdot]$  is i.i.d. and  $a_0 = b_0 = 1$
  - Second order  $c_x(\tau) \stackrel{\text{def}}{=} E\{x[n]x[n+\tau]\}$  can be used to identify  $a_k$ :

$$\sum_{k=0}^p a_k c_x(\tau - k) = 0, \quad \forall \tau > q$$

- Then compute the residue and identify  $b_\ell$  with HOS (cf. slide 269)
- Also possible with HOS only for AR part [?]

# MIMO ARMA mixtures (1)

## Results of SISO case can be extended

- Take a  $K$ -dimensional ARMA model: Define I/O relation:

$$\sum_{i=0}^p \mathbf{A}_i \mathbf{x}[n-i] = \sum_{j=0}^q \mathbf{B}_j \mathbf{w}[n-j]$$

where  $w[\cdot]$  is i.i.d. and  $\mathbf{A}_0 = \mathbf{I}$  and  $\mathbf{B}_0$  invertible

- For instance at order 4, AR identification is based on:
  - $\sum_{j=1}^p \mathbf{A}_j \bar{\mathbf{c}}_x(t, \tau - j) = -\bar{\mathbf{c}}_x(t, \tau), \forall \tau > q, \forall t$
  - with  $\bar{\mathbf{c}}_x(i, j) \stackrel{\text{def}}{=} \mathbf{Unvec}_K(\text{Cum}\{\mathbf{x}[n], \mathbf{x}[n], \mathbf{x}[n+i], \mathbf{x}[n+j]\})$



# MIMO ARMA mixtures (2)

## Limitations

- Sources need to be *linear processes*
- $\mathbf{B}_0$  needs to be invertible
- AR residuals need to be computed (MA filtering) to compute  $\mathbf{B}_i$
- One can compute MA residuals (AR filtering) if input  $\mathbf{s}[n]$  is requested  $\rightarrow$  but might be unstable

## Part VI

### Algorithms for under-determined mixtures

## Back to Essential uniqueness

- Recall the general model (22) to fit (here 3rd order):

$$\varepsilon \stackrel{\text{def}}{=} \left\| \mathbf{T} - \sum_{q=1}^{\text{rank}\{\mathbf{T}\}} \mathbf{a}^{(q)} \circ \mathbf{b}^{(q)} \circ \mathbf{c}^{(q)} \right\|^2 \quad (43)$$

- For instance, if  $(\mathbf{A}, \mathbf{B}, \mathbf{C})$  is solution, so is  $(\mathbf{A}\mathbf{P}\mathbf{\Lambda}, \mathbf{B}\mathbf{P}\mathbf{\Delta}, \mathbf{C}\mathbf{P}\mathbf{\Lambda}^{-1}\mathbf{\Delta}^{-1})$
- Essential uniqueness*: uniqueness up to a common scale-permutation ambiguity.
- The scale indetermination can be fixed by introducing a diagonal tensor  $\mathbf{\Delta}$  and imposing unit-norm columns in the matrices:

$$\mathbf{T} \approx \mathbf{\Delta} \bullet_1 \mathbf{A} \bullet_2 \mathbf{B} \bullet_3 \mathbf{C}$$

# Essential uniqueness

**Sufficient condition** The *Kruskal rank* of a matrix  $\mathbf{A}$  is the maximum number  $k_A$ , such that any subset of  $k_A$  columns are linearly independent.

**Kruskal's bound** [KRU77] [SB00] [SS07] gives *sufficient conditions*. Essential uniqueness is ensured if the tensor rank  $R$  is below an upper bound:

- $2R + 2 \leq k_A + k_B + k_C$ ,
- or generically, for a tensor of order  $d$  and dimensions  $N_\ell$ :

$$2R + d - 1 \leq \sum_{\ell=1}^d \min(N_\ell, R)$$

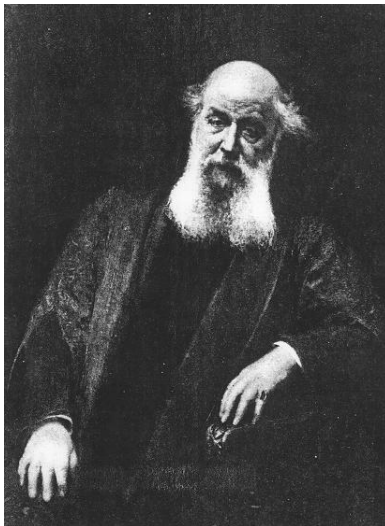
## Essential uniqueness (cont'd)

**Necessary and sufficient condition** **Essential uniqueness** has been proved via *local identifiability*, under the condition that the rank is *sub-generic*:

$$\text{rank}\{\mathbf{T}\} \leq \left\lceil \frac{\prod_{\ell} N_{\ell}}{\sum_{\ell} (N_{\ell} - 1) + 1} \right\rceil$$

This condition is necessary and sufficient up to some exceptions, for which the maximal rank should be decreased by 1. The proof is numerical for the general case [CtB08], but algebraic in the symmetric case [CGLM08].

**Questions:** What algorithms, and under what conditions?



James Joseph Sylvester (1814–1897)

## Binary case

### Construction of the CanD (1)

#### Sylvester's theorem in $\mathbb{R}$

- A binary quantic  $p(x, y) = \sum_{i=0}^d \gamma_i c(i) x^i y^{d-i}$  can be decomposed in  $\mathbb{R}[x, y]$  into a sum of  $r$  powers as  $p(x, y) = \sum_{j=1}^r \lambda_j (\alpha_j x + \beta_j y)^d$  if and only if the form

$$q_c(x, y) = \prod_{j=1}^r (\beta_j x - \alpha_j y) = \sum_{l=0}^r g_l x^l y^{r-l}$$

satisfies

$$\begin{bmatrix} \gamma_0 & \gamma_1 & \cdots & \gamma_r \\ \gamma_1 & \gamma_2 & \cdots & \gamma_{r+1} \\ \vdots & & & \vdots \\ \gamma_{d-r} & & \cdots & \gamma_d \end{bmatrix} \begin{bmatrix} g_0 \\ g_1 \\ \vdots \\ g_r \end{bmatrix} = 0.$$

and has distinct real roots.

- Valid even in non generic cases.

## Construction of the CanD (2)

### Sylvester's theorem in $\mathbb{C}$

A binary quantic  $p(x, y) = \sum_{i=0}^d c(i) \gamma_i x^i y^{d-i}$  can be written as a sum of  $d$ th powers of  $r$  distinct linear forms:

$$p(x, y) = \sum_{j=1}^r \lambda_j (\alpha_j x + \beta_j y)^d, \quad (44)$$

if and only if **(i)** there exists a vector  $\mathbf{g}$  of dimension  $r + 1$ , with components  $g_\ell$ , such that

$$\begin{bmatrix} \gamma_0 & \gamma_1 & \cdots & \gamma_r \\ \vdots & & & \vdots \\ \gamma_{d-r} & \cdots & \gamma_{d-1} & \gamma_d \end{bmatrix} \mathbf{g}^* = \mathbf{0}. \quad (45)$$

and **(ii)** the polynomial  $q(x, y) \stackrel{\text{def}}{=} \sum_{\ell=0}^r g_\ell x^\ell y^{r-\ell}$  admits  $r$  distinct roots



# Algorithm

- Start with  $r = 1$  ( $d \times 2$  matrix) and increase  $r$  until it loses its column rank

|   |   |
|---|---|
| 1 | 2 |
| 2 | 3 |
| 3 | 4 |
| 4 | 5 |
| 5 | 6 |
| 6 | 7 |
| 7 | 8 |



|   |   |   |
|---|---|---|
| 1 | 2 | 3 |
| 2 | 3 | 4 |
| 3 | 4 | 5 |
| 4 | 5 | 6 |
| 5 | 6 | 7 |
| 6 | 7 | 8 |



|   |   |   |   |
|---|---|---|---|
| 1 | 2 | 3 | 4 |
| 2 | 3 | 4 | 5 |
| 3 | 4 | 5 | 6 |
| 4 | 5 | 6 | 7 |
| 5 | 6 | 7 | 8 |



# Symmetric tensors of larger dimension

We have seen the link with polynomials in slide 124. The idea is to extend Sylvester's algorithm to more than 2 variables.

- $Xx$

- $Xx$

xx Tsigaridas Mourrain Comon

# Iterative algorithms

Continue to keep 3rd order case to illustrate the reasoning. Define

$$\mathbf{p} \stackrel{\text{def}}{=} \begin{bmatrix} \text{vec}\{\mathbf{A}^T\} \\ \text{vec}\{\mathbf{B}^T\} \\ \text{vec}\{\mathbf{C}^T\} \end{bmatrix}, \text{ and the gradient } \mathbf{g} = \begin{bmatrix} \mathbf{g}_A \\ \mathbf{g}_B \\ \mathbf{g}_C \end{bmatrix}$$

# Gradient (1)

- Newton update rule:  $\mathbf{p}(k+1) = \mathbf{p}(k) - \mathbf{H}(k)^{-1} \mathbf{g}(k)$   
Pure gradient:  $\mathbf{p}(k+1) = \mathbf{p}(k) - \mu(k) \mathbf{g}(k)$
- Systematic step variation:
  - $\mu(k)$  constant if  $\varepsilon(k) - \varepsilon(k+1) > 0.005 \varepsilon(k)$
  - $\mu$  increased via  $\mu(k+1) = 1.1 \mu(k)$  if  $0 \leq \varepsilon(k) - \varepsilon(k+1) \leq 0.005 \varepsilon(k)$
  - $\mu$  decreased via  $\mu(k+1) = \mu(k)/2$  if  $\varepsilon(k) < \varepsilon(k+1)$

## Gradient (2)

Closed-form expressions of the gradients of  $\varepsilon$  (43):

$$\begin{aligned} \mathbf{g}_A &= [\mathbf{I}_A \otimes (\mathbf{C}^H \mathbf{C} \oslash \mathbf{B}^H \mathbf{B})] \text{vec}\{\mathbf{A}^T\} - [\mathbf{I}_A \otimes (\mathbf{C} \odot \mathbf{B})] \text{vec}\{\mathbf{T}_{KJ \times I}\} \\ \mathbf{g}_B &= [\mathbf{I}_B \otimes (\mathbf{A}^H \mathbf{A} \oslash \mathbf{C}^H \mathbf{C})] \text{vec}\{\mathbf{B}^T\} - [\mathbf{I}_A \otimes (\mathbf{A} \odot \mathbf{C})] \text{vec}\{\mathbf{T}_{IK \times J}\} \\ \mathbf{g}_C &= [\mathbf{I}_C \otimes (\mathbf{B}^H \mathbf{B} \oslash \mathbf{A}^H \mathbf{A})] \text{vec}\{\mathbf{C}^T\} - [\mathbf{I}_C \otimes (\mathbf{B} \odot \mathbf{A})] \text{vec}\{\mathbf{T}_{JI \times K}\} \end{aligned}$$

where  $\oslash$  denotes the elementwise product (Hadamard)

# Quasi-Newton (1)

Define Jacobians with respect to matrices **A**, **B** and **C**, and the joint Jacobian:

$$\mathbf{J}_A = \mathbf{I}_A \otimes (\mathbf{C} \odot \mathbf{B})$$

$$\mathbf{J}_B = \Pi_1 [\mathbf{I}_B \otimes (\mathbf{A} \odot \mathbf{C})]$$

$$\mathbf{J}_C = \Pi_2 [\mathbf{I}_C \otimes (\mathbf{B} \odot \mathbf{A})]$$

where  $\Pi_i$  are appropriately chosen permutations, and

$$\mathbf{J} = [\mathbf{J}_A, \mathbf{J}_B, \mathbf{J}_C]$$

## Quasi-Newton (2)

- Quasi-Newton iteration:

$$\mathbf{p}(k+1) = \mathbf{p}(k) - [\mathbf{J}(k)^H \mathbf{J}(k) + \mathbf{M}(k)]^{-1} \mathbf{g}(k)$$

where matrix  $\mathbf{M}(k)$  is updated from  $\mathbf{J}(k)$ ,  $\mathbf{M}(k)$ ,  $\mathbf{g}(k)$  and  $\mathbf{p}(k)$ .

- The Levenberg-Marquardt iteration takes the form:

$$\mathbf{p}(k+1) = \mathbf{p}(k) - [\mathbf{J}(k)^H \mathbf{J}(k) + \lambda(k) \mathbf{I}]^{-1} \mathbf{g}(k)$$

where  $\lambda(k)$  is updated according to a specific rule, depending on the quality of the approximation of the objective:

$$\varepsilon(\mathbf{p} + \delta) - \varepsilon(\mathbf{p}) \approx \delta^H \mathbf{g} + \frac{1}{2} \delta^H (\mathbf{J}^H \mathbf{J} + \lambda \mathbf{I}) \delta$$

## Gradient algorithms for tensors with symmetries

In the presence of symmetries, the gradient takes a simpler form, given here for clarity in the case of 3rd order tensors, with symmetry in the first 2 modes, i.e. [?]:

$$T_{i,j,k} = T_{\sigma(i,j),k}$$

We have two matrices to determine, **A** and **C** since:

$$\varepsilon = \left\| \mathbf{T} - \sum_{q=1}^R \mathbf{a}(q) \circ \mathbf{a}(q) \circ \mathbf{c}(q) \right\|^2$$

The gradient and the Jacobian are of the form

$$\begin{aligned} \mathbf{g} &= \begin{bmatrix} \mathbf{g}_A + \mathbf{g}_B \\ \mathbf{g}_C \end{bmatrix} \\ \mathbf{J} &= [\mathbf{J}_A + \mathbf{J}_B, \mathbf{J}_C] \end{aligned}$$

where **B** is set to **B** = **A**.



## Other minimization algorithms

Algorithms using explicit expressions of the Hessian

- Newton:  $\mathbf{p}(k+1) = \mathbf{p}(k) - \mathbf{H}(k)^{-1} \mathbf{g}(k)$
- Conjugate Gradient: e.g. the “Multilinear Engine” [PAA99]
- etc...

➡ More costly in terms of memory and complexity per iteration, but fewer iterations needed.

➡ Do not solve the problem of *local minima*

## Compact writing of Objective

The objective function (43) can be written as:

$$\varepsilon = \|\mathbf{T}_{I \times KJ} - \mathbf{A}(\mathbf{C} \odot \mathbf{B})^T\|^2 \quad (46)$$

**Advantage:** compact writing of the best matrix  $\mathbf{A}$ , for fixed  $\mathbf{B}$  and  $\mathbf{C}$ , since (46) is quadratic in  $\mathbf{A}$  [HL94]:

$$\hat{\mathbf{A}} = \mathbf{T}_{I \times KJ} \cdot \{(\mathbf{C} \odot \mathbf{B})^T\}^\dagger$$

where  $\dagger$  denotes pseudo-inverse.

Similarly:

$$\begin{aligned} \|\mathbf{T}_{J \times IK} - \mathbf{B}(\mathbf{A} \odot \mathbf{C})^T\|^2 &\rightarrow \hat{\mathbf{B}} = \mathbf{T}_{J \times IK} \cdot \{(\mathbf{A} \odot \mathbf{C})^T\}^\dagger \\ \|\mathbf{T}_{K \times JI} - \mathbf{C}(\mathbf{B} \odot \mathbf{A})^T\|^2 &\rightarrow \hat{\mathbf{C}} = \mathbf{T}_{K \times JI} \cdot \{(\mathbf{B} \odot \mathbf{A})^T\}^\dagger \end{aligned}$$

# Alternating Least Squares algorithm (1)

Start with arbitrary  $\mathbf{B}(0)$  and  $\mathbf{C}(0)$

For  $k = 1 \dots k_{max}$ ,

- $\mathbf{A}(k+1) = \mathbf{T}_{I \times KJ} \cdot \{(\mathbf{C}(k) \odot \mathbf{B}(k))^T\}^\dagger$
- $\mathbf{B}(k+1) = \mathbf{T}_{J \times IK} \cdot \{(\mathbf{A}(k+1) \odot \mathbf{C}(k))^T\}^\dagger$
- $\mathbf{C}(k+1) = \mathbf{T}_{K \times JI} \cdot \{(\mathbf{B}(k+1) \odot \mathbf{A}(k+1))^T\}^\dagger$

Hence the ALS algorithm also needs that:

$$R \leq \min(JK, IK, IJ)$$

According to Kruskal [KRU89], this inequality is always satisfied.

## Alternating Least Squares algorithm (2)

**Another compact writing** [COM04a]: jointly diagonalize slices of lower order:

$$\varepsilon = \sum_{i=1}^I \|\mathbf{T}[i] - \mathbf{B}\mathbf{\Lambda}[i]\mathbf{C}^T\|^2$$

where  $\mathbf{\Lambda}[i] = \mathbf{Diag}\{A_{i1}, \dots, A_{iR}\}$ . Let  $\mathbf{\Lambda}[i]$  denote the vector containing the diagonal of  $\mathbf{\Lambda}[i]$ , and  $\mathbf{t}[i] \stackrel{\text{def}}{=} \mathbf{vec}\{\mathbf{T}[i]\}$ . Hence:

$$\varepsilon = \sum_i \left\| \mathbf{t}[i] - \sum_{q=1}^R \lambda_q[i] \mathbf{c}[q] \otimes \mathbf{b}[q] \right\|^2 \stackrel{\text{def}}{=} \sum_i \left\| \mathbf{t}[i] - \mathbf{M} \boldsymbol{\lambda}[i] \right\|^2 \quad (47)$$

Then stationary values are:

$$\begin{aligned} \mathbf{B} &= \left\{ \sum_k \mathbf{T}[k] \mathbf{C} \mathbf{\Lambda}[k] \right\} \left\{ \sum_\ell \mathbf{\Lambda}[\ell] \mathbf{C}^T \mathbf{C} \mathbf{\Lambda}[\ell] \right\}^{-1} \\ \mathbf{C} &= \left\{ \sum_k \mathbf{T}[k]^T \mathbf{B} \mathbf{\Lambda}[k] \right\} \left\{ \sum_\ell \mathbf{\Lambda}[\ell] \mathbf{B}^T \mathbf{B} \mathbf{\Lambda}[\ell] \right\}^{-1} \end{aligned}$$

# ALS for symmetric tensors (1)

For clarity, take a symmetric tensor  $\mathbf{T}$  of order 4:

- One can force symmetry in the iteration of page 307:

Start with arbitrary  $\mathbf{A}(0)$ ,  $\mathbf{A}(1)$ ,  $\mathbf{A}(2)$

For  $k = 2 \dots k_{\max}$ ,

**Soft forcing:**

$$\mathbf{A}(k+1) = \mathbf{T}_{I \times I^3} \cdot \{(\mathbf{A}(k) \odot \mathbf{A}(k-1) \odot \mathbf{A}(k-2))^{\mathsf{T}}\}^{\dagger}$$

**Hard forcing:**  $\mathbf{A}(k+1) = \mathbf{T}_{I \times I^3} \cdot \{(\mathbf{A}(k) \odot \mathbf{A}(k) \odot \mathbf{A}(k))^{\mathsf{T}}\}^{\dagger}$

Obviously applies at any order  $d \geq 3$  [?].

## ALS for symmetric tensors (2)

**More tricky iteration** based on compact writing of page 308.  
When  $\mathbf{T}$  is real symmetric:

$$\varepsilon = \sum_i \|\mathbf{T}[i] - \mathbf{B}\boldsymbol{\Lambda}[i]\mathbf{B}^\top\|^2 \stackrel{\text{def}}{=} \sum_i \|\mathbf{t}[i] - \mathbf{M}\boldsymbol{\lambda}[i]\|^2$$

- One shows that [COM04a] [YER02]

$$\boldsymbol{\lambda}[i] = \{\mathbf{M}^\top \mathbf{M}\}^{-1} \mathbf{M}^\top \mathbf{t}[i]$$

and each column of  $\mathbf{B}$  is the dominant eigenvector of the real symmetric matrix:

$$\mathbf{P}[\ell] = \frac{1}{2} \sum_k \lambda_\ell[k] \{ \tilde{\mathbf{T}}[k; \ell]^\top + \tilde{\mathbf{T}}[k; \ell] \}$$

where  $\tilde{\mathbf{T}}[k; \ell] \stackrel{\text{def}}{=} \mathbf{T}[k] - \sum_{n \neq \ell} \lambda_n[k] \mathbf{b}[n] \mathbf{b}[n]^\top$ .

# ALS drawbacks

- 1 Fairly slow convergence when reaching plateaux
- 2 May be stuck about local minima

# ALS with extrapolation

Attempt to face the first drawback [BA98] [BRO97].

- Compute stationary values  $\hat{\mathbf{A}}$ ,  $\hat{\mathbf{B}}$  and  $\hat{\mathbf{C}}$  as in page 306
- At every other iteration, set:

$$\mathbf{A}(k+1) = \hat{\mathbf{A}} + \mu(k)(\mathbf{A}(k) - \hat{\mathbf{A}})$$

$$\mathbf{B}(k+1) = \hat{\mathbf{B}} + \mu(k)(\mathbf{B}(k) - \hat{\mathbf{B}})$$

$$\mathbf{C}(k+1) = \hat{\mathbf{C}} + \mu(k)(\mathbf{C}(k) - \hat{\mathbf{C}})$$

where one may take  $\mu(k) = k^{1/3}$ .

- and otherwise  $\mathbf{A}(k+1) = \hat{\mathbf{A}}$ ,  $\mathbf{B}(k+1) = \hat{\mathbf{B}}$  and  $\mathbf{C}(k+1) = \hat{\mathbf{C}}$ .



# ALS with Enhanced Line Search (ELS)

Attempt to face both drawbacks [RCH08] [RC05]

- Compute stationary values  $\hat{\mathbf{A}}$ ,  $\hat{\mathbf{B}}$  and  $\hat{\mathbf{C}}$  as in page 306
- At every other iteration, set:

$$\mathbf{A}(k+1) = \hat{\mathbf{A}} + \mu (\mathbf{A}(k) - \hat{\mathbf{A}})$$

$$\mathbf{B}(k+1) = \hat{\mathbf{B}} + \mu (\mathbf{B}(k) - \hat{\mathbf{B}})$$

$$\mathbf{C}(k+1) = \hat{\mathbf{C}} + \mu (\mathbf{C}(k) - \hat{\mathbf{C}})$$

where  $\mu = \text{Arg min}_{\mu} \|\mathbf{T} - \mathbf{A}(k+1) \bullet \mathbf{B}(k+1) \bullet \mathbf{C}(k+1)\|^2$ .

- and otherwise  $\mathbf{A}(k+1) = \hat{\mathbf{A}}$ ,  $\mathbf{B}(k+1) = \hat{\mathbf{B}}$  and  $\mathbf{C}(k+1) = \hat{\mathbf{C}}$ .

**NB:**  $\mu(k)$  is obtained by rooting a polynomial of degree 5.  $\Rightarrow$  one gets the *absolute minimum* along the search direction  $\Rightarrow$  increased capability to escape from local minima.

## ELS applied to other iterative algorithms

The same principle applies to any iterative algorithm [?]:

- Compute a search direction  $[\Delta \mathbf{A}, \Delta \mathbf{B}, \Delta \mathbf{C}]$ , which can be the gradient  $\mathbf{g}$ , a direction  $\mathbf{H}^{-1}\mathbf{g}$ , or a difference  $[\hat{\mathbf{A}}, \hat{\mathbf{B}}, \hat{\mathbf{C}}] - [\mathbf{A}(k), \mathbf{B}(k), \mathbf{C}(k)]...$
- Compute the 6 first coefficients of the 6th degree polynomial  $\varepsilon(\mu)$ , defined by replacing  $[\mathbf{A}, \mathbf{B}, \mathbf{C}]$  by  $[\mathbf{A} + \mu \delta \mathbf{A}, \mathbf{B} + \mu \delta \mathbf{B}, \mathbf{C} + \mu \delta \mathbf{C}]$
- Compute the 5 roots of its derivative
- Select the root  $\mu_o$  yielding the smallest minimum of  $\varepsilon(\mu)$
- Update:  $\mathbf{A}(k+1) = \mathbf{A}(k) + \mu_o \delta \mathbf{A}$ ,  
 $\mathbf{B}(k+1) = \mathbf{B}(k) + \mu_o \delta \mathbf{B}$ ,  $\mathbf{C}(k+1) = \mathbf{C}(k) + \mu_o \delta \mathbf{C}$ .

Can be executed at every iteration, or less often.

# Definition of c.f.'s

## Characteristic functions

First:  $\Phi_x(\mathbf{u}) \stackrel{\text{def}}{=} E\{\exp(j \mathbf{u}^T \mathbf{x})\}$

Second:  $\Psi_x(\mathbf{u}) \stackrel{\text{def}}{=} \log \Phi_x(\mathbf{u})$

## Generating functions

First:  $\Phi_x(\mathbf{u}) \stackrel{\text{def}}{=} E\{\exp(\mathbf{u}^T \mathbf{x})\}$

Second:  $\Psi_x(\mathbf{u}) \stackrel{\text{def}}{=} \log \Phi_x(\mathbf{u})$

## Key property

If  $\mathbf{s}$  has statistically independent components

$$\Psi_s(\mathbf{u}) = \sum_p \Psi_{s_p}(u_p)$$

# Characteristic function of a linear mixture

- If  $s_p$  independent,  $E\{\prod_p f(s_p)\} = \prod_p E\{f(s_p)\}$ .
- Hence if  $\mathbf{x} = \mathbf{H}\mathbf{s}$ , then

$$\begin{aligned}\Phi_{\mathbf{x}}(\mathbf{u}) &\stackrel{\text{def}}{=} E\{\exp(\mathbf{u}^T \mathbf{H} \mathbf{s})\} = E\{\exp(\sum_{p,q} u_q H_{qp} s_p)\} \\ &= \prod_p E\{\exp(\sum_q u_q H_{qp} s_p)\}\end{aligned}$$

- Thus we have the *core equation*:

$$\psi_s(\mathbf{u}) = \sum_p \psi_{s_p} \left( \sum_q u_q H_{qp} \right)$$

# Putting the problem in tensor form (1)

**Goal:** Find a matrix  $\mathbf{H}$  such that the  $K$ -variate function  $\Psi_{\mathbf{x}}(\mathbf{u})$  decomposes into a sum of  $P$  univariate functions  $\psi_p \stackrel{\text{def}}{=} \Psi_{s_p}$ .

- Assumption: functions  $\psi_p$ ,  $1 \leq p \leq P$  admit finite derivatives up to order  $r$  in a neighborhood of the origin.
- Then, Taking  $r = 3$  as a working example:

$$\frac{\partial^3 \Psi_{\mathbf{x}}}{\partial u_i \partial u_j \partial u_k}(\mathbf{u}) = \sum_{p=1}^P H_{ip} H_{jp} H_{kp} \psi_p^{(3)}\left(\sum_{q=1}^K u_q H_{qp}\right)$$

# Putting the problem in tensor form (1)

## Several equivalent writings:

- A decomposition into a sum of rank-1 terms:

$$T_{ijk\ell} = \sum_p H_{ip} H_{jp} H_{kp} B_{\ell p}$$

- A joint diagonalization of matrix slices via a common rectangular transform

$$\mathbf{T}[k, \ell] = \mathbf{H} \cdot \mathbf{Diag}\{\mathbf{H}(k, :)\} \mathbf{Diag}\{\mathbf{B}(\ell, :)\} \cdot \mathbf{H}^T$$

- The cumulant tensor case: only one point  $\mathbf{u} = 0$ , i.e.  $\ell = 1$  and matrix  $\mathbf{B}$  disappears.

# Putting the problem in tensor form (3)

## Use of several orders simultaneously:

- Order 3:

$$T_{ijkl}^{(3)} = \sum_p H_{ip} H_{jp} H_{kp} B_{\ell p}$$

- Order 4:

$$T_{ijkml}^{(4)} = \sum_p H_{ip} H_{jp} H_{kp} H_{mp} C_{\ell p}$$

- Orders 3 and 4:

$$T_{ijkl}[m] = \sum_p H_{ip} H_{jp} H_{kp} D_{\ell p}[m]$$

with  $D_{\ell p}[m] = H_{mp} C_{\ell p}$  and  $D_{\ell p}[0] = B_{\ell p}$ .

# BIOME algorithms

- These algorithms work with a cumulant tensor of even order  $2r > 4$
- We take the case  $2r = 6$  for the presentation, and denote

$$\mathcal{C}_{ijk}^{\ell mn} \stackrel{\text{def}}{=} \text{Cum}\{x_i, x_j, x_k, x_l^*, x_m^*, x_n^*\} \quad (48)$$

- In that case, we have

$$\mathcal{C}_{x,ijk}^{\ell mn} = \sum_{p=1}^P H_{ip} H_{jp} H_{kp} H_{\ell p}^* H_{mp}^* H_{np}^* \Delta_p$$

where  $\Delta_p^{(6)} \stackrel{\text{def}}{=} \text{Cum}\{s_p, s_p, s_p, s_p^*, s_p^*, s_p^*\}$  denote the diagonal entries of a  $P \times P$  diagonal matrix,  $\mathbf{\Delta}^{(6)}$



## Writing in matrix form

- Tensor  $\mathcal{C}_x$  is of dimensions  $K \times K \times K \times K \times K \times K$  and enjoys symmetries and Hermitian symmetries.
- Tensor  $\mathcal{C}_x$  can be stored in a  $K^3 \times K^3$  Hermitian matrix,  $\mathbf{C}_x^{(6)}$ , called the *hexacovariance*. With an appropriate storage of the tensor entries, we have

$$\mathbf{C}_x^{(6)} = \mathbf{H}^{\odot 3} \mathbf{\Delta}^{(6)} \mathbf{H}^{\odot 3H} \quad (49)$$

- Because  $\mathbf{C}_x^{(6)}$  is Hermitian,  $\exists \mathbf{V}$  unitary, such that

$$(\mathbf{C}_x^{(6)})^{1/2} = \mathbf{H}^{\odot 3} (\mathbf{\Delta}^{(6)})^{1/2} \mathbf{V} \quad (50)$$

- **Idea:** Use an invariance property existing between blocks of  $(\mathbf{C}_x^{(6)})^{1/2}$ .

## Using the invariance to estimate $\mathbf{V}$

- Cut the  $K^3 \times P$  matrix  $(\mathbf{C}_x^{(6)})^{1/2}$  into  $K$  blocks of size  $K^2 \times P$ .
- Each of these blocks,  $\mathbf{\Gamma}[n]$ , satisfies:

$$\mathbf{\Gamma}[n] = (\mathbf{H} \odot \mathbf{H}^H) \mathbf{D}[n] (\mathbf{\Delta}^{(6)})^{1/2} \mathbf{V}$$

where  $\mathbf{D}[n]$  is the  $P \times P$  diagonal matrix containing the  $n$ th row of  $\mathbf{H}$ ,  $1 \leq n \leq K$ .

- Hence matrices  $\mathbf{\Gamma}[n]$  share the same common right singular space
- **Algorithm:** compute the joint EVD of the  $K(K-1)$  matrices

$$\mathbf{\Theta}[m, n] \stackrel{\text{def}}{=} \mathbf{\Gamma}[m]^\dagger \mathbf{\Gamma}[n]$$

as:  $\mathbf{\Theta}[m, n] = \mathbf{V} \mathbf{\Lambda}[m, n] \mathbf{V}^H$ .

# Estimation of $\mathbf{H}$

Matrices  $\mathbf{\Lambda}[m, n]$  cannot be used directly because  $(\mathbf{\Delta}^{(6)})^{1/2}$  is unknown. But we use  $\mathbf{V}$  to obtain the estimate of  $\mathbf{H}^{\odot 3}$  up to a scale factor:

$$\widehat{\mathbf{H}^{\odot 3}} = (\mathbf{C}_x^{(6)})^{1/2} \mathbf{V} \quad (51)$$

Then several possibilities exist to get  $\mathbf{H}$  from  $\mathbf{H}^{\odot 3}$  [ACCF04]. The best is as follows:

- Build  $K^2$  matrices  $\Xi[m]$  of size  $K \times P$  from conjugates rows of  $\widehat{\mathbf{H}^{\odot 3}}$
- From  $\Xi[m]$  find matrices  $\mathbf{D}[m]$  and  $\widehat{\mathbf{H}}$  in the LS sense:

xx

# Conditions of identifiability

- $X_X$  [ACCF04] [AFCC03]
- $X_X$

# FOOBI algorithms

XX

XX

■  $X_x$

■  $X_x$

$XX$ 

- $X_x VI$
- $X_x$

## Part VII

# Conclusions



# False beliefs

- 1 BSS always requires High-Order Statistics (HOS)  
→ *Second-order can (rarely) suffice*
- 2 Sources must be statistically independent  
→ *Correlated sources can be sometimes separated (e.g. Discrete/CM sources, Pairwise cumulants...)*
- 3 HOS are always required when sources are *i.i.d.*  
→ *Second-order BSS algorithms exist*
- 4 Even local maxima of a contrast function yield good solutions  
→ *sometimes local maxima correspond to bad solutions*
- 5 There should be at least as many sensors as sources:  $K \geq P$  (sufficient diversity)  
→ *Underdetermined mixtures can be identified*

## False beliefs (cont'd)

- 6 Perfect source extraction is impossible if  $K < P$   
→ *Discrete sources can often be perfectly extracted from under-determined mixtures (insufficient diversity)*
- 7 Conditions of application of Parafac are mild  
→ *except when one dimension = 2, the typical rank always exceeds the Parafac bound for uniqueness*
- 8 Approximate a tensor by another of lower rank is as easy as for matrices  
→ *beside for rank 1, there is a lack of closeness*
- 9 The Constant Modulus (CM) property is the best way to handle PSK sources  
→ *The whole alphabet can be taken into account in order to define a contrast function*

# Part VIII

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